

SCHOOL OF ARTIFICIAL INTELLIGENCE, NANJING UNIVERSITY



Stochastic Processes 随机过程

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Email: qianc@nju.edu.cn Homepage: http://www.lamda.nju.edu.cn/qianc/ A stochastic process is a collection $\{X(t) \mid t \in T\}$ of random variables

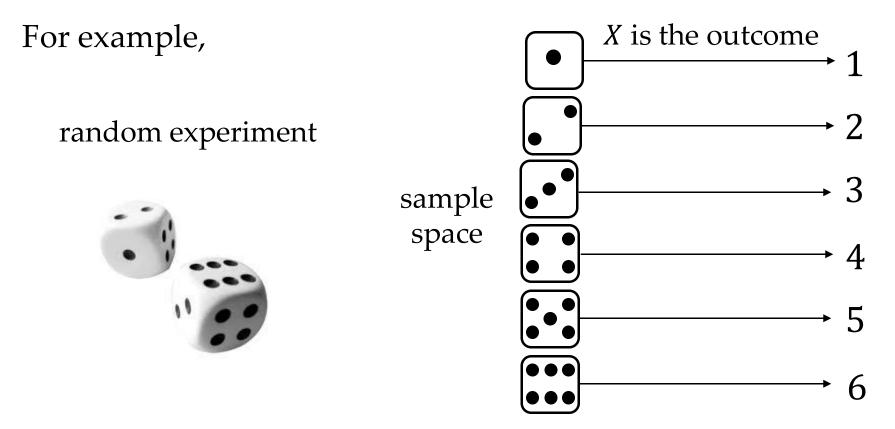
- *X*(*t*) is a random variable
- *t* is often interpreted as time, and *X*(*t*) is called the state of the process at time *t*
- Discrete-time stochastic process:

The index set *T* is a countable set

Continuous-time stochastic process:

The index set *T* is a continuum

A random variable $X: S \rightarrow R$ is a function that assigns a real value to each outcome in the sample space *S*

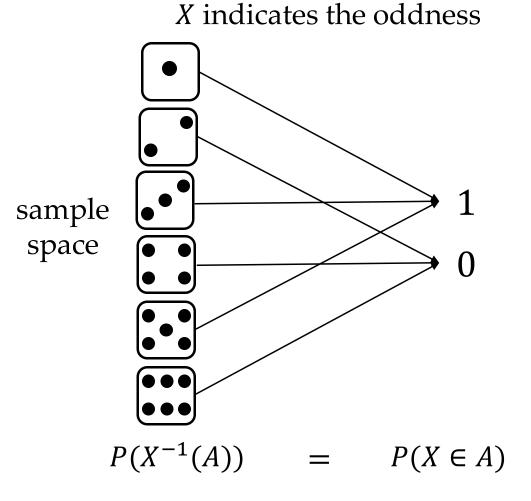


Random variable

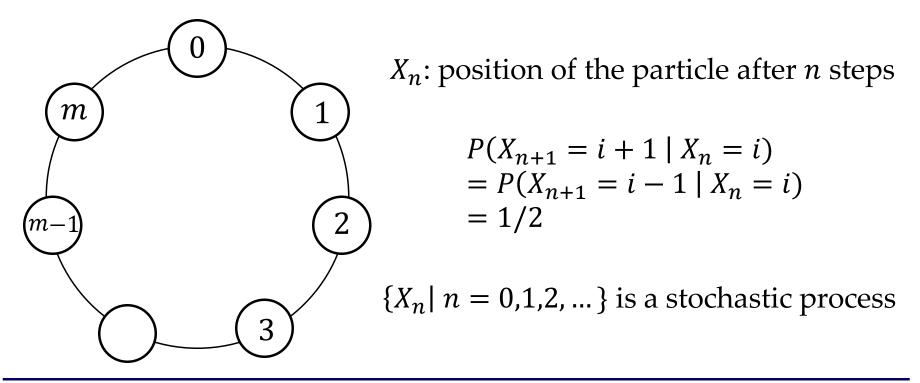
For example,

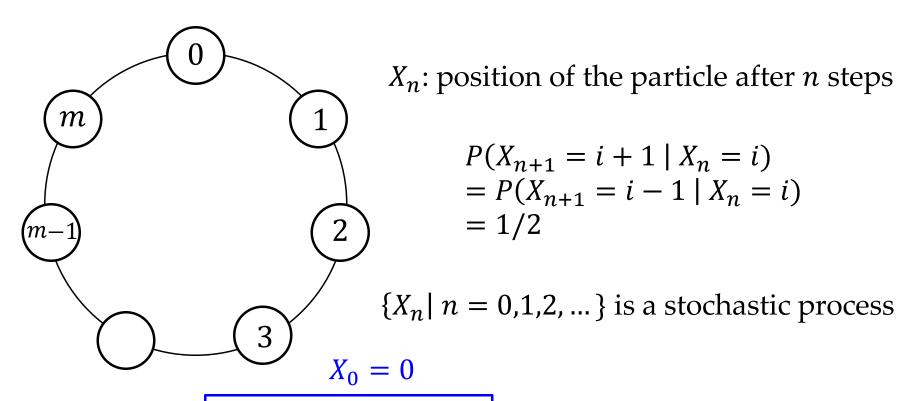
random experiment





Example: Consider a particle that moves along a set of m + 1 nodes, labelled 0, 1, ..., m, that are arranged around a circle. At each step the particle is equally likely to move one position in either the clockwise or counterclockwise direction.





