

Convex optimization problems (II)

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Outline

- Linear Optimization Problems
- Quadratic Optimization Problems
- Geometric Programming
- Generalized Inequality Constraints
- Vector Optimization



Linear Optimization Problems

□ Linear Program (LP)

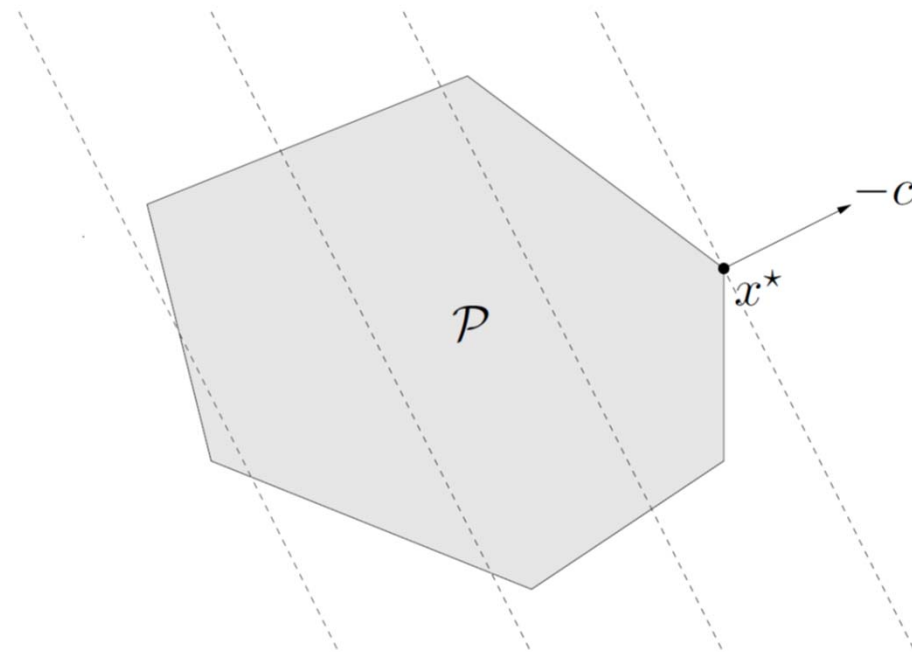
$$\begin{array}{ll} \min & c^T x + d \\ \text{s. t.} & Gx \preceq h \\ & Ax = b \end{array}$$

- $G \in \mathbf{R}^{m \times n}$ and $A \in \mathbf{R}^{p \times n}$
- It is common to omit the constant d
- Maximization problem with affine objective and constraint functions is also an LP
- The feasible set of LP is a polyhedron \mathcal{P}



Linear Optimization Problems

□ Geometric Interpretation of an LP



- The objective $c^T x$ is linear, so its level curves are hyperplanes orthogonal to c
- x^* is as far as possible in the direction $-c$



Two Special Cases of LP

□ Standard Form LP

$$\begin{array}{ll} \min & c^T x \\ \text{s. t.} & Ax = b \\ & x \geq 0 \end{array}$$

- The only inequalities are $x \geq 0$

□ Inequality Form LP

$$\begin{array}{ll} \min & c^T x \\ \text{s. t.} & Ax \leq b \end{array}$$

- No equality constraint



Converting to Standard Form

□ Conversion

$$\begin{array}{ll} \min & c^T x + d \\ \text{s. t.} & Gx \leq h \\ & Ax = b \end{array} \quad \Rightarrow \quad \begin{array}{ll} \min & c^T x \\ \text{s. t.} & Ax = b \\ & x \geq 0 \end{array}$$

- To use an algorithm for standard LP

□ Introduce Slack Variables s

$$\begin{array}{ll} \min & c^T x + d \\ \text{s. t.} & Gx \leq h \\ & Ax = b \end{array} \quad \Rightarrow \quad \begin{array}{ll} \min & c^T x + d \\ \text{s. t.} & Gx + s = h \\ & Ax = b \\ & s \geq 0 \end{array}$$



Converting to Standard Form

□ Decompose x

$$x = x^+ - x^-, \quad x^+, x^- \geq 0$$

□ Standard Form LP

$$\begin{array}{ll} \min & c^\top x + d \\ \text{s. t.} & Gx + s = h \\ & Ax = b \\ & s \geq 0 \end{array} \quad \rightarrow \quad \begin{array}{ll} \min & c^\top x^+ - c^\top x^- + d \\ \text{s. t.} & Gx^+ - Gx^- + s = h \\ & Ax^+ - Ax^- = b \\ & x^+ \geq 0, x^- \geq 0, s \geq 0 \end{array}$$



Example

□ Diet Problem

- Choose nonnegative quantities x_1, \dots, x_n of n foods
- One unit of food j contains amount a_{ij} of nutrient i , and costs c_j
- Healthy diet requires nutrient i in quantities at least b_i
- Determine the cheapest diet that satisfies the nutritional requirements

$$\begin{array}{ll} \min & c^T x \\ \text{s. t.} & Ax \geq b \\ & x \geq 0 \end{array}$$



Example

□ Chebyshev Center of a Polyhedron

- Find the largest Euclidean ball that lies in the polyhedron

$$\mathcal{P} = \{x \in \mathbf{R}^n \mid a_i^\top x \leq b_i, i = 1, \dots, m\}$$

- The center of the optimal ball is called the **Chebyshev center** of the polyhedron
- Represent the ball as $\mathcal{B} = \{x_c + u \mid \|u\|_2 \leq r\}$
- $x_c \in \mathbf{R}^n$ and r are variables, and we wish to maximize r subject to $\mathcal{B} \in \mathcal{P}$
- $\forall x \in \mathcal{B}, a_i^\top x \leq b_i \Leftrightarrow a_i^\top (x_c + u) \leq b_i, \|u\|_2 \leq r \Leftrightarrow a_i^\top x_c + \sup\{a_i^\top u \mid \|u\|_2 \leq r\} \leq b_i \Leftrightarrow a_i^\top x_c + r\|a_i\|_2 \leq b_i$



Example

□ Chebyshev Center of a Polyhedron

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- $x_c \in \mathbf{R}^n$ and r are variables, and we wish to maximize r subject to $\mathcal{B} \in \mathcal{P}$

$$\begin{aligned} \max \quad & r \\ \text{s.t.} \quad & a_i^\top x_c + r \|a_i\|_2 \leq b_i, \quad i = 1, \dots, m \end{aligned}$$



Example

□ Chebyshev Inequalities

- x is a random variable on $\{u_1, \dots, u_n\} \subseteq \mathbf{R}$
- $p_i = \mathbf{prob}(x = u_i)$, $p \geq 0$, $\mathbf{1}^\top p = 1$
- $\mathbf{E}f(x) = \sum_{i=1}^n p_i f(u_i)$ is a linear function of p
- Prior knowledge is given as

$$\alpha_i \leq a_i^\top p \leq \beta_i, \quad i = 1, \dots, m$$

- To find a lower bound of $\mathbf{E}f_0(x) = a_0^\top p$

$$\min \quad a_0^\top p$$

$$\text{s. t.} \quad p \geq 0, \mathbf{1}^\top p = 1$$

$$\alpha_i \leq a_i^\top p \leq \beta_i, \quad i = 1, \dots, m$$



Example

□ Piecewise-linear Minimization

- Consider the (unconstrained) problem

$$f(x) = \max_{i=1,\dots,m} (a_i^\top x + b_i)$$

- The epigraph problem

$$\begin{aligned} \min \quad & t \\ \text{s. t.} \quad & \max_{i=1,\dots,m} (a_i^\top x + b_i) \leq t \end{aligned}$$

