



南 京 大 学  
人 工 智 能 学 院

SCHOOL OF ARTIFICIAL INTELLIGENCE, NANJING UNIVERSITY



# Towards Safe Evolutionary Optimization

Chao Qian

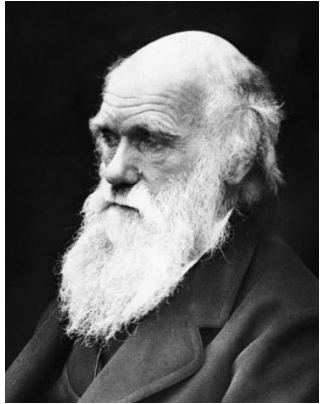
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# Evolutionary algorithms

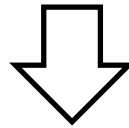
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Charles Darwin  
1809-1882

C. Darwin, after collecting abundant evidence, developed a theory about how species evolve

**reproduction with variation + nature selection**



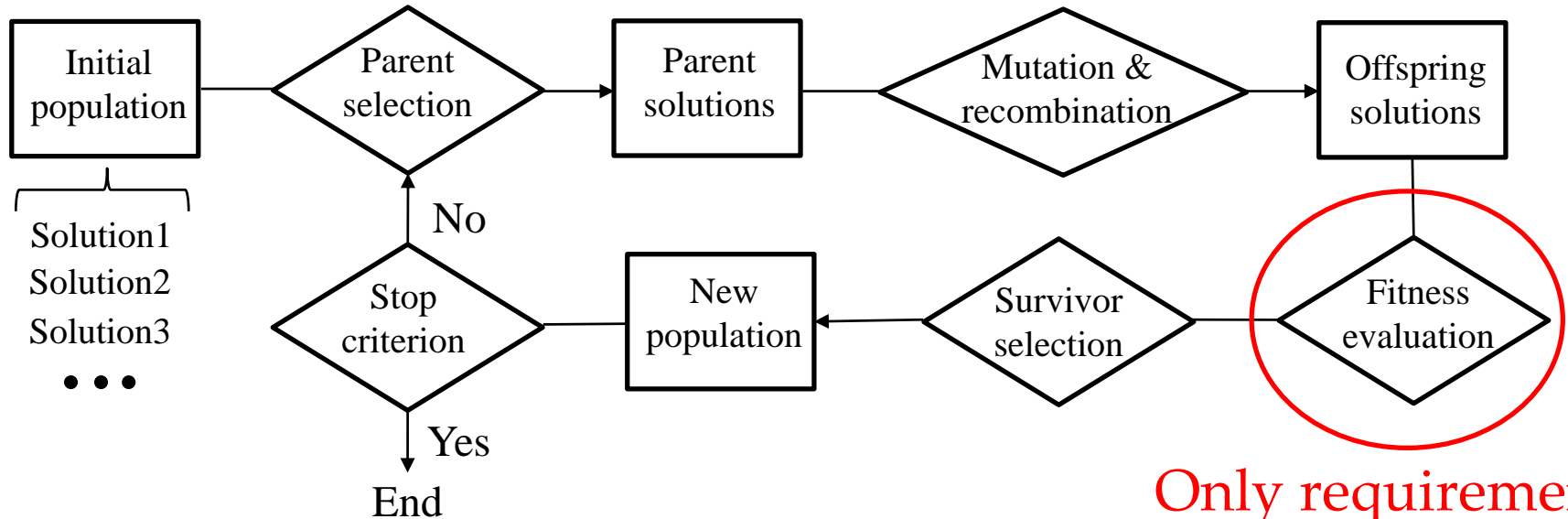
**Evolutionary algorithms:** a kind of nature-inspired randomized heuristic optimization algorithms

genetic algorithms, evolutionary strategies, evolutionary programming, particle swarm optimization, .....

# Evolutionary algorithms

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## Common structure of evolutionary algorithms (EAs)



EAs can be applied to solve complex optimization problems

- non-differentiable and non-continuous problems
- problems with multiple objective functions
- problems without explicit objective function formulation

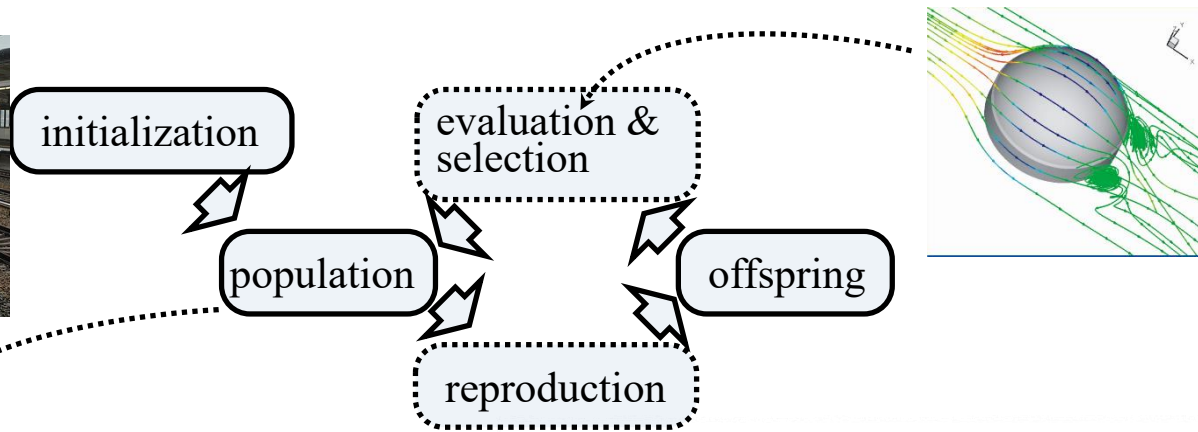
# Application - high-speed train head design



Series 700



Series N700



## Technological overview of the next generation Shinkansen high-speed train Series N700

M. Ueno<sup>1</sup>, S. Usui<sup>1</sup>, H. Tanaka<sup>1</sup>, A. Watanabe<sup>2</sup>

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### Abstract

In March 2005, Central Japan Railway Company (JR Central) has completed prototype trainset of the Series N700, the next generation Shinkansen high-speed rolling stock, developed aiming at easy handling system, and subjected to the problem of aerodynamic pressure waves and other issues related to environmental compatibility such as external noise. To combat this, an aero double-wing-type has been adopted for nose shape (Fig. 3). This nose shape, which boasts the most appropriate aerodynamic performance, has been newly developed for railway rolling stock using the latest analytical technique (i.e. genetic algorithms) used to develop the main wings of airplanes. The shape resembles a bird in flight, suggesting a feeling of boldness and speed.

the Tokaido Shinkansen line, Series N700 cars save 19% energy than Series 700 cars, achieve a 30% increase in the output of their traction equipment for higher-speed operation (Fig. 4).

this nose ... has been newly developed ... using the latest analytical technique (i.e. **genetic algorithms**)

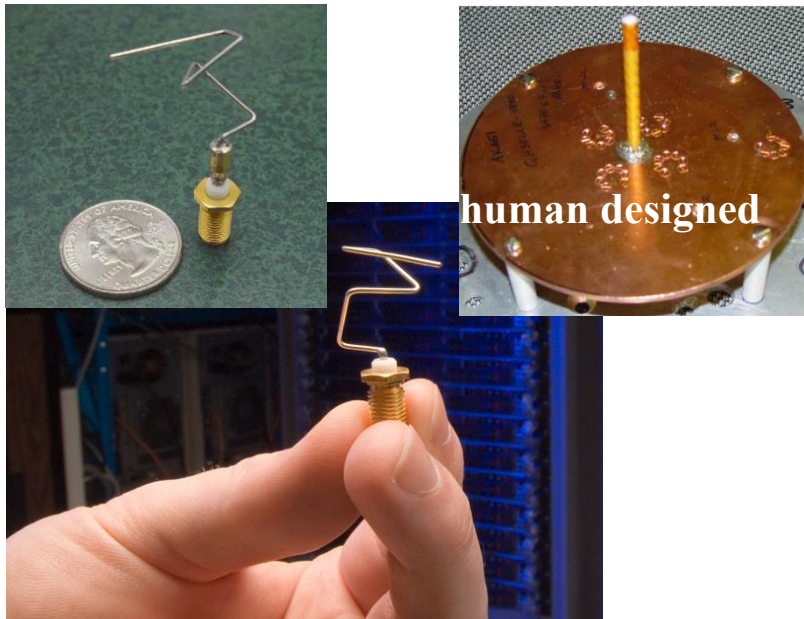
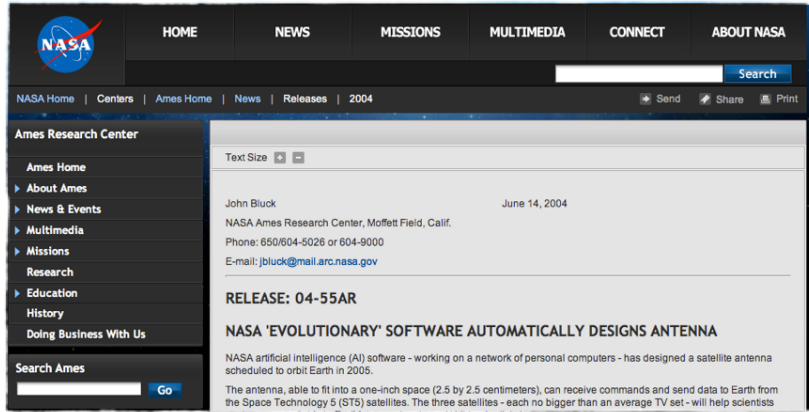
N700 cars save **19%** energy ... **30%** increase in the output...