Suppose that the particle starts at 0 and continues to move around according to the above rules until all the nodes have been visited.

What is the probability that node *i* is the last one visited?

Solution:

Target event *E*_{*i*}: *i* is the last visited node until visiting all nodes

Random variable *T_k***:** the first time that the particle visits *k*

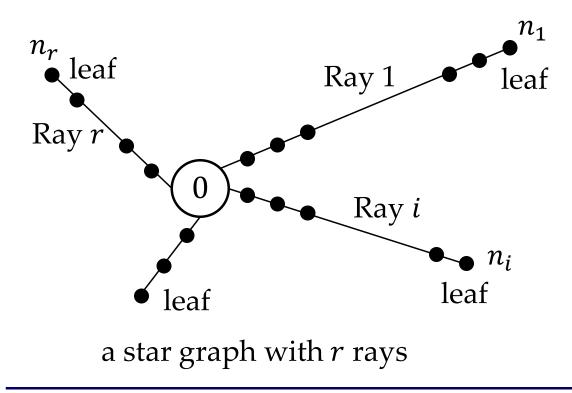
$$P(E_i) = P(E_i | T_{i-1} < T_{i+1}) P(T_{i-1} < T_{i+1}) + P(E_i | T_{i-1} > T_{i+1}) P(T_{i-1} > T_{i+1})$$

before node *i* is visited, $= P(W_m)P(T_{i-1} < T_{i+1}) + P(W_m)P(T_{i-1} > T_{i+1})$ node *i* + 1 is visidted $= P(W_m)$

$$\sum_{i=1}^{m} P(E_i) = 1 \implies P(E_i) = \frac{1}{m}$$

Event W_m : a gambler who starts with 1 unit, and wins 1 when a fair coin turns up heads and loses 1 when it turns up tails, will have his fortune go up by m - 1 before he goes broke

Example: Consider a particle moves along the vertices of the graph so that it is equally likely to move from whichever vertex it is presently at to any of the neighbors of that vertex



 X_n : position of the particle after *n* steps

 $\{X_n | n = 0, 1, 2, ...\}$ is a stochastic process

Starting at 0, what is the probability that the first visited leaf is on ray *i*?

Solution:

Target event *E_i*: the first visited leaf is on ray *i* **Event** *C_k*: the first visited ray is *k* $P(E_i) = \sum_{k=1}^{r} P(E_i | C_k) \frac{1}{r}$ $P(E_i | C_i) = P(W_{n_i}) + \left(1 - P(W_{n_i})\right)P(E_i) = \frac{1}{n_i} + \left(1 - \frac{1}{n_i}\right)P(E_i)$ $\forall j \neq i: P(E_i | C_j) = 0 + \left(1 - P(W_{n_j})\right)P(E_i) = \left(1 - \frac{1}{n_j}\right)P(E_i)$ $P(E_i)r = \frac{1}{n_i} + \sum_{k=1}^r \left(1 - \frac{1}{n_k}\right) P(E_i) \implies P(E_i) = \frac{\frac{1}{n_i}}{\sum_{k=1}^r \frac{1}{n_k}}$

Event W_m : a gambler who starts with 1 unit, and wins 1 when a fair coin turns up heads and loses 1 when it turns up tails, will have his fortune go up by m - 1 before he goes broke

A stochastic process $\{X(t) \mid t \in T\}$ is said to have independent increments if $\forall t_0 < t_1 < t_2 < \cdots < t_n$, the random variables

$$X(t_1) - X(t_0), X(t_2) - X(t_1), \dots, X(t_n) - X(t_{n-1})$$

are independent

A stochastic process $\{X(t) \mid t \in T\}$ is said to have stationary increments if $\forall s > 0$,

