# Introduction

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### Outline

Mathematical Optimization

- Least-squares
- Linear Programming
- Convex Optimization
- Nonlinear Optimization
- □ Summary



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# Mathematical Optimization (1)

Optimization Problem min  $f_0(x)$ 

- s.t.  $f_i(x) \le b_i$ , i = 1, ..., m
- Optimization Variable:  $x = (x_1, ..., x_n)$
- Objective Function:  $f_0: \mathbf{R}^n \to \mathbf{R}$
- Constraint Functions:  $f_i: \mathbf{R}^n \to \mathbf{R}$
- $\Box x^*$  is called optimal or a solution
  - $f_i(x^*) \le b_i, \ i = 1, \dots, m$
  - For any z with  $f_i(z) \le b_i$ , we have  $f_0(z) \ge f_0(x^*)$



# Mathematical Optimization (2)

Linear Program

$$f_i(\alpha x + \beta y) = \alpha f_i(x) + \beta f_i(y)$$

for all  $x, y \in \mathbf{R}^n$  and all  $\alpha, \beta \in \mathbf{R}$ 

#### Nonlinear Program

- If the optimization problem is not linear
- Convex Optimization Problem

 $f_i(\alpha x + \beta y) \le \alpha f_i(x) + \beta f_i(y)$ 

for all  $x, y \in \mathbf{R}^n$  and all  $\alpha, \beta \in \mathbf{R}$  with  $\alpha + \beta = 1, \ \alpha \ge 0, \ \beta \ge 0$ 



### Applications

min 
$$f_0(x)$$
  
s.t.  $f_i(x) \le b_i$ ,  $i = 1, ..., m$ 

#### Abstraction

- x represents the choice made
- $f_i(x) \le b_i$  represent firm requirements that limit the possible choices
- $f_0(x)$  represents the cost of choosing x
- A solution corresponds to a choice that has minimum cost, among all choices that meet the requirements



# Portfolio Optimization

#### Variables

- $x_i$  represents the investment in the *i*-th asset
- $x \in \mathbf{R}^n$  describes the overall portfolio allocation across the set of asset
- Constraints
  - A limit on the budget the requirement
  - Investments are nonnegative
  - A minimum acceptable value of expected return for the whole portfolio
- □ Objective
  - Minimize the variance of the portfolio return



### **Device Sizing**

#### Variables

•  $x \in \mathbf{R}^n$  describes the widths and lengths of the devices

#### Constraints

- Limits on the device sizes
- Timing requirements
- A limit on the total area of the circuit

#### Objective

Minimize the total power consumed by the circuit



# Data Fitting

#### Variables

•  $x \in \mathbf{R}^n$  describes parameters in the model

#### Constraints

- Prior information
- Required limits on the parameters (such as nonnegativity)

#### Objective

Minimize the prediction error between the observed data and the values predicted by the model



# Solving Optimization Problem

#### General Optimization Problem

- Very difficult to solve
- Constraints can be very complicated, the number of variables can be very large
- Methods involve some compromise, e.g., computation time, or suboptimal solution

#### Exceptions

- Least-squares problems
- Linear programming problems
- Convex optimization problems



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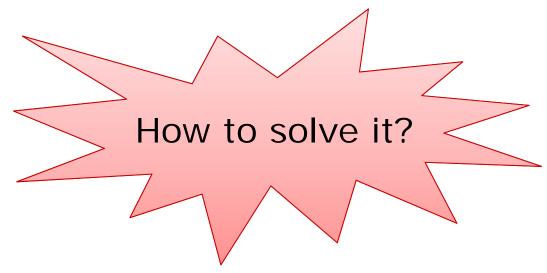
### Least-squares Problems (1)

□ The Problem

min 
$$||Ax - b||_2^2 = \sum_{i=1}^k (a_i^T x - b_i)^2$$

•  $A \in \mathbf{R}^{k \times n}$ ,  $a_i^{\top}$  is the *i*-th row of A,  $b \in \mathbf{R}^k$ 