- An LP problem

$$\begin{aligned} \min \quad & t \\ \text{s. t.} \quad & a_i^\top x + b_i \leq t, \quad i = 1, \dots, m \end{aligned}$$



Linear-fractional Programming

□ Linear-fractional Program

$$\begin{aligned} \min \quad & f_0(x) \\ \text{s. t.} \quad & Gx \preceq h \\ & Ax = b \end{aligned}$$

- The objective function is a ratio of affine functions

$$f_0(x) = \frac{c^\top x + d}{e^\top x + f}$$

- The domain is

$$\text{dom } f_0 = \{x \mid e^\top x + f > 0\}$$

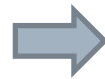
- A quasiconvex optimization problem



Linear-fractional Programming

□ Transforming to a linear program

$$\begin{aligned} \min \quad & f_0(x) = \frac{c^\top x + d}{e^\top x + f} \\ \text{s. t.} \quad & Gx \leq h \\ & Ax = b \end{aligned}$$



$$\begin{aligned} \min \quad & c^\top y + dz \\ \text{s. t.} \quad & Gy - hz \leq 0 \\ & Ay - bz = 0 \\ & e^\top y + fz = 1 \\ & z \geq 0 \end{aligned}$$

■ Proof

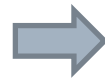
x is feasible in LFP $\Rightarrow y = \frac{x}{e^\top x + f}$, $z = \frac{1}{e^\top x + f}$ is feasible in LP, $c^\top y + dz = f_0(x) \Rightarrow$ the optimal value of LFP is greater than or equal to the optimal value of LP



Linear-fractional Programming

□ Transforming to a linear program

$$\begin{aligned} \min \quad & f_0(x) = \frac{c^\top x + d}{e^\top x + f} \\ \text{s. t.} \quad & Gx \leq h \\ & Ax = b \end{aligned}$$



$$\begin{aligned} \min \quad & c^\top y + dz \\ \text{s. t.} \quad & Gy - hz \leq 0 \\ & Ay - bz = 0 \\ & e^\top y + fz = 1 \\ & z \geq 0 \end{aligned}$$

■ Proof

(y, z) is feasible in LP and $z \neq 0 \Rightarrow x = y/z$ is feasible in LFP, $f_0(x) = c^\top y + dz \Rightarrow$ the optimal value of LFP is less than or equal to the optimal value of LP

(y, z) is feasible in LP, $z = 0$ and x_0 is feasible in LFP $\Rightarrow x = x_0 + ty$ is feasible in LFP for all $t \geq 0$,

$$\lim_{t \rightarrow \infty} f_0(x_0 + ty) = c^\top y + dz$$

Generalized Linear-fractional Programming



□ Generalized Linear-fractional Program

$$f_0(x) = \max_{i=1, \dots, r} \frac{c_i^\top x + d_i}{e_i^\top x + f_i}$$

- $\text{dom } f_0 = \{x \mid e_i^\top x + f_i > 0, i = 1, \dots, r\}$
- A quasiconvex optimization problem

□ Von Neumann Growth Problem

$$\begin{aligned} \max \quad & \min_{i=1, \dots, n} x_i^+ / x_i \\ \text{s. t.} \quad & x^+ \succeq 0 \\ & Bx^+ \preceq Ax \end{aligned}$$

Generalized Linear-fractional Programming



□ Von Neumann Growth Problem

$$\begin{aligned} \max \quad & \min_{i=1, \dots, n} x_i^+ / x_i \\ \text{s.t.} \quad & x^+ \geq 0 \\ & Bx^+ \leq Ax \end{aligned}$$

- $x, x^+ \in \mathbf{R}^n$: activity levels of n sectors, in current and next period
- $(Ax)_i, (Bx^+)_i$: produced and consumed amounts of good i
- $Bx^+ \leq Ax$: goods consumed in the next period cannot exceed the goods produced in the current period
- x_i^+ / x_i growth rate of sector i



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Quadratic Optimization Problems



□ Quadratic Program (QP)

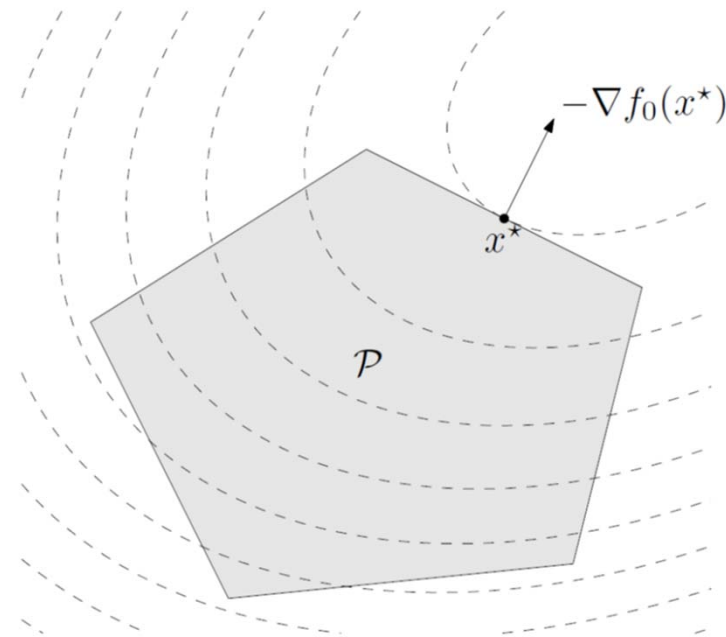
$$\begin{aligned} \min \quad & (1/2)x^\top Px + q^\top x + r \\ \text{s. t.} \quad & Gx \preceq h \\ & Ax = b \end{aligned}$$

- $P \in \mathbf{S}_+^n, G \in \mathbf{R}^{m \times n}$ and $A \in \mathbf{R}^{p \times n}$
- The objective function is (convex) quadratic
- The constraint functions are affine
- When $P = 0$, QP becomes LP

Quadratic Optimization Problems



□ Geometric Illustration of QP



- The feasible set \mathcal{P} is a polyhedron
- The contour lines of the objective function are shown as dashed curves.