$$X(t+s) - X(t)$$

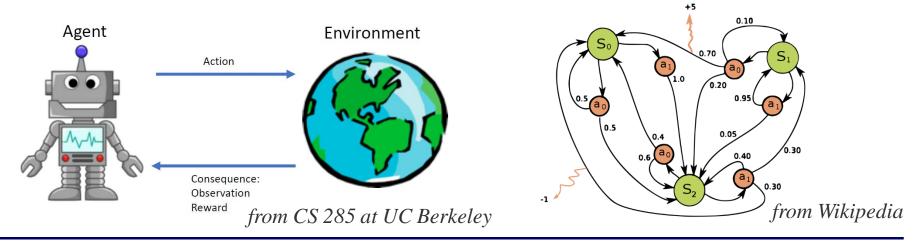
has the same distribution for all *t*

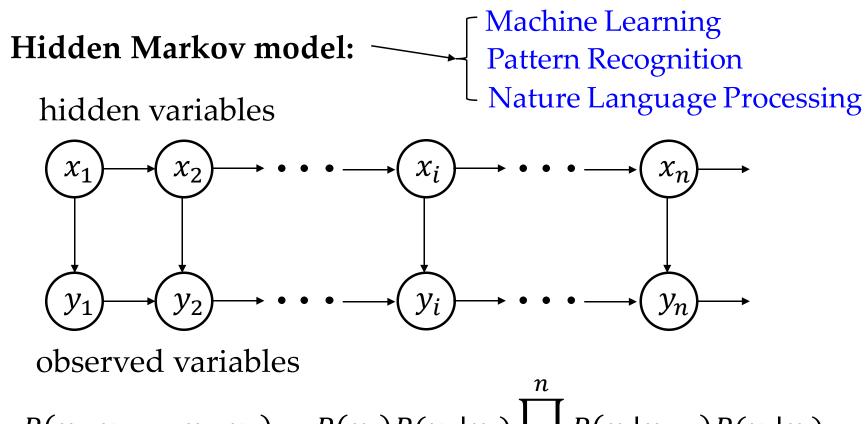
Markov decision process in reinforcement learning:

• State space *S*

Stochastic Process

- Action space *A*
- Transition function *P*(*s*' | *s*, *a*) : the probability of transitioning into state *s*' upon taking action *a* in state *s*
- Reward function *R*(*s*, *a*): the immediate reward associated with taking action *a* in state *s*

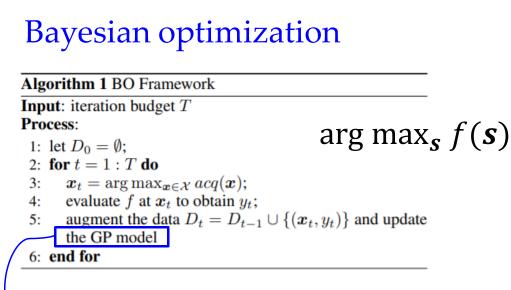




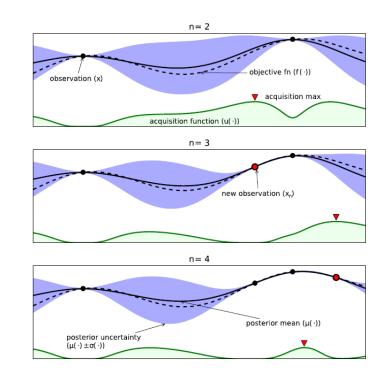
$$P(x_1, y_1, \dots, x_n, y_n) = P(x_1)P(y_1|x_1) \prod_{i=2}^{n} P(x_i|x_{i-1})P(y_i|x_i)$$

 $\{x_n | n \ge 1\}$ and $\{y_n | n \ge 1\}$ are two stochastic processes

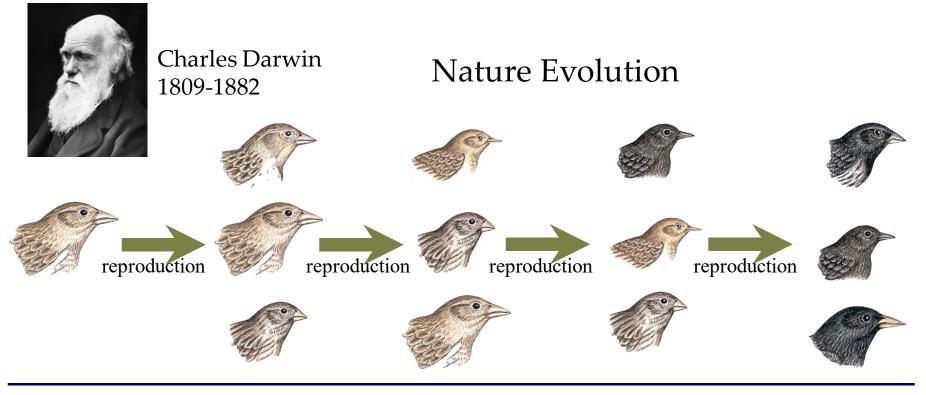
Gaussian process (GP) is a collection of random variables, such that every finite collection of those random variables has a multivariate normal distribution <u>Stochastic Process</u>



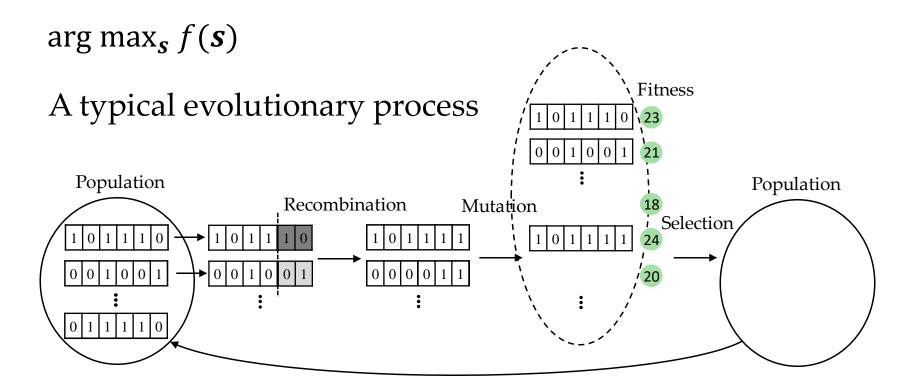
regards the *f* value at each data point as a random variable, and assumes satisfying a joint Gaussian distribution