•  $x \in \mathbf{R}^n$  is the optimization variable





### Least-squares Problems (1)

□ The Problem

min 
$$||Ax - b||_2^2 = \sum_{i=1}^k (a_i^T x - b_i)^2$$

A ∈ R<sup>k×n</sup>, a<sub>i</sub><sup>T</sup> is the *i*-th row of A, b ∈ R<sup>k</sup>
 x ∈ R<sup>n</sup> is the optimization variable
 Setting the gradient to be 0

$$2A^{\top}(Ax - b) = 0$$
  

$$\Rightarrow A^{\top}Ax = A^{\top}b$$
  

$$\Rightarrow x = (A^{\top}A)^{-1}A^{\top}b$$



### Least-squares Problems (2)

- $\Box A \text{ Set of Linear Equations} \\ A^{\mathsf{T}}Ax = A^{\mathsf{T}}b$
- □ Solving least-squares problems
  - Reliable and efficient algorithms and software
  - Computation time proportional to  $n^2k \ (A \in \mathbf{R}^{k \times n})$ ; less if structured
  - A mature technology
  - Challenging for extremely large problems



# Using Least-squares

Easy to RecognizeWeighted least-squares

$$\sum_{i=1}^{\kappa} w_i (a_i^{\mathsf{T}} x - b_i)^2$$

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Different importance



### Using Least-squares

Easy to Recognize Weighted least-squares  $\sum_{i=1}^{k} w_i (a_i^{\mathsf{T}} x - b_i)^2 = \sum_{i=1}^{k} (\sqrt{w_i} a_i^{\mathsf{T}} x - \sqrt{w_i} b_i)^2$ Different importance Regularization  $\sum_{i=1}^{n} (a_i^{\mathsf{T}} x - b_i)^2 + \rho \sum_{i=1}^{n} x_i^2$ More stable



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# Linear Programming

 $\square \text{ The Problem} \\ \min c^{\top} x \\ \text{s.t.} a_i^{\top} x \leq b_i, \qquad i = 1, \dots, m$ 

 $c, a_1, \dots, a_m \in \mathbf{R}^n, b_1, \dots, b_m \in \mathbf{R}$ 

#### □ Solving Linear Programs

- No analytical formula for solution
- Reliable and efficient algorithms and software
- Computation time proportional to  $n^2m$  if  $m \ge n$ ; less with structure
- A mature technology
- Challenging for extremely large problems



## Using Linear Programming

Not as easy to recognize
 Chebyshev Approximation Problem

min	$\max_{i=1,\dots,k}  a_i^{T}x - b_i $
min s.t.	$t = \max_{i=1,\dots,k}  a_i^{T} x - b_i $

$$\iff \begin{array}{l} \min \quad t \\ \text{s.t.} \quad t \ge \left| a_i^{\mathsf{T}} x - b_i \right|, i = 1, \dots, k \end{array}$$

 $\iff \begin{array}{l} \min \quad t \\ \text{s.t.} \quad -t \leq a_i^{\mathsf{T}} x - b_i \leq t, i = 1, \dots, k \end{array}$ 



### Outline

Mathematical Optimization

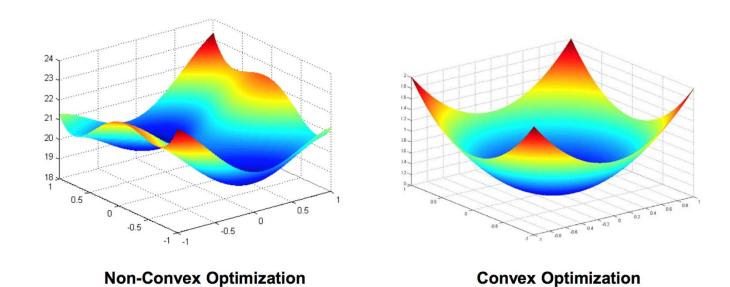
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## **Convex Optimization**

#### □ Why Convexity?

" The great watershed in optimization isn't between linearity and nonlinearity, but convexity and nonconvexity." — R. Rockafellar, SIAM Review 1993





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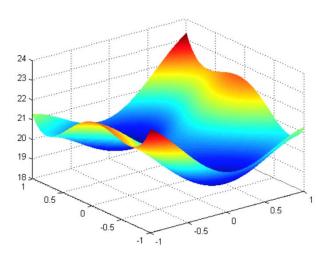
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# **Convex Optimization**

#### □ Why Convexity?

" The great watershed in optimization onlinearity, but convexity and no -R.