Quadratic Optimization Problems



□ Quadratically Constrained Quadratic Program (QCQP)

$$\begin{aligned} \min \quad & (1/2)x^\top P_0 x + q_0^\top x + r_0 \\ \text{s.t.} \quad & (1/2)x^\top P_i x + q_i^\top x + r_i \leq 0, \quad i = 1, \dots, m \\ & Ax = b \end{aligned}$$

- $P_i \in \mathbf{S}_+^n, i = 0, \dots, m$
- The inequality constraint functions are (convex) quadratic
- The feasible set is the intersection of ellipsoids (when $P_i \succ 0$) and an affine set
- Include QP as a special case



Examples

□ Least-squares and Regression

$$\min \|Ax - b\|_2^2 = x^T A^T A x - 2b^T A x + b^T b$$

- Analytical solution: $x = A^\dagger b$
- Can add linear constraints, e.g., $l \preceq x \preceq u$

□ Distance Between Polyhedra

$$\begin{aligned} \min \quad & \|x_1 - x_2\|_2^2 \\ \text{s. t.} \quad & A_1 x_1 \preceq b_1, \quad A_2 x_2 \preceq b_2 \end{aligned}$$

- Find the distance between the polyhedra $\mathcal{P}_1 = \{x | A_1 x \preceq b_1\}$ and $\mathcal{P}_2 = \{x | A_2 x \preceq b_2\}$

$$\text{dist}(\mathcal{P}_1, \mathcal{P}_2) = \inf\{\|x_1 - x_2\|_2 | x_1 \in \mathcal{P}_1, x_2 \in \mathcal{P}_2\}$$



Example

□ Bounding Variance

- x is a random variable on $\{u_1, \dots, u_n\} \subseteq \mathbf{R}$
- $p_i = \mathbf{prob}(x = u_i)$, $p \geq 0$, $\mathbf{1}^\top p = 1$
- The variance of a random variable $f(x)$

$$\mathbf{E}f^2 - (\mathbf{E}f)^2 = \sum_{i=1}^n f_i^2 p_i - \left(\sum_{i=1}^n f_i p_i \right)^2$$

- Maximize the variance

$$\begin{aligned} \max \quad & \sum_{i=1}^n f_i^2 p_i - \left(\sum_{i=1}^n f_i p_i \right)^2 \\ \text{s. t.} \quad & p \geq 0, \mathbf{1}^\top p = 1 \\ & \alpha_i \leq a_i^\top p \leq \beta_i, i = 1, \dots, m \end{aligned}$$

Second-order Cone Programming



□ Second-order Cone Program (SOCP)

$$\begin{aligned} \min \quad & f^\top x \\ \text{s. t.} \quad & \|A_i x + b_i\|_2 \leq c_i^\top x + d_i, \quad i = 1, \dots, m \\ & Fx = g \end{aligned}$$

- $A_i \in \mathbf{R}^{n_i \times n}, F \in \mathbf{R}^{p \times n}$
- Second-order Cone (SOC) constraint:
 $\|Ax + b\|_2 \leq c^\top x + d$ where $A \in \mathbf{R}^{k \times n}$, is same as requiring $(Ax + b, c^\top x + d) \in \text{SOC}$ in \mathbf{R}^{k+1}

$$\begin{aligned} \text{SOC} &= \{(x, t) \in \mathbf{R}^{k+1} \mid \|x\|_2 \leq t\} \\ &= \left\{ \begin{bmatrix} x \\ t \end{bmatrix} \mid \begin{bmatrix} x \\ t \end{bmatrix}^\top \begin{bmatrix} I & 0 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} x \\ t \end{bmatrix} \leq 0, t \geq 0 \right\} \end{aligned}$$

Second-order Cone Programming



□ Second-order Cone Program (SOCP)

$$\begin{aligned} \min \quad & f^\top x \\ \text{s. t.} \quad & \|A_i x + b_i\|_2 \leq c_i^\top x + d_i, \quad i = 1, \dots, m \\ & Fx = g \end{aligned}$$

- $A_i \in \mathbf{R}^{n_i \times n}, F \in \mathbf{R}^{p \times n}$
- Second-order Cone (SOC) constraint:
 $\|Ax + b\|_2 \leq c^\top x + d$ where $A \in \mathbf{R}^{k \times n}$, is same as requiring $(Ax + b, c^\top x + d) \in \text{SOC}$ in \mathbf{R}^{k+1}
- If $c_i = 0, i = 1, \dots, m$, it reduces to QCQP by squaring each inequality constraint
- More general than QCQP and LP



Example

□ Robust Linear Programming

$$\begin{aligned} \min \quad & c^\top x \\ \text{s. t.} \quad & a_i^\top x \leq b_i, \quad i = 1, \dots, m \end{aligned}$$

- There can be uncertainty in a_i
- Assume a_i are known to lie in ellipsoids
 $a_i \in \mathcal{E}_i = \{\bar{a}_i + P_i u \mid \|u\|_2 \leq 1\}, P_i \in R^{n \times n}$
- The constraints must hold for all $a_i \in \mathcal{E}_i$

$$\begin{aligned} \min \quad & c^\top x \\ \text{s. t.} \quad & a_i^\top x \leq b_i \text{ for all } a_i \in \mathcal{E}_i, \quad i = 1, \dots, m \end{aligned}$$

$$\begin{aligned} \min \quad & c^\top x \\ \text{s. t.} \quad & \sup\{a_i^\top x \mid a_i \in \mathcal{E}_i\} \leq b_i, \quad i = 1, \dots, m \end{aligned}$$



Example

- Note that

$$\begin{aligned}\sup\{a_i^\top x \mid a_i \in \mathcal{E}_i\} &= \bar{a}_i^\top x + \sup\{u^\top P_i^\top x \mid \|u\|_2 \leq 1\} \\ &= \bar{a}_i^\top x + \|P_i^\top x\|_2\end{aligned}$$

- Robust linear constraint

$$\bar{a}_i^\top x + \|P_i^\top x\|_2 \leq b_i$$

- SOCP

$$\begin{aligned}\min \quad & c^\top x \\ \text{s. t.} \quad & \bar{a}_i^\top x + \|P_i^\top x\|_2 \leq b_i, \quad i = 1, \dots, m\end{aligned}$$



Example

- Linear Programming with Random Constraints
 - Suppose that a_i is independent Gaussian random vectors with mean \bar{a}_i and covariance Σ_i
 - Require each constraint $a_i^\top x \leq b_i$ holds with probability exceeding $\eta \geq 0.5$

$$\begin{array}{ll} \min & c^\top x \\ \text{s. t.} & \text{prob}(a_i^\top x \leq b_i) \geq \eta, \quad i = 1, \dots, m \end{array}$$



Example

□ Linear Programming with Random Constraints (SOCP)

$$\begin{aligned} \min \quad & c^T x \\ \text{s.t.} \quad & \bar{a}_i^T x + \Phi^{-1}(\eta) \left\| \Sigma_i^{1/2} x \right\|_2 \leq b_i, \quad i = 1, \dots, m \end{aligned}$$