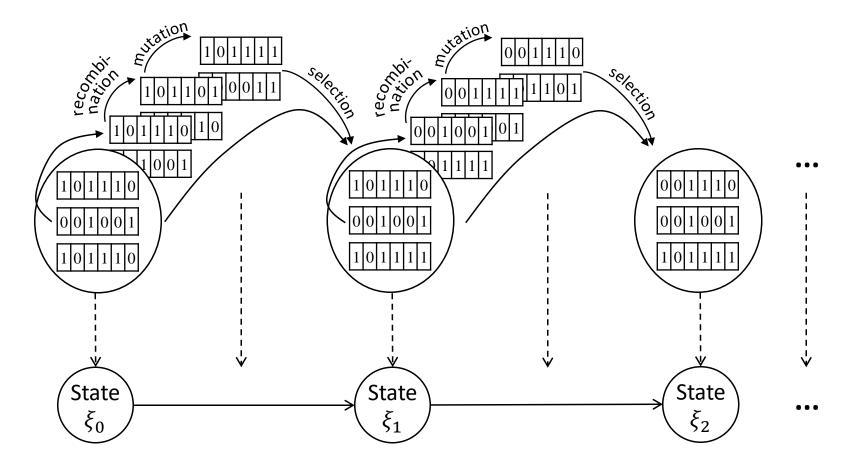
Evolutionary algorithms (EAs) are a kind of randomized heuristic optimization algorithms, inspired by nature evolution (reproduction with variation + nature selection)



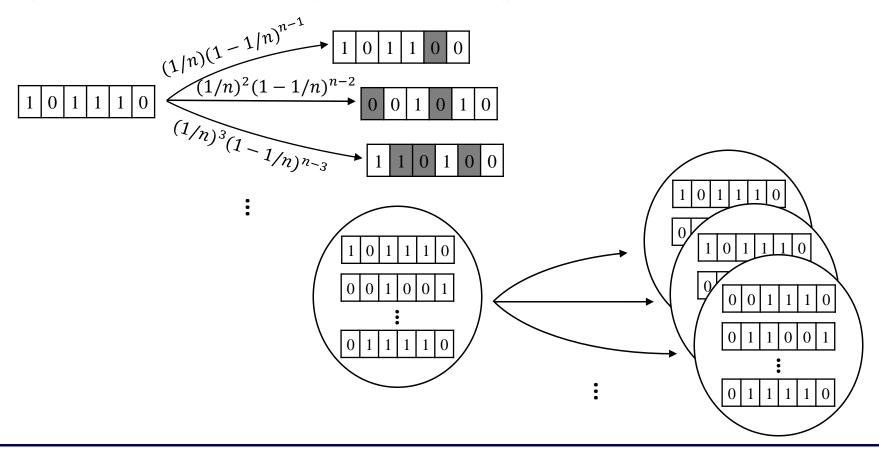
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Evolutionary algorithms

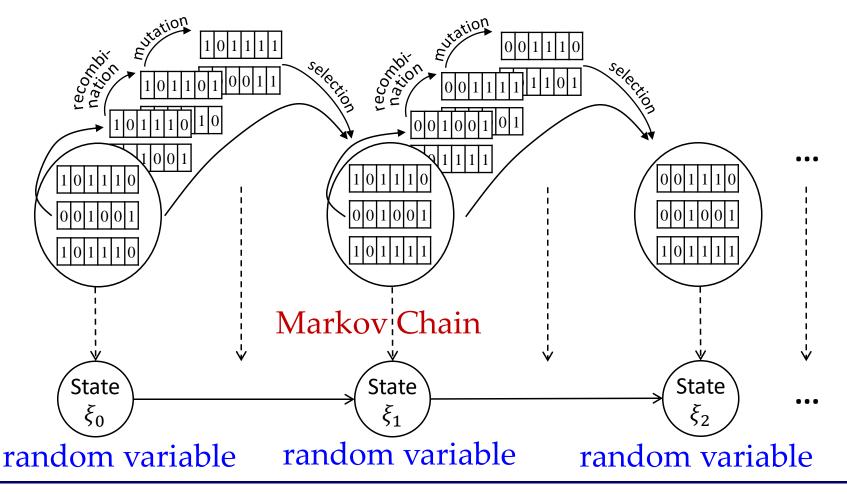


The process of generating new population is randomized e.g., bit-wise mutation: flips each bit with prob. 1/n



Evolutionary algorithms

Stochastic Process



- Lecture 1: Preliminaries
- Lecture 2: Poisson process
- Lecture 3: Renewal process
- Lecture 4: Markov chain
- Lecture 5: Martingale
- Lecture 6: Random walk
- **Lecture 7: Brownian motion**



- An Introduction to Stochastic Modeling, 4th edition, 2010 by Mark Pinsky and Samuel Karlin
- Basic Stochastic Processes, 1999 by Zdzislaw Brzezniak and Tomasz Zastawniak
- Stochastic Processes, 2nd edition, 1995 by Sheldon M. Ross
- Introduction to Stochastic Processes, 2013 *by Erhan Cinlar*
- Essentials of Stochastic Processes, 2nd edition, 2012 by Richard Durrett