This is a result of adopting the ... nose shape

This is a result of adopting the aerodynamically excellent nose shape, reduced running resistance thanks to the drastically smoothened car body and under-floor equipment, effective



# Application - antenna design



## Computer-Automated Evolution of an X-Band Antenna for NASA's Space Technology 5 Mission

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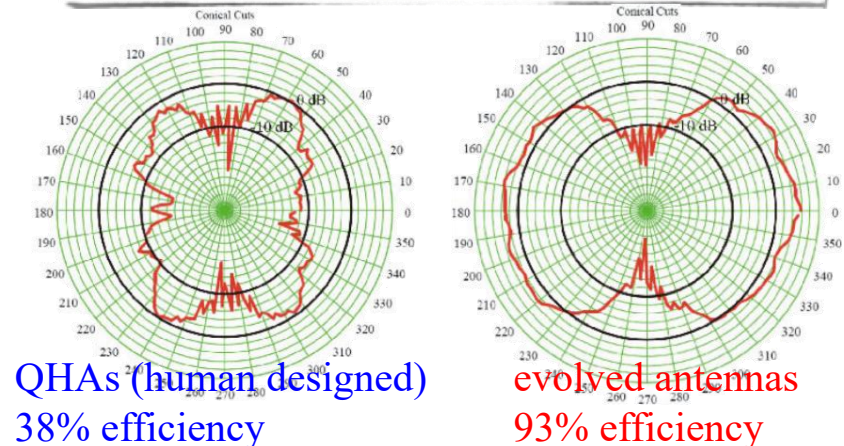
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dclinden@jemengineering.com

Since there are two antennas on each spacecraft, and not just one, it is important to measure the overall gain pattern with two antennas mounted on the spacecraft. For this, different combinations of the two evolved antennas and the QHA were tried on the ST5 mock-up and measured in an anechoic chamber. **With two QHAs 38% efficiency was achieved, using a QHA with an evolved antenna resulted in 80% efficiency, and using two evolved antennas resulted in 93% efficiency.** Here "efficiency" means how much power is being radiated versus how much power is being eaten up in resistance, with greater efficiency resulting in a stronger signal and greater range. Figure 11



# Application - biological evolution



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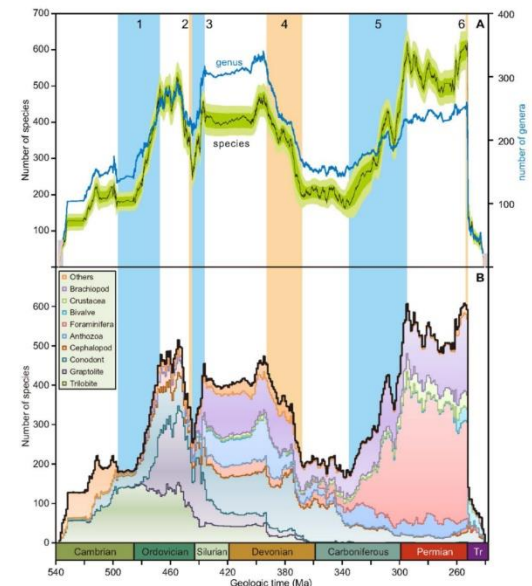
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首页 - 综合新闻

2020-01-17 作者: 地球科学与工程学院 来源: 地球科学与工程学院

## 《Science》刊登南京大学地球科学与工程学院研究成果：大数据和超算揭秘古生代海洋生物多样性演化

北京时间1月17日，国际权威期刊《Science》以研究长文的形式在线发表了南京大学、中国科学院南京地质古生物所樊隽轩教授、沈树忠院士等的论文“A high-resolution summary of Cambrian to Early Triassic marine invertebrate biodiversity”。该研究利用古生物大数据、超算和遗传算法等全新的方法和手段，基于化石记录重现了生命演化历史，改变了当前对古生代海洋生物多样性演化的认知。



# Safe evolutionary algorithms

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EAs have yielded encouraging empirical outcomes,  
but lack theoretical guarantee

Reported results

Data set 1	✓
Data set 2	✓
Data set 3	✓
Data set 4	✓
Data set 5	✓

How about the performance on other data?

Data set 6	?
Data set 7	?
Data set 8	?
Data set 9	?
Data set 10	?

Not safe

**Theoretical guarantee:** for any instance of a given problem,

function value of the output solution  $\leftarrow \textcircled{f(x)} \geq \alpha \cdot \textcircled{\text{OPT}} \rightarrow$  Optimal function value

# Safe evolutionary algorithms

---

EAs have yielded encouraging empirical outcomes,  
but lack theoretical guarantee

Reported results

Data set 1	✓
Data set 2	✓
Data set 3	✓
Data set 4	✓
Data set 5	✓

How about the performance on other data?

Data set 6	?
Data set 7	?
Data set 8	?
Data set 9	?
Data set 10	?

Not safe

Can we design “safe” EAs, i.e., EAs with provable approximation guarantee?



# Outline

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Safe evolutionary optimization?



Develop running time analysis tools for EAs

# Convergence analysis of EAs

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Convergence analysis

$$\lim_{t \rightarrow +\infty} P(\xi_t \in \mathcal{X}^*) = 1 ?$$

An EA that

1. uses global operators
2. preserves the best solution

converges to the optimal solutions

[Rudolph, 1998]

the probability of finding optimum in each step

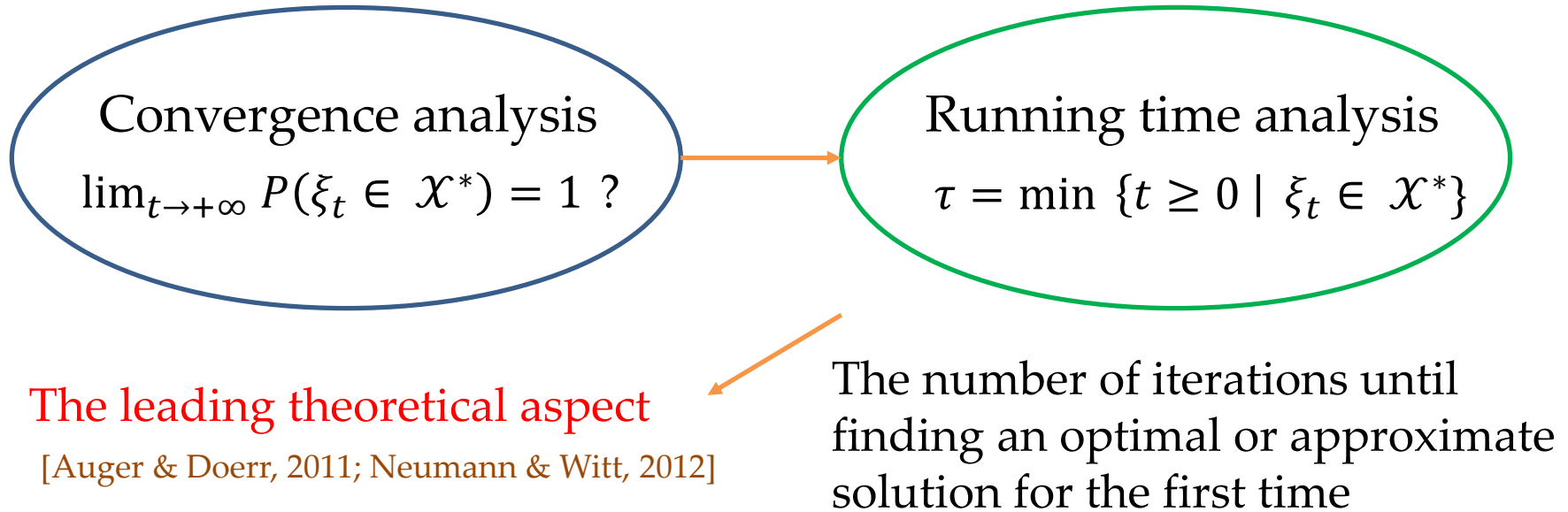
$$\Rightarrow \forall x: P(\xi_{t+1} \in \mathcal{X}^* \mid \xi_t = x) > 0$$



$$\lim_{t \rightarrow +\infty} P(\xi_t \in \mathcal{X}^*) = 1$$

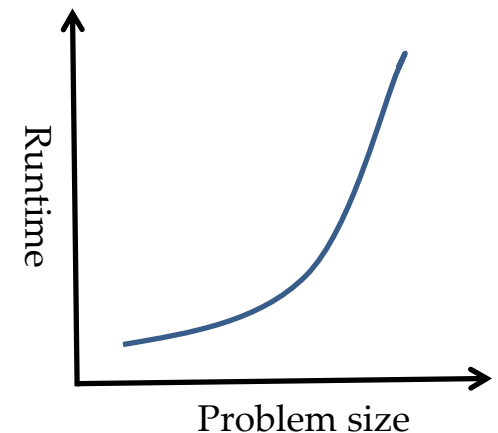
But life is limited! How fast does it converge?