■ Analysis

$$\text{prob}(a_i^T x \leq b_i) = \text{prob}\left(\frac{a_i^T x - \bar{a}_i^T x}{\left\| \Sigma_i^{1/2} x \right\|_2} \leq \frac{b_i - \bar{a}_i^T x}{\left\| \Sigma_i^{1/2} x \right\|_2}\right) \geq \eta \Leftrightarrow$$

$$\Phi\left(\frac{b_i - \bar{a}_i^T x}{\left\| \Sigma_i^{1/2} x \right\|_2}\right) \geq \eta \Leftrightarrow \frac{b_i - \bar{a}_i^T x}{\left\| \Sigma_i^{1/2} x \right\|_2} \geq \Phi^{-1}(\eta) \Leftrightarrow \bar{a}_i^T x +$$

$$\Phi^{-1}(\eta) \left\| \Sigma_i^{1/2} x \right\|_2 \leq b_i$$

$$\text{where } \Phi(z) = (1/\sqrt{2\pi}) \int_{-\infty}^z e^{-t^2/2} dt$$



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Definitions

□ Monomial Function

$$f(x) = cx_1^{a_1} x_2^{a_2} \dots x_n^{a_n}$$

- $f: \mathbf{R}^n \rightarrow \mathbf{R}$, $\text{dom } f = \mathbf{R}_{++}^n$, $c > 0$ and $a_i \in \mathbf{R}$
- Closed under multiplication, division, and nonnegative scaling

□ Posynomial Function

$$f(x) = \sum_{k=1}^K c_k x_1^{a_{1k}} x_2^{a_{2k}} \dots x_n^{a_{nk}}$$

- Closed under addition, multiplication, and nonnegative scaling



Geometric Programming (GP)

□ The Problem

$$\begin{aligned} \min \quad & f_0(x) \\ \text{s. t.} \quad & f_i(x) \leq 1, \quad i = 1, \dots, m \\ & h_i(x) = 1, \quad i = 1, \dots, p \end{aligned}$$

- f_0, \dots, f_m are posynomials
- h_1, \dots, h_p are monomials
- Domain of the problem

$$\mathcal{D} = \mathbf{R}_{++}^n$$

- Implicit constraint: $x \succ 0$



Extensions of GP

- f is a posynomial and h is a monomial

$$f(x) \leq h(x) \Leftrightarrow \frac{f(x)}{h(x)} \leq 1$$

- h_1 and h_2 are nonzero monomials

$$h_1(x) = h_2(x) \Leftrightarrow \frac{h_1(x)}{h_2(x)} = 1$$

- Maximize a nonzero monomial objective function by minimizing its inverse

$$\begin{array}{ll} \max & x/y \\ \text{s.t.} & 2 \leq x \leq 3 \\ & x^2 + 3y/z \leq \sqrt{y} \\ & x/y = z^2 \end{array} \quad \Leftrightarrow$$

$$\begin{array}{ll} \min & x^{-1}y \\ \text{s.t.} & 2x^{-1} \leq 1, (1/3)x \leq 1 \\ & x^2y^{-1/2} + y^{1/2}z^{-1} \leq 1 \\ & xy^{-1}z^{-2} = 1 \end{array}$$



GP in Convex Form

□ Change of Variables $y_i = \log x_i$

- f is the monomial function

$$f(x) = cx_1^{a_1} x_2^{a_2} \dots x_n^{a_n}, \quad x_i = e^{y_i}$$

$$f(x) = f(e^{y_1}, \dots, e^{y_n}) = c(e^{y_1})^{a_1} \dots (e^{y_n})^{a_n}$$

$$= e^{a_1 y_1 + \dots + a_n y_n + \log c} = e^{a^\top y + b}$$

- f is the posynomial function

$$f(x) = \sum_{k=1}^K c_k x_1^{a_{1k}} x_2^{a_{2k}} \dots x_n^{a_{nk}}$$

$$f(x) = \sum_{k=1}^K e^{a_k^\top y + b_k}$$



GP in Convex Form

□ New Form

$$\begin{aligned} \min \quad & \sum_{k=1}^{K_0} e^{a_{0k}^\top y + b_{0k}} \\ \text{s. t.} \quad & \sum_{k=1}^{K_i} e^{a_{ik}^\top y + b_{ik}} \leq 1, \quad i = 1, \dots, m \\ & e^{g_i^\top y + h_i} = 1, \quad i = 1, \dots, p \end{aligned}$$

□ Taking the Logarithm

$$\begin{aligned} \min \quad & \tilde{f}_0(y) = \log \left(\sum_{k=1}^{K_0} e^{a_{0k}^\top y + b_{0k}} \right) \\ \text{s. t.} \quad & \tilde{f}_i(y) = \log \left(\sum_{k=1}^{K_i} e^{a_{ik}^\top y + b_{ik}} \right) \leq 0, \quad i = 1, \dots, m \\ & \tilde{h}_i(y) = g_i^\top y + h_i = 0, \quad i = 1, \dots, p \end{aligned}$$



Example

□ Frobenius Norm Diagonal Scaling

- Given a matrix $M \in \mathbf{R}^{n \times n}$
- Choose a diagonal matrix D such that DMD^{-1} is small

$$\begin{aligned}\|DMD^{-1}\|_F^2 &= \text{tr}((DMD^{-1})^\top (DMD^{-1})) = \sum_{i,j=1}^n (DMD^{-1})_{ij}^2 \\ &= \sum_{i,j=1}^n M_{ij}^2 d_i^2 / d_j^2\end{aligned}$$

- Unconstrained GP

$$\min \sum_{i,j=1}^n M_{ij}^2 d_i^2 / d_j^2$$



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Generalized Inequality Constraints



□ Convex Optimization Problem with Generalized Inequality Constraints

$$\begin{aligned} \min \quad & f_0(x) \\ \text{s. t.} \quad & f_i(x) \preceq_{K_i} 0, \quad i = 1, \dots, m \\ & Ax = b \end{aligned}$$

- $f_0: \mathbf{R}^n \rightarrow \mathbf{R}$ is convex;
- $K_i \subseteq \mathbf{R}^{k_i}$ are proper cones
- $f_i: \mathbf{R}^n \rightarrow \mathbf{R}^{k_i}$ is K_i -convex w.r.t. proper cone $K_i \subseteq \mathbf{R}^{k_i}$

Generalized Inequality Constraints



□ Convex Optimization Problem with Generalized Inequality Constraints

$$\begin{aligned} \min \quad & f_0(x) \\ \text{s. t.} \quad & f_i(x) \preceq_{K_i} 0, \quad i = 1, \dots, m \\ & Ax = b \end{aligned}$$

- The feasible set, any sublevel set, and the optimal set are convex
- Any locally optimal is globally optimal
- The optimality condition for differentiable f_0 holds without change



Conic Form Problems

□ Conic Form Problems

$$\begin{array}{ll} \min & c^\top x \\ \text{s. t.} & Fx + g \preceq_K 0 \\ & Ax = b \end{array}$$

- A linear objective
- One inequality constraint function which is **affine**
- A generalization of linear programs



Conic Form Problems

□ Conic Form Problems

$$\begin{array}{ll} \min & c^\top x \\ \text{s. t.} & Fx + g \preceq_K 0 \\ & Ax = b \end{array}$$

□ Standard Form

$$\begin{array}{ll} \min & c^\top x \\ \text{s. t.} & x \succeq_K 0 \\ & Ax = b \end{array}$$

□ Inequality Form

$$\begin{array}{ll} \min & c^\top x \\ \text{s. t.} & Fx + g \preceq_K 0 \end{array}$$