课程时间:周五下午14:00-16:00 课程主页: http://www.lamda.nju.edu.cn/SP22/ 课程讨论QQ群: 569309099 随机过程 不随机过 助教: 卞超、王雨桐 答疑时间:周五下午16:00-17:30、逸A-502

成绩计算:4次平时作业(15%)、期末考试(40%)



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Stochastic Processes Lecture 1: Preliminaries

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Email: qianc@nju.edu.cn Homepage: http://www.lamda.nju.edu.cn/qianc/ Two random variables *X* and *Y* are said to be **independent** if

$$P(X = x, Y = y) = P(X = x) \cdot P(Y = y)$$

for all *x* and *y*

Random variables $X_1, X_2, ..., X_n$ are said to be **mutually independent** if for any subset $I \subseteq \{1, 2, ..., n\}$ and any x_i ,

$$P\left(\bigwedge_{i\in I} (X_i = x_i)\right) = \prod_{i\in I} P(X_i = x_i)$$

Expectation of random variables

The **expectation** of a discrete random variable *X* is

$$E[X] = \sum_{x} x \cdot P(X = x)$$

where the sum is over all *x* in the range of *X*

Two common ways of calculating *E*[*X*]:

• Let
$$X = X_1 + X_2 + \dots + X_n$$
, then $E[X] = \sum_{i=1}^n E[X_i]$

• E[X] = E[E[X | Y]]

How to calculate the expectation

Let
$$X = X_1 + X_2 + \dots + X_n$$
, then $E[X] = \sum_{i=1}^n E[X_i]$

Proof:

$$E[X + Y] = \sum_{x} \sum_{y} (x + y)P(X = x, Y = y)$$

= $\sum_{x} x \sum_{y} P(X = x, Y = y) + \sum_{y} y \sum_{x} P(x = x, Y = y)$
= $\sum_{x} x P(X = x) + \sum_{y} y P(Y = y)$
= $E[X] + E[Y]$
 \downarrow
 $E[X_{1} + \dots + X_{n}] = E[X_{1} + \dots + X_{n-1}] + E[X_{n}] = \dots = \sum_{i=1}^{n} E[X_{i}]$

How to calculate the expectation

$$E[X] = E[E[X | Y]]$$

Proof:

$$E[E[X | Y]] = \sum_{y} E[X | Y = y]P(Y = y)$$

$$= \sum_{y} \sum_{x} x P(X = x | Y = y)P(Y = y)$$

$$= \sum_{x} x \sum_{y} P(X = x, Y = y)$$

$$= \sum_{x} x P(X = x)$$

$$= E[X]$$

Example

Example: There are *n* keys with the same shape, where only one can unlock the door. Each key is selected randomly without replacement. Let *X* denote the number of selected keys until unlocking the door. Calculate E[X].

Solution1: By the definition of expectation

$$E[X] = \sum_{k=1}^{n} k \cdot P(X = k)$$

= $\sum_{k=1}^{n} k \cdot \frac{n-1}{n} \times \frac{n-2}{n-1} \times \dots \times \frac{n-(k-1)}{n-(k-2)} \times \frac{1}{n-(k-1)}$
= $\sum_{k=1}^{n} k \cdot \frac{1}{n} = \frac{n+1}{2}$



Let
$$X = X_1 + X_2 + \dots + X_n$$
, then $E[X] = \sum_{i=1}^n E[X_i]$
Solution2:

Let
$$X_i = \begin{cases} 1 & \text{the first } i - 1 \text{ tries all fail} \\ 0 & \text{otherwise} \end{cases}$$

Then,
$$X_1 = 1$$

 $\forall i \ge 2, E[X_i] = P(X_i = 1) = \frac{n-1}{n} \times \frac{n-2}{n-1} \times \dots \times \frac{n-(i-1)}{n-(i-2)}$
 $= \frac{n-(i-1)}{n}$
Thus, $E[X] = 1 + \sum_{i=2}^{n} \frac{n-(i-1)}{n} = \frac{n+1}{2}$



$$E[X] = E[E[X | Y]]$$

Solution3:

Let $Y = \begin{cases} 1 & \text{the first try succeeds} \\ 0 & \text{otherwise} \end{cases}$,

and X_n denote the random variable X corresponding to n keys

Then,
$$E[X_n] = E[E[X_n | Y]]$$

$$= \frac{1}{n} E[X_n | Y = 1] + (1 - \frac{1}{n}) E[X_n | Y = 0]$$

$$= \frac{1}{n} + (1 - \frac{1}{n})(1 + E[X_{n-1}]) = 1 + (1 - \frac{1}{n}) E[X_{n-1}]$$

$$\implies E[X_n] = \frac{n+1}{2}$$

A discrete random variable *X* is said to have a **Poisson distribution with parameter** λ , $\lambda > 0$, if

$$P(X = k) = \frac{\lambda^k e^{-\lambda}}{k!} \quad \text{for } k = 0, 1, 2, \dots$$

Expectation: $E[X] = \lambda$ Variance: $Var[X] = \lambda$

A discrete random variable *X* is said to have a **Binomial distribution with parameters** n **and** $p \in [0, 1]$, if

$$P(X = k) = {n \choose k} p^k (1 - p)^{n-k}$$
 for $k = 0, 1, 2, ..., n$

Expectation: E[X] = np Variance: Var[X] = np(1-p)

Poisson and binomial distribution

Poisson distribution with parameter $\lambda = np$ can be used as an approximation to binomial distribution with parameters n and p if n is sufficiently large and p is sufficiently small.