# Running time analysis of EAs



## Running time complexity

- Usually grows with the problem size and expressed in asymptotic notations  
e.g., (1+1)-EA solving LeadingOnes:  $O(n^2)$



# Running time analysis tools of EAs

- Analyses starting from scratch are quite difficult
- We need **general running time analysis tools**

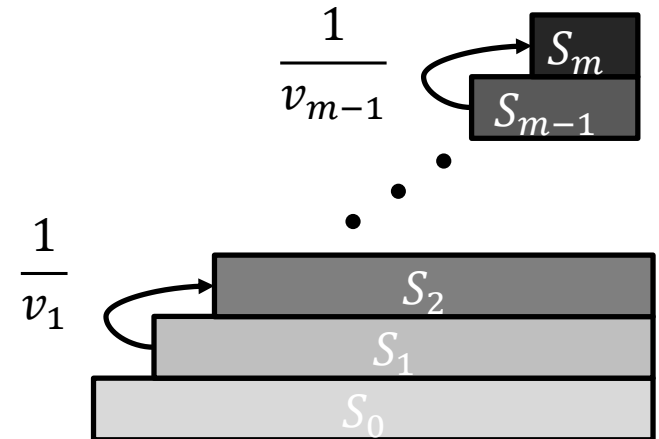
## Fitness level method [Droste et al., TCS'02]

- Jumping probability lower bound:

$$P(\xi_{t+1} \in \cup_{j=i+1}^m S_j | \xi_t \in S_i) \geq v_i$$



- Running time upper bound:  $\sum_{i=j}^{m-1} 1/v_i$



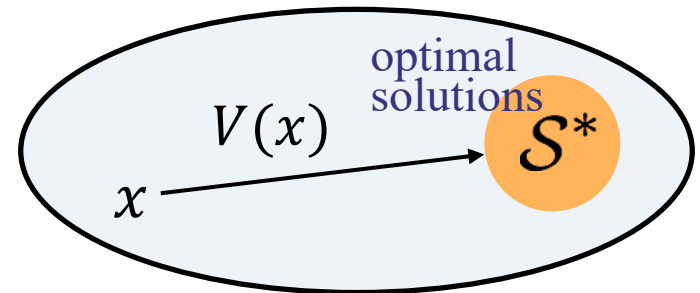
## Drift analysis [He & Yao, AIJ'01]

- Expected drift in one step:

$$E[V(\xi_t) - V(\xi_{t+1}) | \xi_t] \geq c$$



- Running time upper bound:  $V(x)/c$



# Switch analysis

## Switch analysis [Yu, Qian & Zhou, TEC'15]

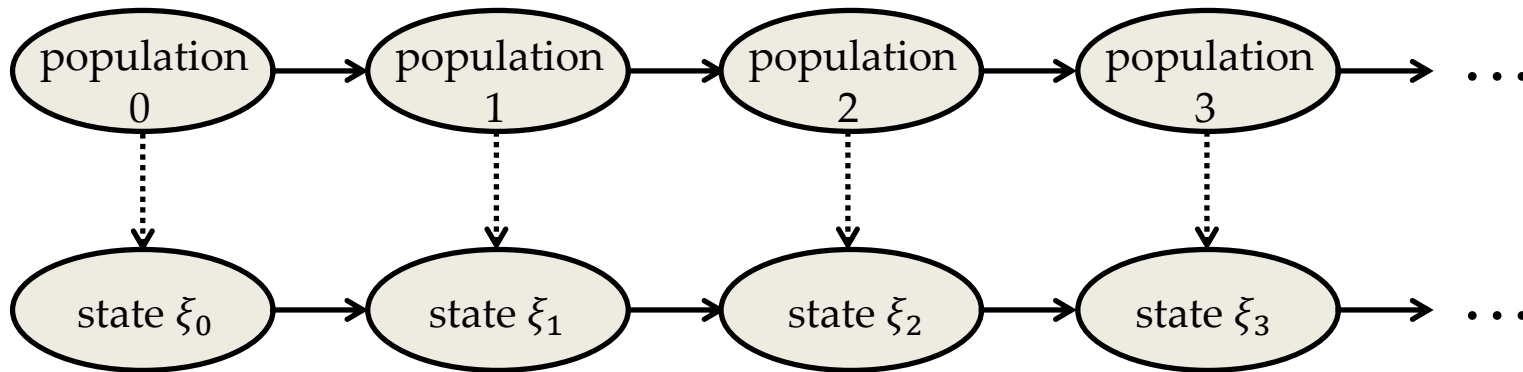
- One step difference with a reference evolutionary process:

$$\forall t: \sum_{x \in \mathcal{X}, y \in \mathcal{Y}} \pi_t(x) P(\xi_{t+1} \in \phi^{-1}(y) | \xi_t = x) E[\tau' | \xi'_0 = y] \\ - \sum_{u, y \in \mathcal{Y}} \pi_t^\phi(u) P(\xi'_1 \in y | \xi'_0 = u) E[\tau' | \xi'_1 = y] \leq \rho_t$$



- Running time upper bound:  $E[\tau | \xi_0 \sim \pi_0] \leq E[\tau' | \xi'_0 \sim \pi_0^\phi] + \sum_{t=0}^{+\infty} \rho_t$

## Main idea:

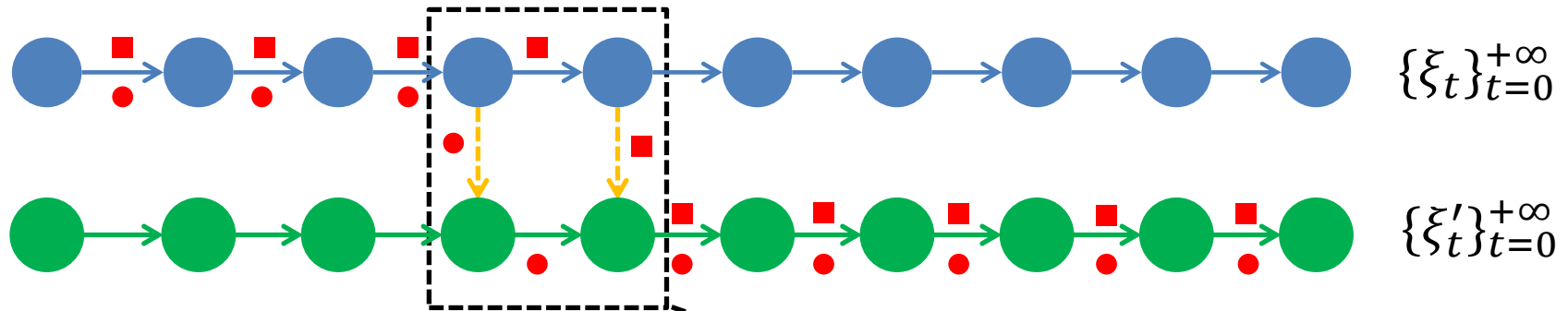




# Switch analysis

**Main idea** [Yu, Qian and Zhou, TEC'15]:

Given EA on the given problem



Reference chain

The expected running time of  $\{\xi_t\}_{t=0}^{+\infty}$ :

$$E[\tau] \leq (\geq) E[\tau'] + \sum_{t=0}^{+\infty} \rho_t$$

The expected running time of  $\{\xi'_t\}_{t=0}^{+\infty}$ , easy to analyze

# Ability of switch analysis

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## Fitness level method [Droste et al., TCS'02]

- Jumping probability lower bound:

$$P(\xi_{t+1} \in \bigcup_{j=i+1}^m S_j | \xi_t \in S_i) \geq v_i$$



- Running time upper bound:  $\sum_{i=j}^{m-1} 1/v_i$

## Drift analysis [He & Yao, AIJ'01]

- Expected drift in one step:

$$E[V(\xi_t) - V(\xi_{t+1}) | \xi_t] \geq c$$



- Running time upper bound:  $V(x)/c$

Reducible to  
switch analysis

Switch analysis can derive at least the same tight bound while requiring no more information

# Application of switch analysis

Application [Bian, Qian and Tang, IJCAI'18]:

GSEMO	Problem	Previous result	Our result
Bi-objective	LOTZ	$O(n^3)$ [Giel, CEC'03]	$\leq 6n^3$
	COCZ	$O(n^2 \log n)$ [Qian et al., AIJ'13]	$\leq 3n^2 \log n$
Many-objective	mCOCZ	$O(n^{m+1})$ [Laumanns et al., TEC'04]	$O(n^m)$ for $m > 4$ , $O(n^3 \log n)$ for $m = 4$
Approximate analysis	WOMM	—	1/n-approximation: $O(n^2(\log_l n + \log_l(w_n/w_1)))$

→ gives the leading constants

→ is asymptotically tighter than



**L. Thiele**

Member of  
Academia  
Europaea

Switch analysis is general and powerful

# Outline

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Safe evolutionary optimization?



Develop running time analysis tools for EAs



Disclose theoretical properties of EAs  
for constrained and noisy optimization

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# How to deal with constraints when EAs are used for constrained optimization?

The optimization problems in real-world applications often come with constraints



# Constrained optimization

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## General formulation:

$$\begin{array}{ll} \min_{x \in \mathcal{X}} & f(x) \\ \text{s.t.} & g_i(x) = 0, \quad 1 \leq i \leq q; \\ & h_i(x) \leq 0, \quad q + 1 \leq i \leq m \end{array}$$

objective function

equality constraints

inequality constraints

**The goal:** find a feasible solution minimizing the objective  $f$

satisfies all constraints

# Constraint handling strategies

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Two common constraint handling strategies:

- ❑ Penalty function [Hadj-Alouane & Bean, OR'97]
  - transform the original constrained optimization problem into an unconstrained optimization problem
- ❑ Multi-objective reformulation [Coello Coello, 2002; Cai & Wang, TEC'06]
  - transform the original constrained optimization problem into a bi-objective optimization problem

# Penalty function

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## Main idea [Hadj-Alouane & Bean, OR'97]

1. transform the original **constrained** optimization problem into an **unconstrained** optimization problem