Semidefinite Programming

□ Semidefinite Program (SDP)

$$\begin{aligned} \min \quad & c^\top x \\ \text{s. t.} \quad & x_1 F_1 + \cdots + x_n F_n + G \preceq 0 \\ & Ax = b \end{aligned}$$

- $K = \mathbf{S}_+^k$
- $G, F_1, \dots, F_n \in \mathbf{S}^k$ and $A \in \mathbf{R}^{p \times n}$
- Linear matrix inequality (LMI)
- If G, F_1, \dots, F_n are all diagonal, LMI is equivalent to a set of n linear inequalities, and SDP reduces to LP



Semidefinite Programming

□ Standard Form SDP

$$\begin{aligned} \min \quad & \text{tr}(CX) \\ \text{s. t.} \quad & \text{tr}(A_i X) = b_i, \quad i = 1, \dots, p \\ & X \succeq 0 \end{aligned}$$

- $X \in \mathbf{S}^n$ is the variable and $C, A_1, \dots, A_p \in \mathbf{S}^n$
- p linear equality constraints
- A nonnegativity constraint

□ Inequality Form SDP

$$\begin{aligned} \min \quad & c^T x \\ \text{s. t.} \quad & x_1 A_1 + \dots + x_n A_n \preceq B \end{aligned}$$

- $B, A_1, \dots, A_p \in \mathbf{S}^k$ and no equality constraint



Semidefinite Programming

□ Multiple LMIs and Linear Inequalities

$$\begin{aligned} \min \quad & c^\top x \\ \text{s. t.} \quad & F^{(i)}(x) = x_1 F_1^{(i)} + \cdots + x_n F_n^{(i)} + G^{(i)} \preceq 0, i = 1, \dots, K \\ & Gx \preceq h, \quad Ax = b \end{aligned}$$

- It is referred as SDP as well

□ Be transformed as

$$\begin{aligned} \min \quad & c^\top x \\ \text{s. t.} \quad & \text{diag}(Gx - h, F^{(1)}(x), \dots, F^{(K)}(x)) \preceq 0 \\ & Ax = b \end{aligned}$$

- A standard SDP



Examples

□ Second-order Cone Programming

$$\begin{aligned} \min \quad & c^\top x \\ \text{s. t.} \quad & \|A_i x + b_i\|_2 \leq c_i^\top x + d_i, \quad i = 1, \dots, m \\ & Fx = g \end{aligned}$$

■ A conic form problem

$$\begin{aligned} \min \quad & c^\top x \\ \text{s. t.} \quad & -(A_i x + b_i, c_i^\top x + d_i) \preceq_{K_i} 0, \quad i = 1, \dots, m \\ & Fx = g \end{aligned}$$

in which

$$K_i = \{(y, t) \in \mathbf{R}^{n_i+1} \mid \|y\|_2 \leq t\}$$



Example

□ Matrix Norm Minimization

$$\min \|A(x)\|_2 = (\lambda_{\max}(A(x)^T A(x)))^{1/2}$$

- $A(x) = A_0 + x_1 A_1 + \cdots + x_n A_n$ and $A_i \in \mathbf{R}^{p \times q}$

- Fact: $\|A\|_2 \leq t \Leftrightarrow A^T A \preceq t^2 I$

□ A New Problem

$$\min \|A(x)\|_2^2 \Leftrightarrow \begin{array}{ll} \min & s \\ \text{s. t.} & \|A(x)\|_2^2 \preceq s \end{array}$$



Example

□ Matrix Norm Minimization

$$\min \|A(x)\|_2 = (\lambda_{\max}(A(x)^T A(x)))^{1/2}$$

- $A(x) = A_0 + x_1 A_1 + \cdots + x_n A_n$ and $A_i \in \mathbf{R}^{p \times q}$

- Fact: $\|A\|_2 \leq t \Leftrightarrow A^T A \preceq t^2 I$

□ A New Problem

$$\begin{array}{ll} \min & s \\ \text{s.t.} & A(x)^T A(x) \preceq sI \end{array} \Leftrightarrow \begin{array}{ll} \min & s \\ \text{s.t.} & A(x)^T A(x) - sI \preceq 0 \end{array}$$

- $A(x)^T A(x) - sI$ is matrix convex



Example

□ Matrix Norm Minimization

$$\min \|A(x)\|_2 = (\lambda_{\max}(A(x)^\top A(x)))^{1/2}$$

- $A(x) = A_0 + x_1 A_1 + \cdots + x_n A_n$ and $A_i \in \mathbf{R}^{p \times q}$

- Fact:

$$\|A\|_2 \leq t \Leftrightarrow A^\top A \preceq t^2 I \Leftrightarrow \begin{bmatrix} tI & A \\ A^\top & tI \end{bmatrix} \succeq 0$$

□ SDP

$$\begin{array}{ll} \min & t \\ \text{s. t.} & \begin{bmatrix} tI & A(x) \\ A(x)^\top & tI \end{bmatrix} \succeq 0 \end{array}$$

- A single linear matrix inequality



Outline

- Linear Optimization Problems
- Quadratic Optimization Problems
- Geometric Programming
- Generalized Inequality Constraints
- Vector Optimization

General and Convex Vector Optimization Problems



□ General Vector Optimization Problem

$$\begin{array}{ll} \min \text{ (w. r. t. } K) & f_0(x) \\ \text{s. t.} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p \end{array}$$

- $f_0: \mathbf{R}^n \rightarrow \mathbf{R}^q$ is a **vector-valued** objective function
- $K \in \mathbf{R}^q$ is a proper cone, which is used to compare objective values
- $f_i: \mathbf{R}^n \rightarrow \mathbf{R}$ are the inequality constraint functions
- $h_i: \mathbf{R}^n \rightarrow \mathbf{R}$ are the equality constraint functions

General and Convex Vector Optimization Problems



□ Convex Vector Optimization Problem

$$\begin{array}{ll} \min \text{ (w.r.t. } K) & f_0(x) \\ \text{s.t.} & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p \end{array}$$

- $f_0: \mathbf{R}^n \rightarrow \mathbf{R}^q$ is K -convex
- $f_i: \mathbf{R}^n \rightarrow \mathbf{R}$ are convex
- $h_i: \mathbf{R}^n \rightarrow \mathbf{R}$ are affine

□ x is better than or equal to y

$$f_0(x) \preceq_K f_0(y)$$

- Could be incomparable



Optimal Points and Values

□ Achievable Objective Values

$$\mathcal{O} = \{f_0(x) | \exists x \in \mathcal{D}, f_i(x) \leq 0, i = 1, \dots, m, h_i(x) = 0, i = 1, \dots, p\}$$