Brun's sieve: Let *X* be a bounded nonnegative integer-valued random variable. If, for all $i \ge 0$,

$$E\left[\binom{X}{i}\right] \approx \lambda^i/i!$$

To show that a binomially distributed random variable *X* satisfies this equation

then

$$P(X = j) \approx \frac{\lambda^{j} e^{-\lambda}}{j!}$$
 for $j = 0, 1, 2, ...$ Poisson

Poisson distribution

Brun's sieve: Let *X* be a bounded nonnegative integer-valued random variable. If, for all $i \ge 0$,

Poisson and binomial distribution

To prove that a binomial random variable *X* with parameters *n* and *p* satisfies $\lceil (X) \rceil$

$$E\left[\binom{X}{i}\right] \approx \lambda^i / i!$$

Proof:

- *X* can be viewed as the number of successes in *n* independent trials where each is a success with probability *p*
- For each of the $\binom{n}{i}$ sets of *i* trials, define

$$\forall j \in \left\{1, 2, \dots, \binom{n}{i}\right\}: X_j = \begin{cases}1 \text{ if all the } i \text{ trails are successes}\\0 \text{ otherwise}\end{cases}$$

To prove that a binomial random variable *X* with parameters *n* and *p* satisfies $\lceil (X) \rceil$

$$E\left[\binom{X}{i}\right] \approx \lambda^i/i!$$

Proof:
$$\binom{X}{i} = \sum_{j=1}^{C_n^i} X_j$$

$$E\left[\binom{X}{i}\right] = E\left[\sum_{j=1}^{C_n^i} X_j\right] = \sum_{j=1}^{C_n^i} E\left[X_j\right] = C_n^i \cdot p^i = \frac{n(n-1)\cdots(n-i+1)}{i!} \cdot p^i$$

Case 1: *i* is small enough relative to *n*, e.g. $i \in o(n)$ $E\left[\binom{X}{i}\right] \approx \frac{n^{i}}{i!} \cdot p^{i} = \lambda^{i}/i!$

Case 2: *i* is not small relative to *n*, e.g. $i \in \theta(n)$ $E[\binom{X}{i}] \leq \frac{n^{i}}{i!} \cdot p^{i} \to 0$ holds when *n* is large enough, *p* is small enough and *np* is not large

Poisson and binomial distribution

How to bound the approximation error for general *n* and *p*?

Let $X = \sum_{i=1}^{n} X_i$, where the X_i are Bernoulli random variables with respective means p_i , i = 1, ..., n. Set $\lambda = \sum_{i=1}^{n} p_i$ and let *A* denote a set of nonnegative integers. To bound

$$P(X \in A) - \sum_{i \in A} \frac{\lambda^i e^{-\lambda}}{i!}$$

Remark:

 X_i are not necessarily independent, and p_i can be different

More general than binomial distribution

Poisson and binomial distribution

To bound
$$P(X \in A) - \sum_{i \in A} \frac{\lambda^i e^{-\lambda}}{i!}$$

Proof: Define a function *g* for which

$$P(X \in A) - \sum_{i \in A} \frac{\lambda^i e^{-\lambda}}{i!} = E[\lambda g(X+1) - Xg(X)]$$

Let
$$g(0) = 0$$
, $\forall j \ge 0$: $g(j+1) = \frac{1}{\lambda} \left[I(j \in A) - \sum_{i \in A} \frac{\lambda^i e^{-\lambda}}{i!} + jg(j) \right]$

Then,
$$\lambda g(j+1) - jg(j) = I(j \in A) - \sum_{i \in A} \frac{\lambda^i e^{-\lambda}}{i!}$$

$$\implies E[\lambda g(X+1) - Xg(X)] = E[I(X \in A)] - \sum_{i \in A} \frac{\lambda^i e^{-\lambda}}{i!}$$
$$= P(X \in A) - \sum_{i \in A} \frac{\lambda^i e^{-\lambda}}{i!}$$

To bound
$$P(X \in A) - \sum_{i \in A} \frac{\lambda^i e^{-\lambda}}{i!}$$

Proof:

• Analyze $E[\lambda g(X + 1) - Xg(X)]$

 $E[\lambda g(X+1)] = E[\sum_{i=1}^{n} p_i g(X+1)] = \sum_{i=1}^{n} p_i E[g(X+1)]$ $E[Xg(X)] = \sum_{i=1}^{n} p_i E[g(X) \mid X_i = 1] = \sum_{i=1}^{n} p_i E[g(V_i + 1)]$ following the lemma next page $P(V_i = k) = P(\sum_{j \neq i} X_j = k \mid X_i = 1)$ $|E[\lambda g(X+1)] - E[Xg(X)]| = |\sum_{i=1}^{n} p_i (E[g(X+1)] - E[g(V_i + 1)])|$