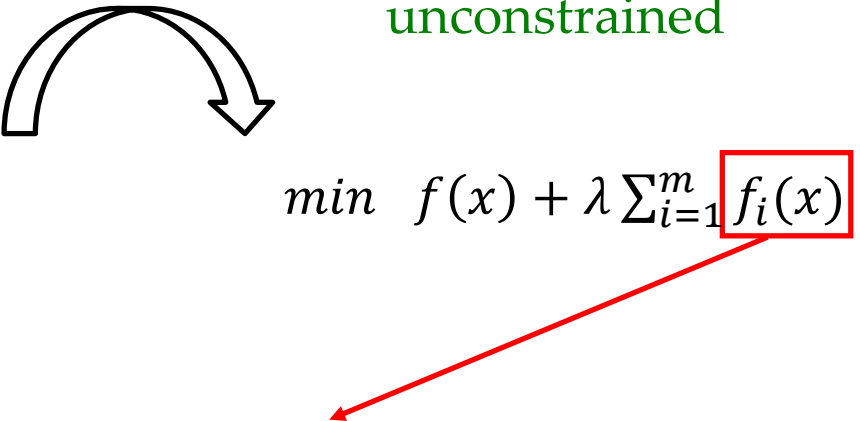
**constrained**

$$\begin{aligned} \min \quad & f(x) \\ \text{s.t.} \quad & g_i(x) = 0, \quad 1 \leq i \leq q; \\ & h_i(x) \leq 0, \quad q + 1 \leq i \leq m \end{aligned}$$

**unconstrained**

$$\min \quad f(x) + \lambda \sum_{i=1}^m f_i(x)$$

constraint violation degree

$$f_i(x) = \begin{cases} |g_i(x)| & 1 \leq i \leq q \\ \max\{0, h_i(x)\} & q + 1 \leq i \leq m \end{cases}$$


# Penalty function


---

## Main idea [Hadj-Alouane & Bean, OR'97]

1. transform the original **constrained** optimization problem into an **unconstrained** optimization problem

**constrained**

$$\begin{aligned} \min \quad & f(x) \\ \text{s.t.} \quad & g_i(x) = 0, \quad 1 \leq i \leq q; \\ & h_i(x) \leq 0, \quad q + 1 \leq i \leq m \end{aligned}$$



**unconstrained**

$$\min \quad f(x) + \lambda \sum_{i=1}^m f_i(x)$$

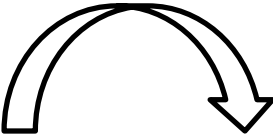
2. employ a single-objective EA to solve the transformed problem

# Multi-objective reformulation

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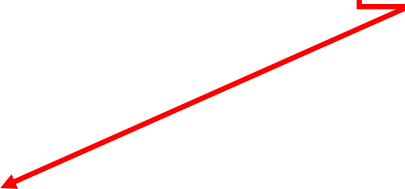
## Main idea [Coello Coello, 2002; Cai & Wang, TEC'06]

1. transform the original **constrained** optimization problem into a **bi-objective** optimization problem

**constrained**  **bi-objective**

$$\begin{array}{ll} \min & f(x) \\ \text{s.t.} & g_i(x) = 0, \quad 1 \leq i \leq q; \\ & h_i(x) \leq 0, \quad q + 1 \leq i \leq m \end{array}$$
$$\min (f(x), \sum_{i=1}^m \boxed{f_i(x)})$$

constraint violation degree  $f_i(x) = \begin{cases} |g_i(x)| & 1 \leq i \leq q \\ \max\{0, h_i(x)\} & q + 1 \leq i \leq m \end{cases}$





# Multi-objective optimization

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**The task:** optimize multiple objectives simultaneously

$$\min_{x \in \mathcal{X}} (f_1(x), f_2(x), \dots, f_m(x))$$

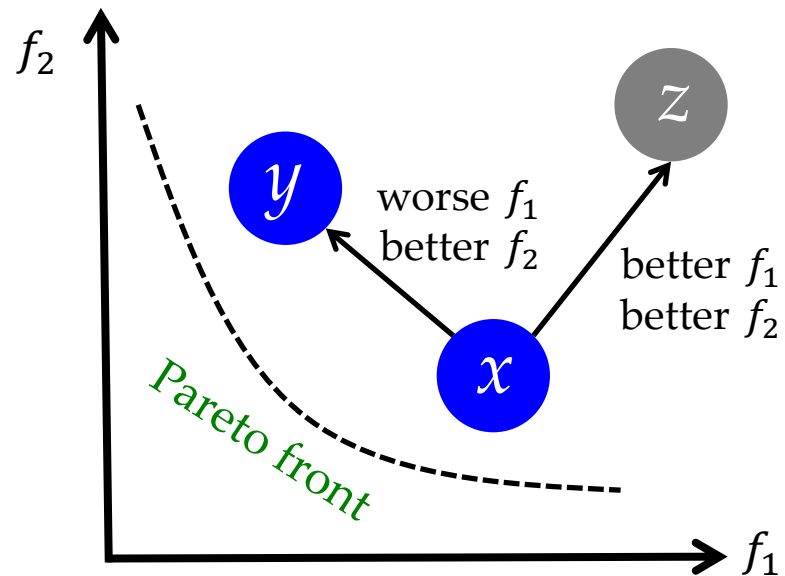
Example: bi-objective minimization

$x$  dominates  $z$  :

$$f_1(x) < f_1(z) \wedge f_2(x) < f_2(z)$$

$x$  is incomparable with  $y$  :

$$f_1(x) > f_1(y) \wedge f_2(x) < f_2(y)$$



# Multi-objective reformulation

---

Main idea [Coello Coello, 2002; Cai & Wang, TEC'06]

1. transform the original **constrained** optimization problem into a **bi-objective** optimization problem

**constrained**

$$\begin{aligned} \min \quad & f(x) \\ \text{s.t.} \quad & g_i(x) = 0, \quad 1 \leq i \leq q; \\ & h_i(x) \leq 0, \quad q + 1 \leq i \leq m \end{aligned}$$

↷

**bi-objective**

$$\min (f(x), \sum_{i=1}^m f_i(x))$$

2. employ a multi-objective EA to solve the transformed problem
3. output the feasible solution from the generated non-dominated solution set

constraint violation degree = 0

# Constraint handling strategies

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Two common constraint handling strategies:

- ❑ Penalty function [Hadj-Alouane & Bean, OR'97]
- ❑ Multi-objective reformulation [Coello Coello, 2002; Cai & Wang, TEC'06]

Previous empirical studies have shown the superior performance of multi-objective reformulation

It is not yet clear whether multi-objective reformulation can be better in theory

# Problems

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- Minimum matroid optimization (P-solvable) [Edmonds, MP'71]  
e.g., minimum spanning tree, maximum bipartite matching

**Definition 1.** Given a matroid  $(U, S)$ , a rank function  $r: 2^U \rightarrow \mathbb{N}$  and a weight function  $w: U \rightarrow \mathbb{N}$ , the problem is formulated as

$$\min_{x \in \{0,1\}^n} \sum_{i=1}^n w_i x_i \quad \text{s.t.} \quad r(x) = r(U)$$

- Minimum cost coverage (NP-hard) [Wolsey, Combinatorica'82]  
e.g., minimum set cover, submodular set cover

**Definition 2.** Given a monotone submodular function  $f: 2^U \rightarrow \mathbb{R}$ , some value  $q \leq f(U)$  and a weight function  $w: U \rightarrow \mathbb{N}$ , the problem is formulated as

$$\min_{x \in \{0,1\}^n} \sum_{i=1}^n w_i x_i \quad \text{s.t.} \quad f(x) \geq q$$

# Theoretical analysis

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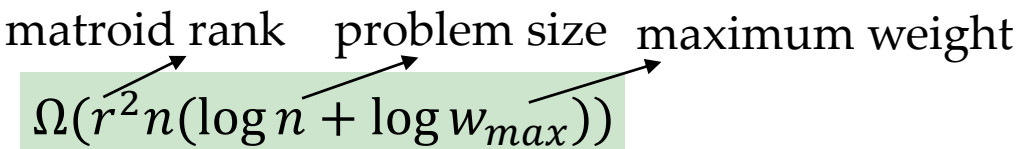
## Penalty function vs. Multi-objective reformulation

[Qian, Yu and Zhou, IJCAI'15]

- Minimum matroid optimization (P-solvable): obtaining an optimal solution

Penalty function:  $\Omega(r^2 n (\log n + \log w_{\max}))$

matroid rank      problem size      maximum weight



Multi-objective reformulation:  $O(rn(\log n + \log w_{\max} + r))$

The running time reduces by a factor

$$\min\{\log n + \log w_{\max}, r\}$$



# Theoretical analysis

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- Minimum matroid optimization (P-solvable): obtaining an optimal solution

Penalty function:  $\Omega(r^2 n (\log n + \log w_{\max}))$

Multi-objective reformulation:  $O(rn (\log n + \log w_{\max} + r))$

The running time reduces by a factor  $\min\{\log n + \log w_{\max}, r\}$

- Minimum cost coverage (NP-hard): obtaining a  $H_q$ -approximate solution

Penalty function: exponential w.r.t.  $n, q, \log w_{\max}$

Multi-objective reformulation:  $O(qn (\log n + \log w_{\max} + q))$

The running time reduces exponentially  polynomial

# Theoretical analysis

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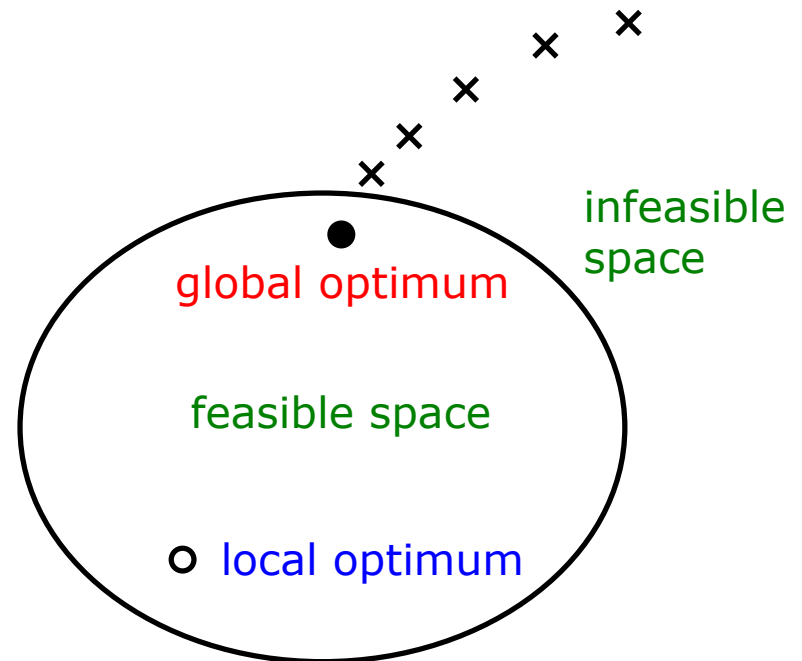
## Findings from the analysis:

### Penalty function

- the penalty prefers feasible solutions
- get trapped in the local optimum, which is far from the global optimum

### Multi-objective reformulation

- the constraint violation objective allows infeasible solutions
- follow a short path from infeasible to feasible to find good solutions



# General problem classes

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[Qian et al., AIJ'19]

- Constrained submodular approximately monotone maximization

Multi-objective reformulation:  $f(x) \geq \left(1 - \frac{1}{e}\right) \cdot (\text{OPT} - k\epsilon)$

- Constrained monotone approximately submodular maximization

Multi-objective reformulation:  $f(x) \geq \frac{1}{1 + \frac{2k\epsilon}{1 - \epsilon}} \left(1 - \frac{1}{e} \left(\frac{1 - \epsilon}{1 + \epsilon}\right)^k\right) \cdot \text{OPT}$

Achieve the best known polynomial-time approximation guarantee [Krause et al., JMLR'08]

[Horel & Singer, NIPS'16]

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# How to improve the robustness when EAs are used for noisy optimization?

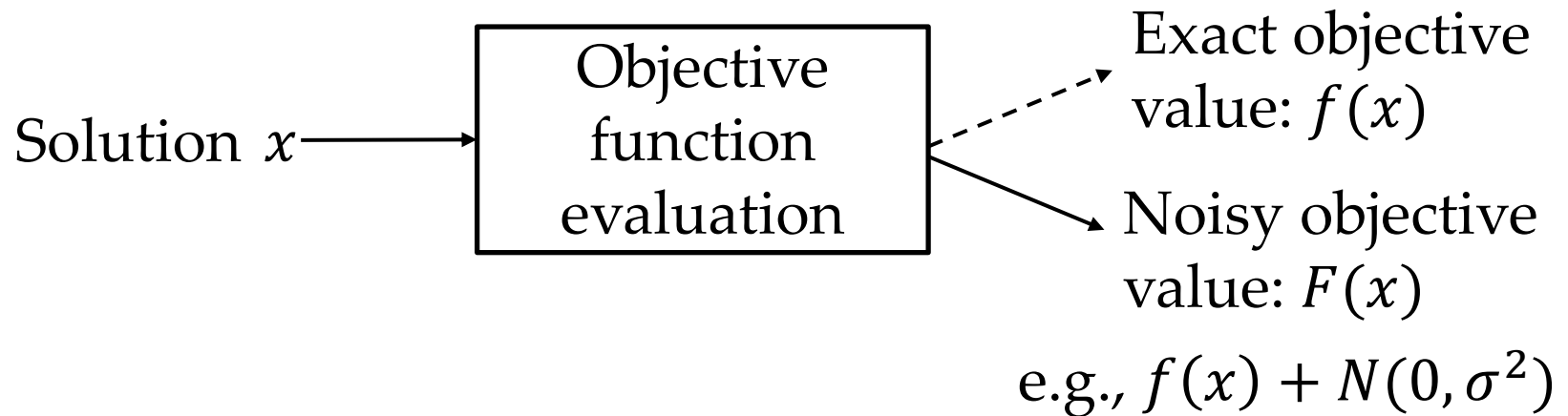
The optimization problems in real-world applications often come with noise

# Noisy optimization

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The objective evaluation is often disturbed by noise

e.g., a prediction model is evaluated only on a limited amount of data



# The influence of noise

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It was believed that noise makes evolutionary optimization harder

many noise handling strategies have been proposed

[Jin & Branke, TEC'05; Goh & Tan, TEC'07]

Some empirical observations have shown that noise can have a positive impact on the performance of local search

[Selman et al., AAAI'94; Hoos & Stutzle, JAR'00]

Can noise make evolutionary optimization easier?

# Theoretical analysis

**A sufficient condition: noise is helpful** [Qian, Yu and Zhou, ECJ'18]

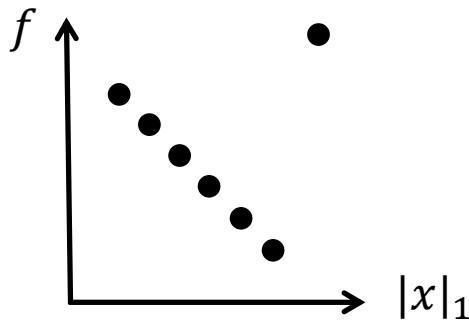
**Theorem 1.** For an EA  $\mathcal{A}$  optimizing a problem  $f$ , which can be modeled by a deceptive Markov chain, if

$$\forall x \notin \mathcal{X}_0 : P_{\xi}^t(x, \mathcal{X}_0) = \sum_{x' \cap \mathcal{S}^* \neq \emptyset} P_{var}(x, x'), \quad (6)$$

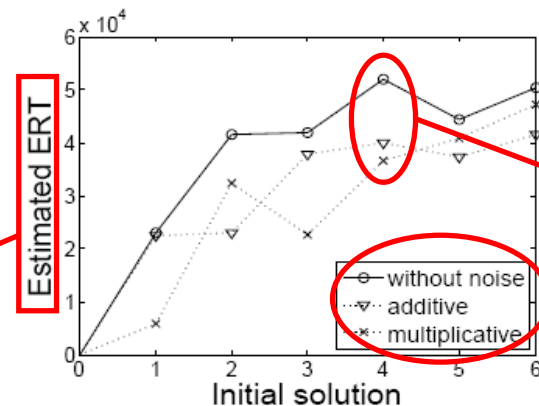
then noise makes  $f$  easier for  $\mathcal{A}$ .

Intuitively, noise can bring some randomness to help the EA escape from local optima

**Example:**  $(1+n)$ -EA on the Trap problem



Running time



Noise  
helpful

# Noise handling strategies

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## Noise is harmful in most cases

Two commonly used noise handling strategies:

- ❑ Re-evaluation [Arnold & Beyer, TEC'02; Jin & Branke, TEC'05]
  - every time we access the fitness of a solution by evaluation  
smooth noise
- ❑ Threshold selection [Markon et al., CEC'01; Bartz-Beielstein & Markon, CEC'02]
  - an offspring solution is accepted only if its fitness is larger than that of the parent solution by at least a threshold  $\tau$   
reduce the risk of accepting a bad solution due to noise



# Theoretical analysis

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the range of noise level such that the running time is polynomial

Example:

(1+1)-EA

OneMax

one-bit noise

Noise handling strategies	PNT
single-evaluation	$[0, 1 - \frac{1}{\Theta(\text{poly}(n))}]$
single-evaluation & $\tau > 0$	$[0, 0]$
re-evaluation	$[0, \Theta(\frac{\log n}{n})]$ (Droste, 2004)
re-evaluation & $\tau = 1$	$[0, 1]$
re-evaluation & $\tau = 2$	$[\frac{1}{\Theta(\text{poly}(n))}, 1 - \frac{1}{\Theta(\text{poly}(n))}]$
re-evaluation & $\tau > 2$	$\emptyset$

combining re-evaluation with proper threshold selection is better [Qian, Yu and Zhou, ECJ'18]

# Outline

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Safe evolutionary optimization?



Develop running time analysis tools for EAs



Disclose theoretical properties of EAs  
for constrained and noisy optimization



Propose EAs with good polynomial-time  
approximation guarantees for (noisy) subset selection

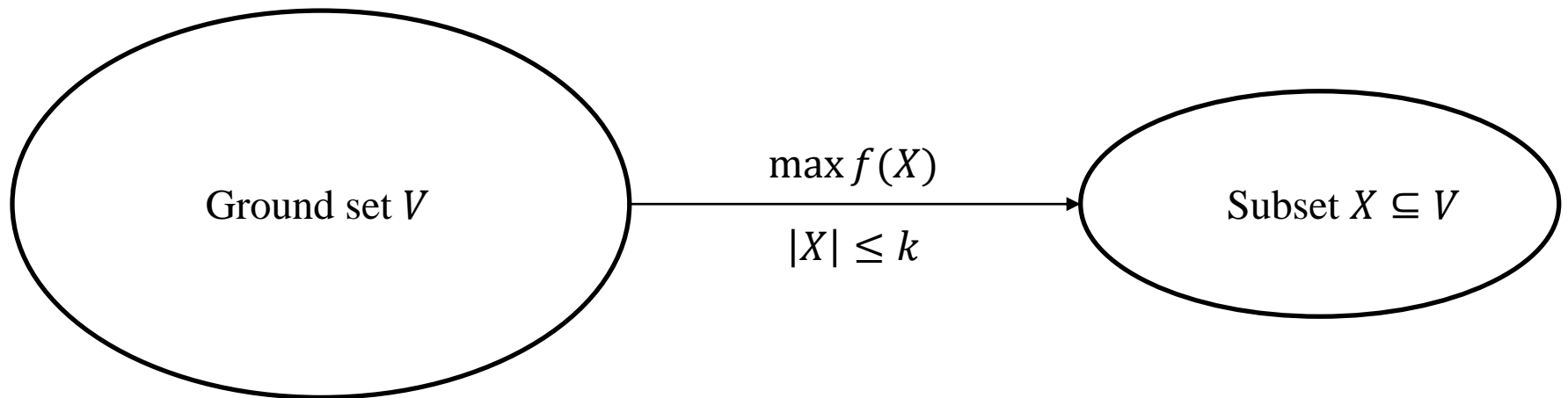
# Subset selection

---

**Subset selection** is to select a subset of size  $k$  from a total set of  $n$  items for optimizing some objective function