Local minimizers are also global minimizers.



**Non-Convex Optimization** 

**Convex Optimization** 



#### **Convex Optimization Problems (1)**

□ The Problem

 $\begin{array}{ll} \min & f_0(x) \\ \text{s.t.} & f_i(x) \le b_i, \qquad i = 1, \dots, m \end{array}$ 

Functions  $f_0, ..., f_m: \mathbb{R}^n \to \mathbb{R}$  are convex:

 $f_i(\alpha x + \beta y) \le \alpha f_i(x) + \beta f_i(y)$ 

for all  $x, y \in \mathbf{R}^n$  and all  $\alpha, \beta \in \mathbf{R}$  with  $\alpha + \beta = 1, \ \alpha \ge 0, \ \beta \ge 0$ 

Least-squares and linear programs as special cases



### **Convex Optimization Problems (2)**

#### Solving Convex Optimization Problems

- No analytical solution
- Reliable and efficient algorithms (e.g., interior-point methods)
- Computation time (roughly) proportional to max{n<sup>3</sup>, n<sup>2</sup>m, F}
  - ✓ F is cost of evaluating the first and second derivatives of  $f'_i$ s
- Almost a technology



Using Convex Optimization

Often difficult to recognize

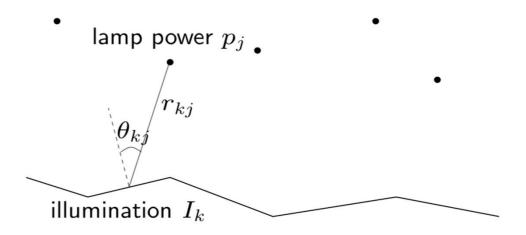
Many tricks for transforming problems into convex form

Surprisingly many problems can be solved via convex optimization



### An Example (1)

#### □ *m* lamps illuminating *n* patches



Intensity I<sub>k</sub> at patch k depends linearly on lamp powers p<sub>i</sub>

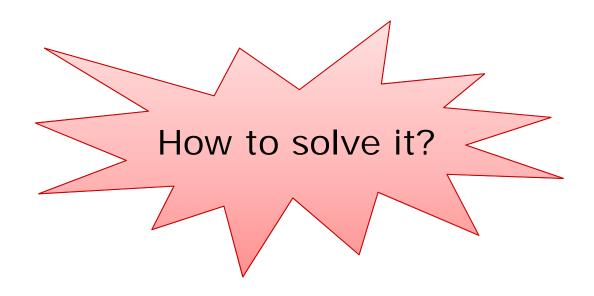
$$I_{k} = \sum_{j=1}^{m} a_{kj} p_{j}, \qquad a_{kj} = r_{kj}^{-2} \max\{\cos\theta_{kj}, 0\}$$



# An Example (2)

Achieve desired illumination I<sub>des</sub> with bounded lamp powers

 $\begin{array}{ll} \min & \max_{k=1,\dots,n} |\log I_k - \log I_{\mathrm{des}}| \\ \mathrm{s.t.} & 0 \leq p_j \leq p_{\mathrm{max}}, j = 1,\dots,m \end{array}$ 





### An Example (3)

1. Use uniform power:  $p_j = p$ , vary p2. Use least-squares

min 
$$\sum_{i=1}^{k} (I_k - I_{des})^2 = \sum_{i=1}^{k} \left( \sum_{j=1}^{m} a_{kj} p_j - I_{des} \right)^2$$

Round  $p_j$  if  $p_j > p_{max}$  or  $p_j < 0$ 

3. Use weighted least-squares

min 
$$\sum_{i=1}^{k} (I_k - I_{des})^2 + \sum_{j=1}^{m} w_j \left( p_j - \frac{p_{max}}{2} \right)^2$$