□ If \mathcal{O} has a minimum element $f_0(x)$

■ x is optimal and $f_0(x)$ is the optimal value

□ x^* is optimal if and only if it is feasible and

$$\mathcal{O} \subseteq f_0(x^*) + K$$



Optimal Points and Values

□ Achievable Objective Values

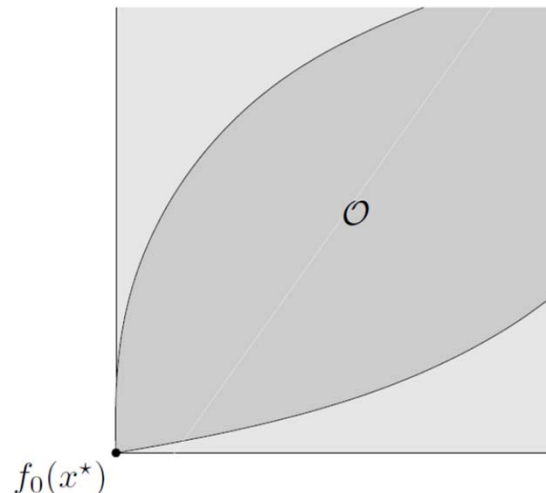
$$\mathcal{O} = \{f_0(x) | \exists x \in \mathcal{D}, f_i(x) \leq 0, i = 1, \dots, m, h_i(x) = 0, i = 1, \dots, p\}$$

□ If \mathcal{O} has a minimum element $f_0(x)$

■ x is optimal and $f_0(x)$ is the optimal value

□ $K = \mathbf{R}_+^2$

$$\mathcal{O} \subseteq f_0(x^*) + K$$





Example

□ Best Linear Unbiased Estimator

- Suppose that $y = Ax + v$, where $v \in \mathbf{R}^m$ is noise, $y \in \mathbf{R}^m$ and $x \in \mathbf{R}^n$
- Estimate x from A and y
- Assume that A has rank n , and $\mathbf{E}v = 0$, $\mathbf{E}vv^T = I$
- A linear estimator $\hat{x} = Fy$
- If $FA = I$, $\hat{x} = Fy$ is an **unbiased** linear estimator of x

$$\mathbf{E}\hat{x} = \mathbf{E}[FAx + Fv] = x$$



Example

□ Best Linear Unbiased Estimator

- The error covariance of an unbiased estimator

$$\mathbf{E}(\hat{x} - x)(\hat{x} - x)^\top = \mathbf{E}Fvv^\top F^\top = FF^\top$$

□ Minimize the covariance

$$\begin{array}{ll} \min (\text{w. r. t. } \mathbf{S}_+^n) & FF^\top \\ \text{s. t.} & FA = I \end{array}$$

- Solution

$$F^* = A^\dagger = (A^\top A)^{-1} A^\top$$

$$F^* F^{*\top} = (A^\top A)^{-1}$$

Pareto Optimal Points and Values



□ Achievable Objective Values

$$\mathcal{O} = \{f_0(x) \mid \exists x \in \mathcal{D}, f_i(x) \leq 0, i = 1, \dots, m, h_i(x) = 0, i = 1, \dots, p\}$$

□ $f_0(x)$ is a minimal element of \mathcal{O}

■ x is Pareto optimal

■ $f_0(x)$ is a Pareto optimal value

□ x is Pareto optimal if and only if it is feasible and

$$(f_0(x) - K) \cap \mathcal{O} = \{f_0(x)\}$$

Pareto Optimal Points and Values



□ Achievable Objective Values

$$\mathcal{O} = \{f_0(x) \mid \exists x \in \mathcal{D}, f_i(x) \leq 0, i = 1, \dots, m, h_i(x) = 0, i = 1, \dots, p\}$$

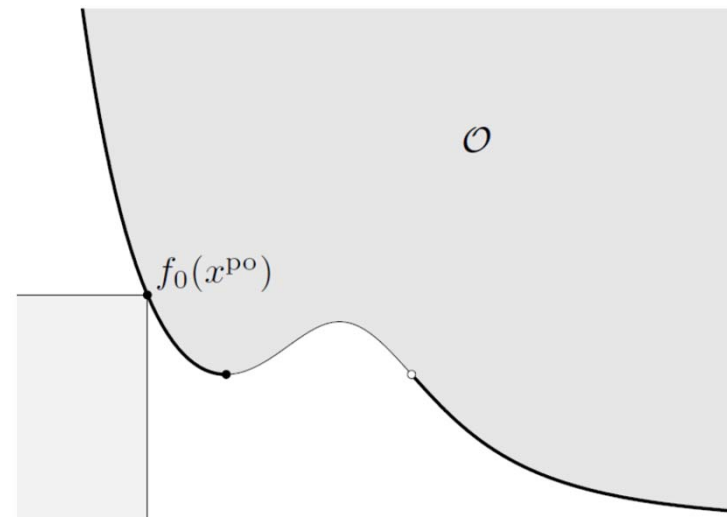
□ $f_0(x)$ is a minimal element of \mathcal{O}

■ x is Pareto optimal

■ $f_0(x)$ is a Pareto optimal value

□ $K = \mathbf{R}_+^2$

$$(f_0(x) - K) \cap \mathcal{O} = \{f_0(x)\}$$



Pareto Optimal Points and Values



□ Achievable Objective Values

$$\mathcal{O} = \{f_0(x) \mid \exists x \in \mathcal{D}, f_i(x) \leq 0, i = 1, \dots, m, h_i(x) = 0, i = 1, \dots, p\}$$

□ $f_0(x)$ is a minimal element of \mathcal{O}

■ x is Pareto optimal

■ $f_0(x)$ is a Pareto optimal value

□ x is Pareto optimal if and only if it is feasible and

$$(f_0(x) - K) \cap \mathcal{O} = \{f_0(x)\}$$

□ Let \mathcal{P} be the set of Pareto optimal values

$$\mathcal{P} \subseteq \mathcal{O} \cap \text{bd}\mathcal{O}$$



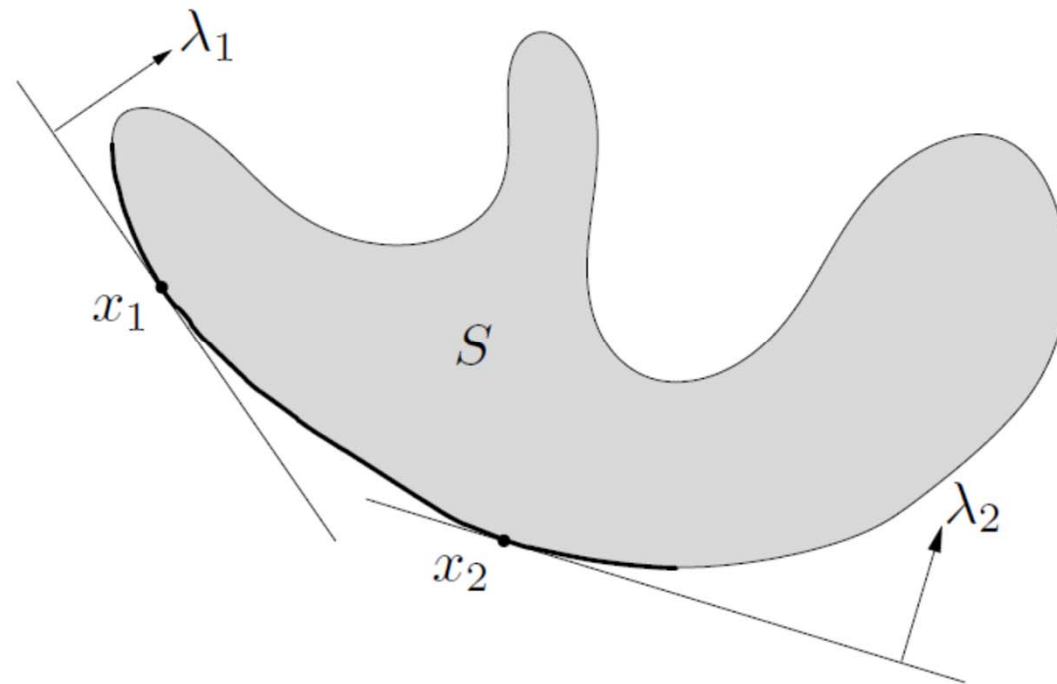
Scalarization

- A standard technique for finding Pareto optimal (or optimal) points
- Find Pareto optimal points for any vector optimization problem by solving the ordinary **scalar** optimization problem
- Characterization of minimum and minimal points via dual generalized inequalities