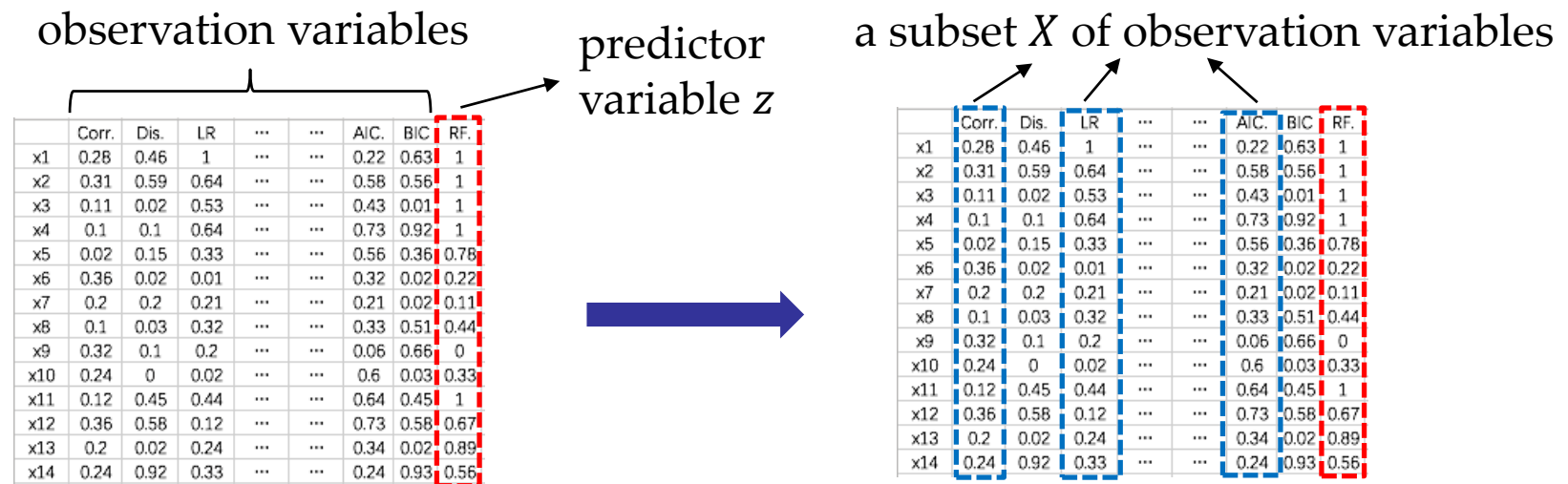
**Formally stated:** given all items  $V = \{v_1, \dots, v_n\}$ , an objective function  $f: 2^V \rightarrow \mathbb{R}$  and a budget  $k$ , to find a subset  $X \subseteq V$  such that

$$\max_{X \subseteq V} f(X) \quad \text{s.t.} \quad |X| \leq k.$$



# Application - sparse regression

**Sparse regression** [Tropp, TIT'04]: select a few observation variables to best approximate the predictor variable by linear regression



Item  $v_i$ : an observation variable

Objective  $f$ : squared multiple correlation  $R_{z,X}^2 = \frac{\text{Var}(z) - \text{MSE}_{z,X}}{\text{Var}(z)}$

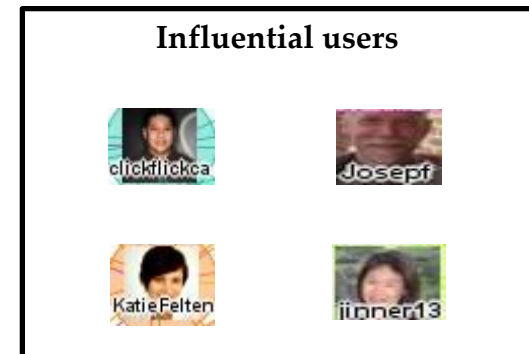
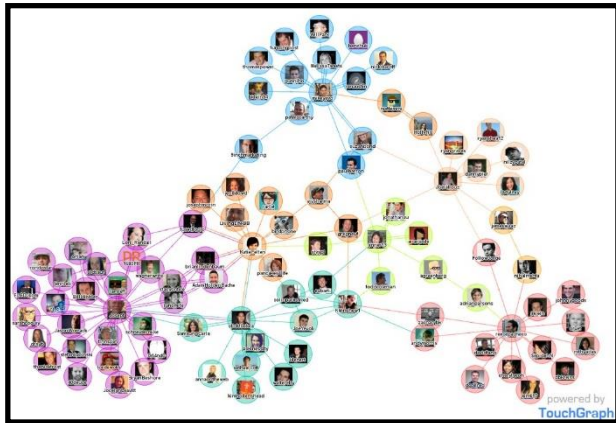
variance

mean squared error

# Application - influence maximization

---

**Influence maximization** [Kempe et al., KDD'03] : select a subset of users from a social network to maximize its influence spread

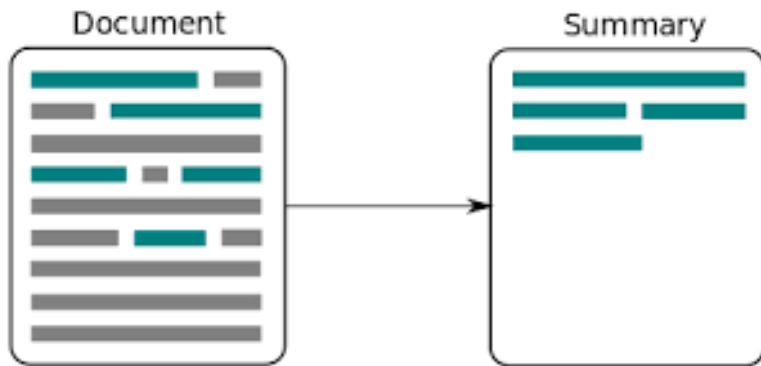


Item  $v_i$ : a social network user

Objective  $f$ : influence spread, measured by the expected number of social network users activated by diffusion

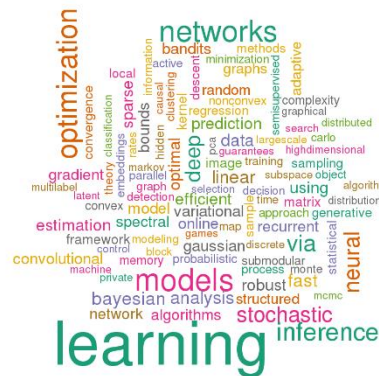
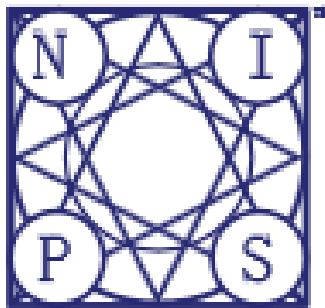
# Application - document summarization

**Document summarization** [Lin & Bilmes, ACL'11] : select a few sentences to best summarize the documents



Item  $v_i$ : a sentence

Objective  $f$ : summary quality

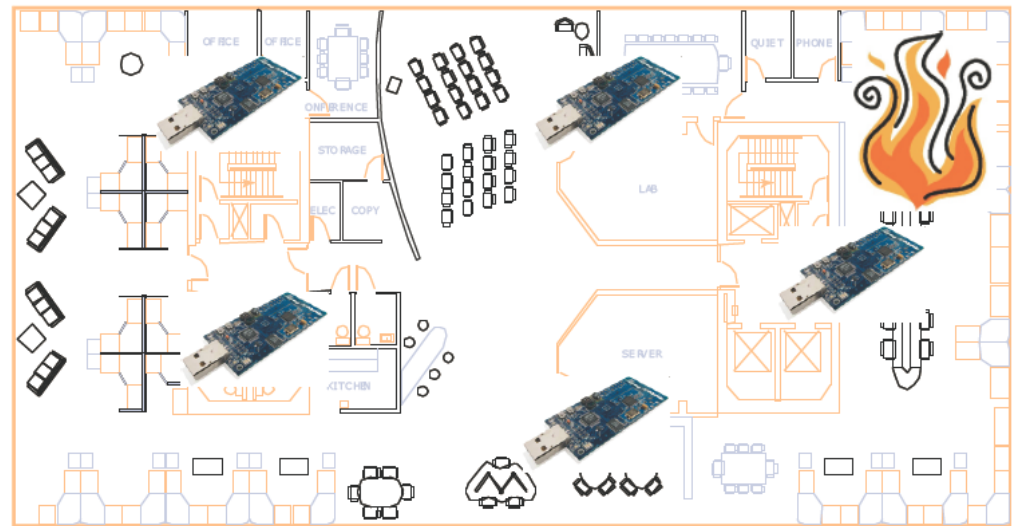


# Application - sensor placement

**Sensor placement** [Krause & Guestrin, IJCAI'09 Tutorial] : select a few places to install sensors such that the information gathered is maximized