Adjust weights  $w_j$  until  $0 \le p_j \le p_{max}$ 



### An Example (4)

4. Use linear programming min  $\max_{k=1,\dots,n} |I_k - I_{des}|$ s.t.  $0 \le p_j \le p_{max}, j = 1,\dots,m$ 

5. Use convex optimization

$$\begin{array}{ll} \min & \max_{k=1,\dots,n} |\log I_k - \log I_{des}| \\ \text{s.t.} & 0 \le p_j \le p_{\max}, j = 1,\dots,m \end{array}$$

$$\iff \min \qquad \max_{k=1,\dots,n} \left| \log \frac{I_k}{I_{des}} \right|$$
  
s.t.  $0 \le p_j \le p_{max}, j = 1,\dots,m$ 



# An Example (5)

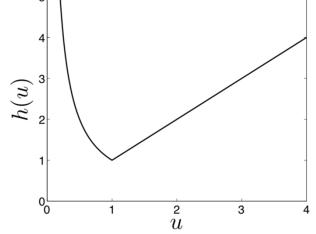
$$\iff \min \max_{k=1,\dots,n} \max\left(\log \frac{I_k}{I_{des}}, -\log \frac{I_k}{I_{des}}\right)$$
  
s.t.  $0 \le p_j \le p_{max}, j = 1, \dots, m$ 

$$\begin{array}{l} \min & \max_{k=1,\dots,n} \max\left(\log \frac{I_k}{I_{des}}, \log \frac{I_{des}}{I_k}\right) \\ \Leftrightarrow & \text{s.t.} \quad 0 \leq p_j \leq p_{\max}, j = 1,\dots,m \end{array}$$

$$\iff \min \max_{k=1,\dots,n} \max\left(\frac{I_k}{I_{des}}, \frac{I_{des}}{I_k}\right)$$
  
s.t.  $0 \le p_j \le p_{max}, j = 1, \dots, m$ 



# An Example (5) min $\max_{k=1,\dots,n} h\left(\frac{I_k}{I_{des}}\right)$ s.t. $0 \le p_j \le p_{\max}, j = 1, ..., m$ m $I_k = \sum_{j=1}^{k} a_{kj} p_j$ $h(u) = \max\left(u, \frac{1}{u}\right)$





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# Nonlinear Optimization

- An optimization problem when the objective or constraint functions are not linear, but not known to be convex
- Sadly, there are no effective methods for solving the general nonlinear programming problem
  - Could be NP-hard

□ We need compromise



### Local Optimization Methods

- □ Find a point that minimizes  $f_0$  among feasible points near it
  - The compromise is to give up seeking the optimal x
- □ Fast, can handle large problems
  - Differentiability
- Require initial guess, Provide no information about distance to (global) optimum, Sensitive to parameter values
- Local optimization methods are more art than technology



### Comparisons

	Problem Formulation	Solving the Problem
Local Optimization Methods for Nonlinear Programming	Straightforward	Art
Convex Optimization	Art	Standard



# Global Optimization (1)

 Find the global solution
 The compromise is efficiency
 Worst-case complexity grows exponentially with problem size

#### Applications

- A small number of variables, where computing time is not critical
- The value of finding the true global solution is very high



# Global Optimization (2)

- Worst-case Analysis of a high value or safety-critical system
  - Variables represent uncertain parameters
  - Objective function is a utility function
  - Constraints represent prior knowledge
  - If the worst-case value is acceptable, we can certify the system as safe or reliable

#### Local optimization methods can be tried

- If finding values that yield unacceptable performance, then the system is not reliable
- But it cannot certify the system as reliable

Role of Convex Optimization in **Nonconvex Problems** 



Initialization for local optimization

- An approximate, but convex, formulation
- Convex heuristics for nonconvex optimization
  - Sparse solutions (compressive sensing)
- Bounds for global optimization
  - Relaxation
  - Lagrangian relaxation



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- Least-squares
  - Closed-form Solution
- Linear Programming
  - Efficient algorithms
- Convex Optimization
  - Efficient algorithms, Modeling is an art
- Nonlinear Optimization
  - Compromises, Optimization is an art