 $\leq \sum_{i=1}^{n} p_i E[|g(X+1) - g(V_i+1)|]$

Lemma: For any random variable *R*,

$$E[XR] = \sum_{i=1}^{n} p_i E[R \mid X_i = 1]$$

]

Proof:

$$E[XR] = E\left[\sum_{i=1}^{n} R X_i\right] = \sum_{i=1}^{n} E[RX_i]$$
$$= \sum_{i=1}^{n} E[E[RX_i \mid X_i]]$$
$$= \sum_{i=1}^{n} p_i E[R \mid X_i = 1]$$

$$\left| P(X \in A) - \sum_{i \in A} \frac{\lambda^{i} e^{-\lambda}}{i!} \right| \le \sum_{i=1}^{n} p_{i} E[|g(X+1) - g(V_{i}+1)|]$$

Lemma: For any λ and A,

$$\begin{aligned} |g(j) - g(j-1)| &\leq \min\{1, 1/\lambda\} \\ |g(X+1) - g(V_i+1)| &= |g(X+1) - g(X) + \dots + g(V_i+2) - g(V_i+1)| \\ &\leq |g(X+1) - g(X)| + \dots + |g(V_i+2) - g(V_i+1)| \\ &= |X - V_i| \cdot \min\{1, 1/\lambda\} \end{aligned}$$

$$\left| P(X \in A) - \sum_{i \in A} \frac{\lambda^{i} e^{-\lambda}}{i!} \right| \leq \sum_{i=1}^{n} p_{i} E[|g(X+1) - g(V_{i}+1)|]$$
$$\leq \sum_{i=1}^{n} p_{i} \cdot E[|X - V_{i}|] \cdot \min\{1, 1/\lambda\}$$

How to bound the approximation error for general *n* and *p*?

Let $X = \sum_{i=1}^{n} X_i$, where the X_i are Bernoulli random variables with respective means p_i , i = 1, ..., n. Set $\lambda = \sum_{i=1}^{n} p_i$ and let *A* denote a set of nonnegative integers.

$$\begin{split} P(X \in A) &- \sum_{i \in A} \frac{\lambda^i e^{-\lambda}}{i!} \middle| \le \sum_{i=1}^n p_i \cdot \min\{1, 1/\lambda\} \cdot E[|X - V_i|] \\ & \text{where } P(V_i = k) = P(\sum_{j \neq i} X_j = k \mid X_i = 1) \end{split}$$

Remark:

 X_i are not necessarily independent, and p_i can be different

How to bound the approximation error for general *n* and *p*?

For binomial random variable *X*,

 X_i are independent, and p_i are the same, denoted as p

$$\begin{vmatrix} P(X \in A) - \sum_{i \in A} \frac{\lambda^i e^{-\lambda}}{i!} \end{vmatrix} \le \sum_{i=1}^n p_i \cdot \min\{1, 1/\lambda\} \cdot \underbrace{E[|X - V_i|]}_{\text{by independence}} \\ = \min\{np^2, p\} \qquad E[|X_i|] = p_i \end{aligned}$$

Exponential distribution

A continuous random variable *X* is said to have an **exponential distribution with parameter** λ , $\lambda > 0$, if its *probability density function* is given by

$$f(x) = \begin{cases} \lambda e^{-\lambda x} & x \ge 0\\ 0 & x < 0 \end{cases}$$

or, equivalently, if its *cumulative distribution function* is

$$F(x) = \begin{cases} 1 - e^{-\lambda x} & x \ge 0\\ 0 & x < 0 \end{cases}$$

Expectation: $E[X] = 1/\lambda$ Variance: $Var[X] = 1/\lambda^2$

An exponentially distributed random variable *X* has the **memoryless** property:

$$\forall s, t \ge 0, P(X > s + t | X > t) = P(X > s)$$

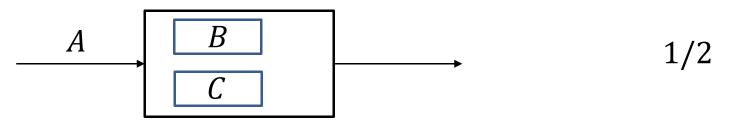
Proof:

$$P(X > s + t \mid X > t) = \frac{P(X > s + t, X > t)}{P(X > t)} = \frac{P(X > s + t)}{P(X > t)}$$
$$= \frac{e^{-\lambda(s+t)}}{e^{-\lambda t}} = e^{-\lambda s} = P(X > s)$$

An exponentially distributed random variable *X* has the **memoryless** property:

 $\forall s, t \ge 0, P(X > s + t | X > t) = P(X > s)$

Application: Consider a post office having two clerks, and suppose that when *A* enters the system, *B* and *C* are being served by these two clerks, respectively. Suppose also that *A* will begin to be served once either *B* or *C* leaves. If the amount of time a clerk spends with a customer is exponentially distributed with parameter λ , what is the probability that *A* is the last to leave the post office?



Failure (Hazard) rate function

Consider a non-negative, continuous random variable *X* having distribution *F* and density *f*, let $\overline{F}(x) = P(X > x) = 1 - F(x)$.