Water contamination detection

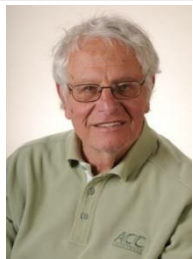
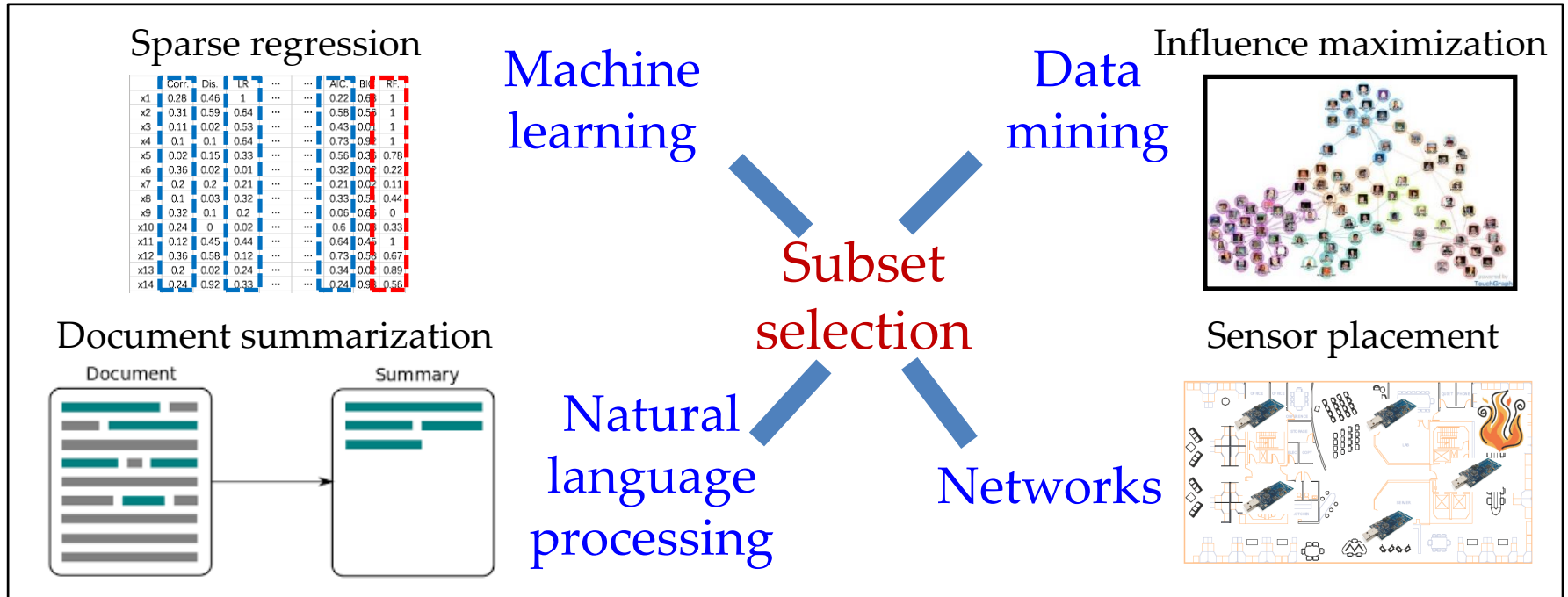


Fire detection

Item  $v_i$ : a place to install a sensor

Objective  $f$ : entropy

# Subset selection



**George Nemhauser**

John Von Neumann  
Theory Prize

[Mathematical Programming 1978]

$f$ : monotone and submodular

The greedy algorithm:

$(1 - 1/e)$ -approximation

Best Paper/Test of  
Time Award:

[Kempe et al., KDD'03]

[Das & Kempe, ICML'11]

[Iyer & Bilmes, NIPS'13]





# POSS algorithm

POSS algorithm [Qian, Yu and Zhou, NIPS'15]

Transformation:

$$\begin{array}{ll} \max_{X \subseteq V} f(X) \quad s.t. \quad |X| \leq k & \text{original} \\ \Downarrow & \\ \min_{X \subseteq V} (-f(X), |X|) & \text{bi-objective} \end{array}$$

---

## Algorithm 1 POSS

---

**Input:** all variables  $V = \{X_1, \dots, X_n\}$ , a given objective  $f$  and an integer parameter  $k \in [1, n]$

**Parameter:** the number of iterations  $T$

**Output:** a subset of  $V$  with at most  $k$  variables

**Process:**

1: Let  $s = \{0\}^n$  and  $P = \{s\}$ .

2: Let  $t = 0$ .

3: **while**  $t < T$  **do**

4:   Select  $s$  from  $P$  uniformly at random.

5:   Generate  $s'$  by flipping each bit of  $s$  with prob.  $\frac{1}{n}$ .

6:   Evaluate  $f_1(s')$  and  $f_2(s')$ .

7:   **if**  $\nexists z \in P$  such that  $z \prec s'$  **then**

8:      $Q = \{z \in P \mid s' \preceq z\}$ .

9:      $P = (P \setminus Q) \cup \{s'\}$ .

10:   **end if**

11:    $t = t + 1$ .

12: **end while**

13: **return**  $\arg \min_{s \in P, |s| \leq k} f_1(s)$

---

**Initialization:** put the special solution  $\{0\}^n$  into the population  $P$

**Reproduction:** pick a solution  $x$  randomly from  $P$ , and flip each bit of  $x$  with prob.  $1/n$  to generate a new solution

**Evaluation & Updating:** if the new solution is not dominated, put it into  $P$  and weed out bad solutions

**Output:** select the best feasible solution

# Sparse regression

**Sparse regression:** given all observation variables  $V = \{v_1, \dots, v_n\}$ , a predictor variable  $z$  and a budget  $k$ , to find a subset  $X \subseteq V$  such that

$$\max_{X \subseteq V} R_{z,X}^2 = \frac{\text{Var}(z) - \text{MSE}_{z,X}}{\text{Var}(z)} \quad \text{s.t.} \quad |X| \leq k$$

$\text{Var}(z)$ : variance of  $z$

$\text{MSE}_{z,X}$ : mean squared error of predicting  $z$  by using observation variables in  $X$

observation variables

	Corr.	Dis.	LR	...	...	AIC.	BIC	RF.
x1	0.28	0.46	1	...	...	0.22	0.63	1
x2	0.31	0.59	0.64	...	...	0.58	0.56	1
x3	0.11	0.02	0.53	...	...	0.43	0.01	1
x4	0.1	0.1	0.64	...	...	0.73	0.92	1
x5	0.02	0.15	0.33	...	...	0.56	0.36	0.78
x6	0.36	0.02	0.01	...	...	0.32	0.02	0.22
x7	0.2	0.2	0.21	...	...	0.21	0.02	0.11
x8	0.1	0.03	0.32	...	...	0.33	0.51	0.44
x9	0.32	0.1	0.2	...	...	0.06	0.66	0
x10	0.24	0	0.02	...	...	0.6	0.03	0.33
x11	0.12	0.45	0.44	...	...	0.64	0.45	1
x12	0.36	0.58	0.12	...	...	0.73	0.58	0.67
x13	0.2	0.02	0.24	...	...	0.34	0.02	0.89
x14	0.24	0.92	0.33	...	...	0.24	0.93	0.56

predictor variable  $z$

a subset  $X$  of observation variables

	Corr.	Dis.	LR	...	...	AIC.	BIC	RF.
x1	0.28	0.46	1	...	...	0.22	0.63	1
x2	0.31	0.59	0.64	...	...	0.58	0.56	1
x3	0.11	0.02	0.53	...	...	0.43	0.01	1
x4	0.1	0.1	0.64	...	...	0.73	0.92	1
x5	0.02	0.15	0.33	...	...	0.56	0.36	0.78
x6	0.36	0.02	0.01	...	...	0.32	0.02	0.22
x7	0.2	0.2	0.21	...	...	0.21	0.02	0.11
x8	0.1	0.03	0.32	...	...	0.33	0.51	0.44
x9	0.32	0.1	0.2	...	...	0.06	0.66	0
x10	0.24	0	0.02	...	...	0.6	0.03	0.33
x11	0.12	0.45	0.44	...	...	0.64	0.45	1
x12	0.36	0.58	0.12	...	...	0.73	0.58	0.67
x13	0.2	0.02	0.24	...	...	0.34	0.02	0.89
x14	0.24	0.92	0.33	...	...	0.24	0.93	0.56

# Experimental results

the size constraint:  $k = 8$

the number of iterations of POSS:  $2ek^2n$

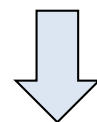
exhaustive search

greedy algorithms

relaxation methods

Data set	OPT	POSS	FR	FoBa	OMP	RFE	MCP
housing	.7437±.0297	.7437±.0297	.7429±.0300●	.7423±.0301●	.7415±.0300●	.7388±.0304●	.7354±.0297●
eunite2001	.8484±.0132	.8482±.0132	.8348±.0143●	.8442±.0144●	.8349±.0150●	.8424±.0153●	.8320±.0150●
svmguide3	.2705±.0255	.2701±.0257	.2615±.0260●	.2601±.0279●	.2557±.0270●	.2136±.0325●	.2397±.0237●
ionosphere	.5995±.0326	.5990±.0329	.5920±.0352●	.5929±.0346●	.5921±.0353●	.5832±.0415●	.5740±.0348●
sonar	–	.5365±.0410	.5171±.0440●	.5138±.0432●	.5112±.0425●	.4321±.0636●	.4496±.0482●
triazines	–	.4301±.0603	.4150±.0592●	.4107±.0600●	.4073±.0591●	.3615±.0712●	.3793±.0584●
coil2000	–	.0627±.0076	.0624±.0076●	.0619±.0075●	.0619±.0075●	.0363±.0141●	.0570±.0075●
mushrooms	–	.9912±.0020	.9909±.0021●	.9909±.0022●	.9909±.0022●	.6813±.1294●	.8652±.0474●
clean1	–	.4368±.0300	.4169±.0299●	.4145±.0309●	.4132±.0315●	.1596±.0562●	.3563±.0364●
w5a	–	.3376±.0267	.3319±.0247●	.3341±.0258●	.3313±.0246●	.3342±.0276●	.2694±.0385●
gisette	–	.7265±.0098	.7001±.0116●	.6747±.0145●	.6731±.0134●	.5360±.0318●	.5709±.0123●
farm-ads	–	.4217±.0100	.4196±.0101●	.4170±.0113●	.4170±.0113●	–	.3771±.0110●
POSS: win/tie/loss	–	–	12/0/0	12/0/0	12/0/0	11/0/0	12/0/0