Dual Characterization of Minimal Elements (1)



- If $\lambda \succ_{K^*} 0$, and x minimizes $\lambda^\top z$ over $z \in S$, then x is minimal.





Scalarization

□ Choose any $\lambda \succ_{K^*} 0$

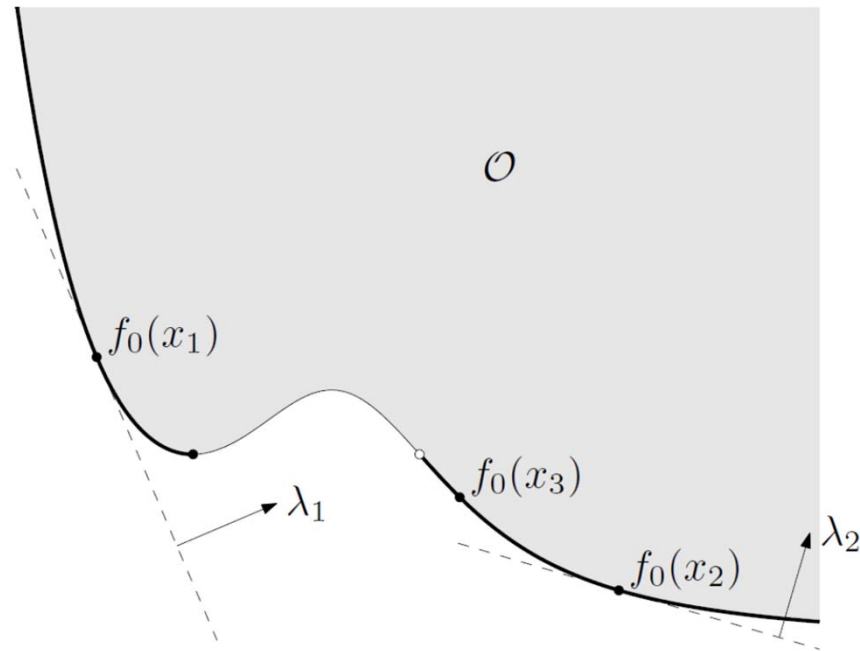
$$\begin{aligned} \min \quad & \lambda^\top f_0(x) \\ \text{s. t.} \quad & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p \end{aligned}$$

- The optimal point x for this scalar problem is Pareto optimal for the vector optimization problem
- λ is called the weight vector
- By varying λ we obtain (possibly) different Pareto optimal solutions



Scalarization

□ $K = \mathbf{R}_+^2$



- Scalarization **cannot** find $f_0(x_3)$

Scalarization of Convex Vector Optimization Problems



□ Choose any $\lambda \succ_{K^*} 0$

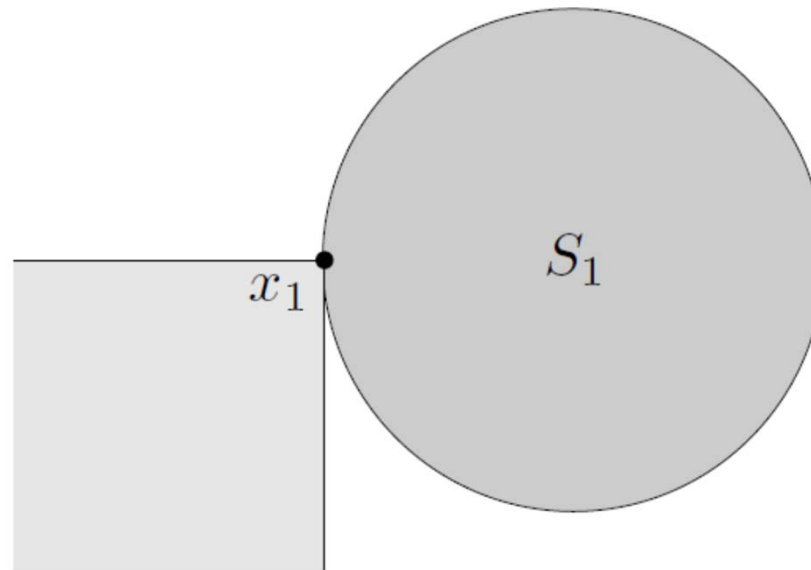
$$\begin{aligned} \min \quad & \lambda^\top f_0(x) \\ \text{s. t.} \quad & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p \end{aligned}$$

- A **convex** optimization problem
- The optimal point x for this scalar problem is Pareto optimal for the vector optimization problem
- λ is called the weight vector
- By varying λ we obtain (possibly) different Pareto optimal solutions

Dual Characterization of Minimal Elements (2)



- If S is convex, for any minimal element x there exists a nonzero $\lambda \succcurlyeq_{K^*} 0$ such that x minimizes $\lambda^T z$ over $z \in S$.



x_1 minimizes $\lambda^T z$ over $z \in S_1$ for $\lambda = (1, 0) \succcurlyeq 0$

Scalarization of Convex Vector Optimization Problems



- For every Pareto optimal point x^{p0} , there is some nonzero $\lambda \succ_{K^*} 0$ such that x^{p0} is a solution of the scalarized problem

$$\begin{aligned} \min \quad & \lambda^\top f_0(x) \\ \text{s. t.} \quad & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p \end{aligned}$$

- It is **not true** that every solution of the scalarized problem, with $\lambda \succ_{K^*} 0$ and $\lambda \neq 0$, is a Pareto optimal point for the vector problem

Scalarization of Convex Vector Optimization Problems



1. Consider all $\lambda \succ_{K^*} 0$

$$\begin{aligned} \min \quad & \lambda^\top f_0(x) \\ \text{s. t.} \quad & f_i(x) \leq 0, \quad i = 1, \dots, m \\ & h_i(x) = 0, \quad i = 1, \dots, p \end{aligned}$$

■ Solve the above problem

2. Consider all $\lambda \succeq_{K^*} 0, \lambda \neq 0, \lambda \not\prec_{K^*} 0$

■ Solve the above problem

■ Verify the solution



Example

□ Minimal Upper Bound on a Set of Matrices

$$\begin{array}{ll} \min (\text{w. r. t. } \mathbf{S}_+^n) & X \\ \text{s. t.} & X \succeq A_i, \quad i = 1, \dots, m \end{array}$$