The failure (or hazard) rate function $\lambda(t)$ is defined by

$$\lambda(t) = \frac{f(t)}{\overline{F}(t)}$$

Intuitive explanation: Let *X* denote the lifetime of some item

$$P(X \in (t, t + dt) \mid X > t) = \frac{P(x \in (t, t + dt), X > t)}{P(X > t)}$$
$$\approx \frac{f(t) \cdot dt}{\overline{F}(t)} = \lambda(t) \cdot dt$$

Failure rate function $\lambda(t)$ for the exponential distribution:

$$\lambda(t) = \frac{f(t)}{\overline{F}(t)} = \frac{\lambda e^{-\lambda t}}{e^{-\lambda t}} = \lambda$$

The failure rate function $\lambda(t)$ uniquely determines the cumulative distribution *F*

Proof:

By definition of
$$\lambda(t)$$
, $\lambda(t) = \frac{f(t)}{\overline{F}(t)} = \frac{-\frac{dT(t)}{dt}}{\overline{F}(t)}$
Then, $\int_0^t -\lambda(t)dt = \log \overline{F}(t) \Big|_0^t = \log \overline{F}(t) - \log \overline{F}(0)$
Thus, $\overline{F}(t) = e^{-\int_0^t \lambda(t)dt}$

 $d\overline{E}(t)$

Markov inequality: If *X* is a nonnegative random variable, then for any a > 0,

$$P(X \ge a) \le \frac{E[X]}{a}$$

Proof:

Let
$$Y = \begin{cases} 1 & \text{if } X \ge a \\ 0 & \text{otherwise} \end{cases}$$
, then $Y \le \frac{X}{a}$

Thus,
$$E[Y] \leq \frac{E[X]}{a}$$

Chernoff bound: If *X* is random variable, then for a > 0,

$$P(X \ge a) \le e^{-ta} E[e^{tX}] \text{ for all } t > 0$$
$$P(X \le a) \le e^{-ta} E[e^{tX}] \text{ for all } t < 0$$
Proof:

For all
$$t > 0$$
: $P(X \ge a) = P(tX \ge ta) = P(e^{tX} \ge e^{ta})$
by Markov inequality $\swarrow \subseteq \frac{E[e^{tX}]}{e^{ta}}$
For all $t < 0$: $P(X \le a) = P(tX \ge ta) = P(e^{tX} \ge e^{ta})$
 $\leq \frac{E[e^{tX}]}{e^{ta}}$

Chernoff bound: If *X* is random variable, then for a > 0,

$$P(X \ge a) \le e^{-ta} E[e^{tX}] \text{ for all } t > 0$$
$$P(X \le a) \le e^{-ta} E[e^{tX}] \text{ for all } t < 0$$

Application: If *X* is Poisson with mean λ , i.e.,

$$P(X = k) = \frac{\lambda^k e^{-\lambda}}{k!}$$
 for $k = 0, 1, 2, ...$

$$\begin{split} P(X \ge j) \le e^{-tj} E[e^{t\alpha}] &= \sum_{k=0}^{\infty} e^{tk} \cdot \frac{X^k e^{-\lambda}}{k!} & (\lambda e^t - j) \cdot e^{\lambda (e^t - 1) - tj} \\ &= e^{-\lambda} \sum_{k=0}^{\infty} \frac{(\lambda e^t)^k}{k!} = e^{-\lambda} \cdot e^{\lambda e^t} = \boxed{e^{\lambda (e^t - 1) - tj}} \end{split}$$

Let $t = \ln \frac{j}{\lambda}$ (assume $j > \lambda$), then $P(X \ge j) \le e^{j-\lambda} \cdot \left(\frac{j}{\lambda}\right)^{-j} = \frac{e^{-\lambda}(e\lambda)^j}{j^j}$

Jensen's inequality: If *f* is a convex function, then $E[f(X)] \ge f(E[X])$

provided the expectations exist

Proof:

Let
$$\mu = E[X]$$

Then, $f(X) = f(\mu) + f'(\mu)(X - \mu) + \frac{f''(\xi)}{2}(X - \xi)^2$
 $\ge f(\mu) + f'(\mu)(X - u)$

Thus, $E(f(X)] \ge f(\mu) + f'(\mu)(E[X] - \mu) = f(\mu) = f(E[X])$

Strong Law of Large Numbers: If $X_1, X_2, ...$ are independent and identically distributed with mean μ , then

$$P\left(\lim_{n\to\infty}(X_1+\cdots+X_n)/n=\mu\right)=1$$

Central Limit Theorem: If $X_1, X_2, ...$ are independent and identically distributed with mean μ and variance σ^2 , then

$$\lim_{n \to \infty} P\left(\frac{X_1 + \dots + X_n - n\mu}{\sigma\sqrt{n}} \le a\right) = \int_{-\infty}^a \frac{1}{\sqrt{2\pi}} e^{-x^2/2} dx$$



- What is stochastic process
- Stochastic process in AI
- Preliminaries

References: Chapter 1 & 10, Stochastic Processes, 2nd edition, 1995, *by Sheldon M. Ross*