- $A_i \in \mathbf{S}^n, i = 1, \dots, m$
- The constraints mean that X is an upper bound on A_1, \dots, A_m
- A Pareto optimal solution is a minimal upper bound on the matrices



Example

□ Scalarization

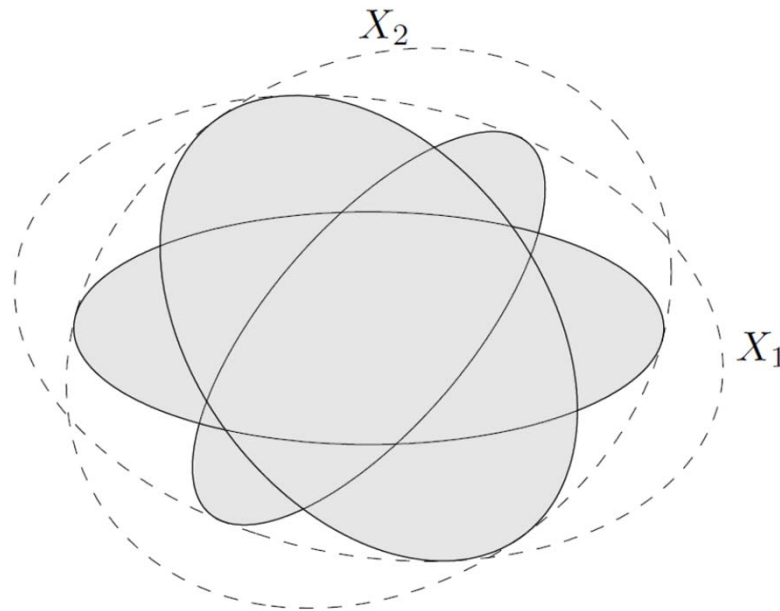
$$\begin{array}{ll} \min & \text{tr}(WX) \\ \text{s. t.} & X \succcurlyeq A_i, \quad i = 1, \dots, m \end{array}$$

- $W \in \mathbf{S}_{++}^n$
- An SDP
- If X is Pareto optimal for the vector problem then it is optimal for the SDP, for some nonzero weight matrix $W \succcurlyeq 0$.



Example

- A Simple Geometric Interpretation
 - Define an ellipsoid centered at the origin as $\mathcal{E}_A = \{u | u^\top A^{-1} u \leq 1\}$
 - $A \preceq B \iff \mathcal{E}_A \subseteq \mathcal{E}_B$





Multicriterion Optimization

□ $K = \mathbf{R}_+^q$

$$f_0(x) = (F_1(x), \dots, F_q(x))$$

- f_0 consists of q different objectives F_i and we want to minimize all F_i
- It is convex if f_1, \dots, f_m are convex, h_1, \dots, h_p are affine, and F_1, \dots, F_q are convex
- Feasible x^* is optimal if
$$y \text{ is feasible} \Rightarrow f_0(x^*) \preceq f_0(y)$$
- Feasible x^{p0} is Pareto optimal if
$$y \text{ is feasible, } f_0(y) \preceq f_0(x^{p0}) \Rightarrow f_0(x^{p0}) = f_0(y)$$



Example

□ Regularized Least-Squares

$$\min (\text{w. r. t. } \mathbf{R}_+^2) \quad f_0(x) = (F_1(x), F_2(x))$$

- $F_1(x) = \|Ax - b\|_2^2$ measures the misfit
- $F_2(x) = \|x\|_2^2$ measures the size
- Our goal is to find x that gives a good fit and that is not large

□ Scalarization

$$\begin{aligned} \lambda^\top f_0(x) &= \lambda_1 F_1(x) + \lambda_2 F_2(x) \\ &= x^\top (\lambda_1 A^\top A + \lambda_2 I)x - 2\lambda_1 b^\top Ax + \lambda_1 b^\top b \end{aligned}$$

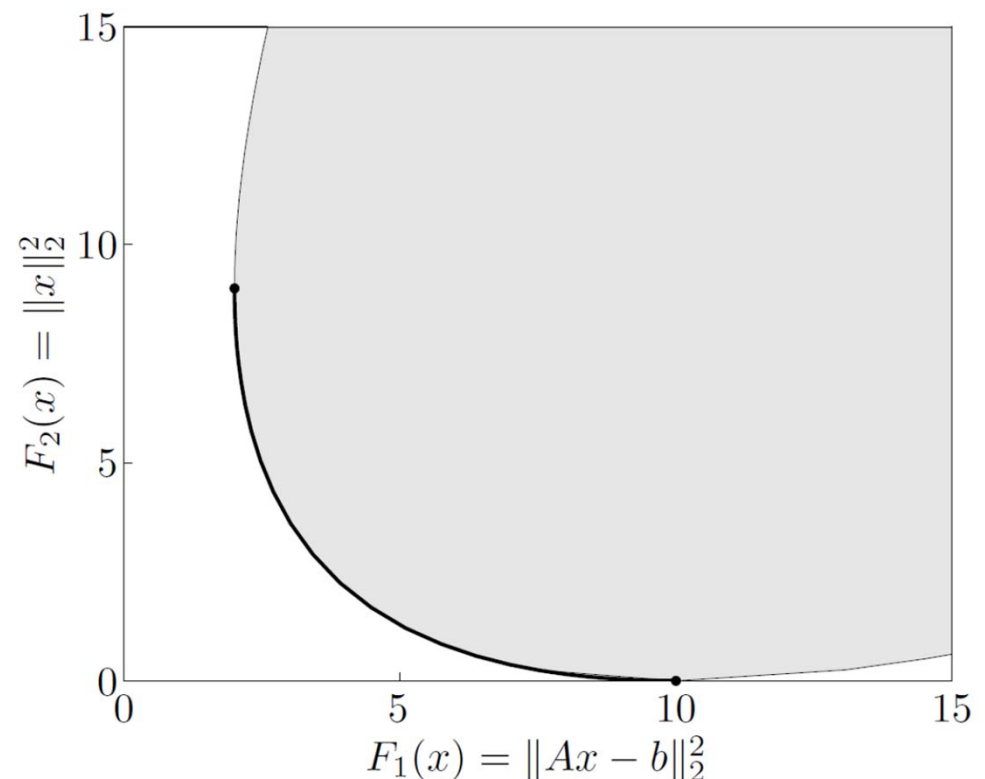


Example

□ Solution

$$x(\mu) = (\lambda_1 A^\top A + \lambda_2 I)^{-1} \lambda_1 A^\top b = (A^\top A + \mu I)^{-1} A^\top b$$

- $\mu = \lambda_2 / \lambda_1$
- $\lambda = (0, 1)$, we get $x = 0$
- With $\lambda \rightarrow (1, 0)$, we get $x = A^\dagger b$





Summary

- Linear Optimization Problems
- Quadratic Optimization Problems
- Geometric Programming
- Generalized Inequality Constraints
- Vector Optimization



Homework 2

- <http://www.lamda.nju.edu.cn/wanyy/optfall2020gra.html>
- Deadline: Dec 8