- denotes that POSS is significantly better by the  $t$ -test with confidence level 0.05



**POSS is significantly better than all the compared state-of-the art algorithms on all data sets**

# Theoretical analysis

---

POSS can achieve the optimal polynomial-time approximation guarantee

**Theorem 1.** For subset selection with monotone objective functions, POSS using  $E[T] \leq 2ek^2n$  finds a solution  $X$  with  $|X| \leq k$  and

$$f(X) \geq (1 - e^{-\gamma}) \cdot OPT.$$

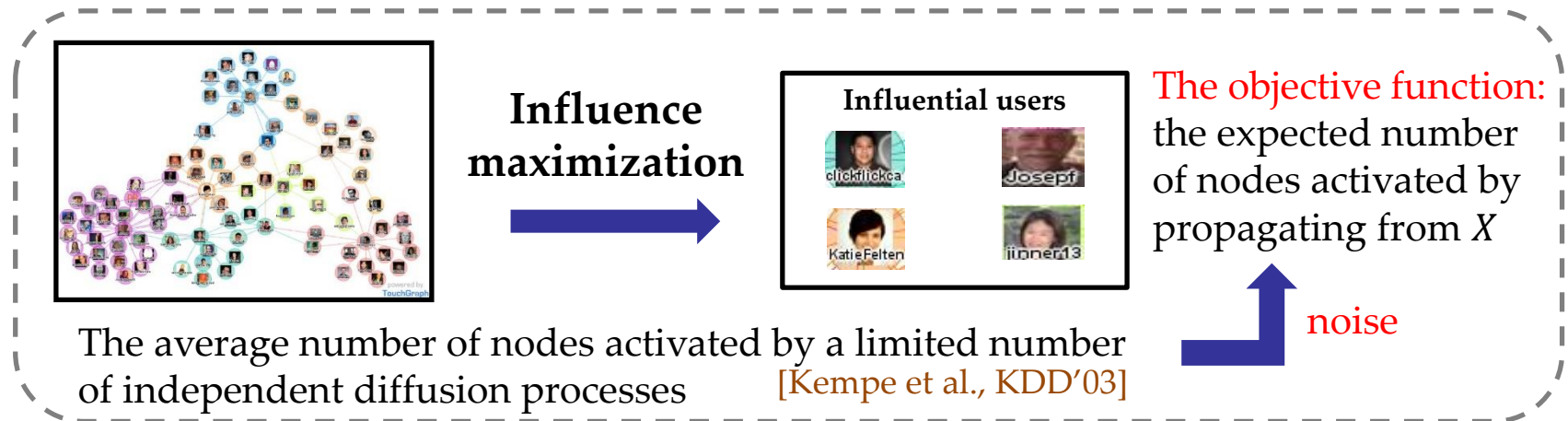


the optimal polynomial-time approximation guarantee  
for monotone  $f$  [Harshaw et al., ICML'19]

# Noise

Previous analyses often assume that the **exact** value of the objective function can be accessed

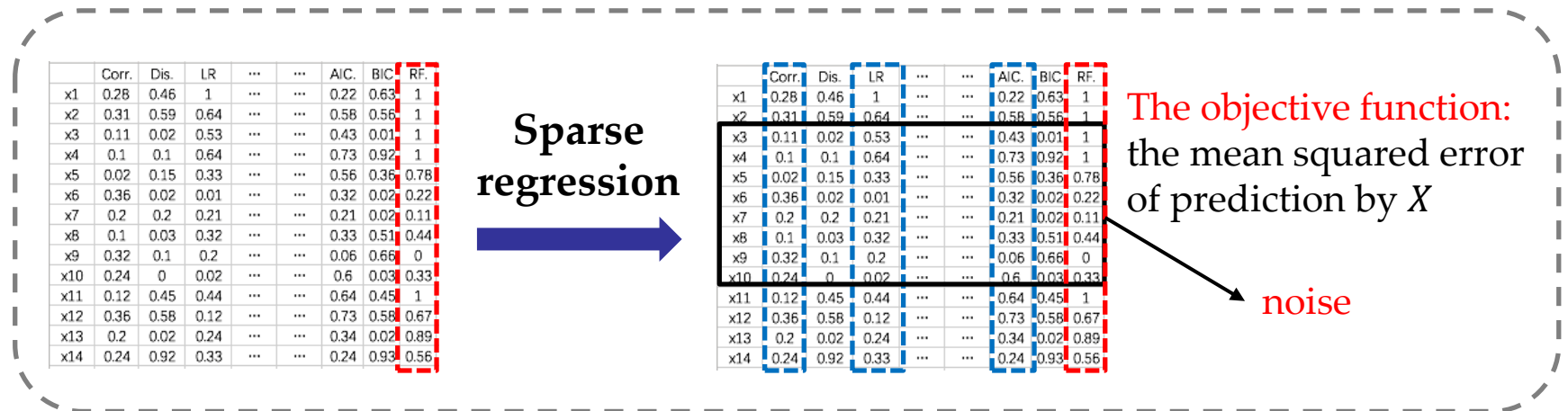
However, in many applications of subset selection, only a **noisy** value of the objective function can be obtained



# Noise

Previous analyses often assume that the **exact** value of the objective function can be accessed

However, in many applications of subset selection, only a **noisy** value of the objective function can be obtained



**How about the performance for noisy subset selection?**

# Noisy subset selection

**Subset selection:** given  $V = \{v_1, \dots, v_n\}$ , an objective function  $f: 2^V \rightarrow \mathbb{R}$  and a budget  $k$ , to find a subset  $X \subseteq V$  such that

$$\max_{X \subseteq V} f(X) \quad \text{s.t.} \quad |X| \leq k$$

exact objective value

noisy objective value

Noise

Multiplicative:  $(1 - \epsilon) \cdot f(X) \leq F(X) \leq (1 + \epsilon) \cdot f(X)$

Additive:  $f(X) - \epsilon \leq F(X) \leq f(X) + \epsilon$

**Applications:** influence maximization, sparse regression

maximizing information gain in graphical models [Chen et al., COLT'15]

crowdsourced image collection summarization [Singla et al., AAAI'16]

# Theoretical analysis

---

Approximation guarantee of POSS

Under multiplicative noise:

$$f(X) \geq \frac{1}{1 + \frac{2\epsilon k}{(1-\epsilon)\gamma}} \left( 1 - \left( \frac{1-\epsilon}{1+\epsilon} \right)^k \left( 1 - \frac{\gamma}{k} \right)^k \right) \cdot OPT$$

noise strength

$\epsilon \leq 1/k$  for a constant  
approximation ratio

Under additive noise:

$$f(X) \geq \left( 1 - \left( 1 - \frac{\gamma}{k} \right)^k \right) \cdot OPT - \left( \frac{2k}{\gamma} - \frac{2k}{\gamma} e^{-\gamma} \right) \epsilon$$

Noiseless approximation guarantee [Qian, Yu and Zhou, NIPS'15]

$$f(X) \geq \left( 1 - \left( 1 - \frac{\gamma}{k} \right)^k \right) \cdot OPT \geq (1 - e^{-\gamma}) \cdot OPT$$

a constant  
approximation ratio

**The performance degrades largely in noisy environments**



# PONSS algorithm

---

**Threshold selection** has theoretically been shown to be robust against noise [Qian, Yu and Zhou, ECJ'18]

$$f(X) \geq f(Y) \longrightarrow f(X) \geq f(Y) + \theta$$

POSS

“better”

$$X \preceq Y \Leftrightarrow \begin{cases} f(X) \geq f(Y) \\ |X| \leq |Y| \end{cases}$$

PONSS [Qian et al., NIPS'17]

Multiplicative:

$$X \preceq Y \Leftrightarrow \begin{cases} f(X) \geq \frac{1+\theta}{1-\theta} f(Y) \\ |X| \leq |Y| \end{cases}$$

Additive:

$$X \preceq Y \Leftrightarrow \begin{cases} f(X) \geq f(Y) + 2\theta \\ |X| \leq |Y| \end{cases}$$

Conservative  
comparison

# Theoretical analysis

---

Under multiplicative noise:

**PONSS**  $f(X) \geq \frac{1-\epsilon}{1+\epsilon} \left( 1 - \left( 1 - \frac{\gamma}{k} \right)^k \right) \cdot OPT$  **Significantly better**

**POSS**  $f(X) \geq \frac{1}{1 + \frac{2\epsilon k}{(1-\epsilon)\gamma}} \left( 1 - \left( \frac{1-\epsilon}{1+\epsilon} \right)^k \left( 1 - \frac{\gamma}{k} \right)^k \right) \cdot OPT$

When  $\gamma = 1$  (submodular),  $\epsilon$  is a constant

**PONSS**      a constant approximation ratio

**POSS**       $\Theta(1/k)$  approximation ratio

# Theoretical analysis

---

Under multiplicative noise:

**PONSS**  $f(X) \geq \frac{1-\epsilon}{1+\epsilon} \left(1 - \left(1 - \frac{\gamma}{k}\right)^k\right) \cdot OPT$  **Significantly better**

**POSS**  $f(X) \geq \frac{1}{1 + \frac{2\epsilon k}{(1-\epsilon)\gamma}} \left(1 - \left(\frac{1-\epsilon}{1+\epsilon}\right)^k \left(1 - \frac{\gamma}{k}\right)^k\right) \cdot OPT$

Under additive noise:

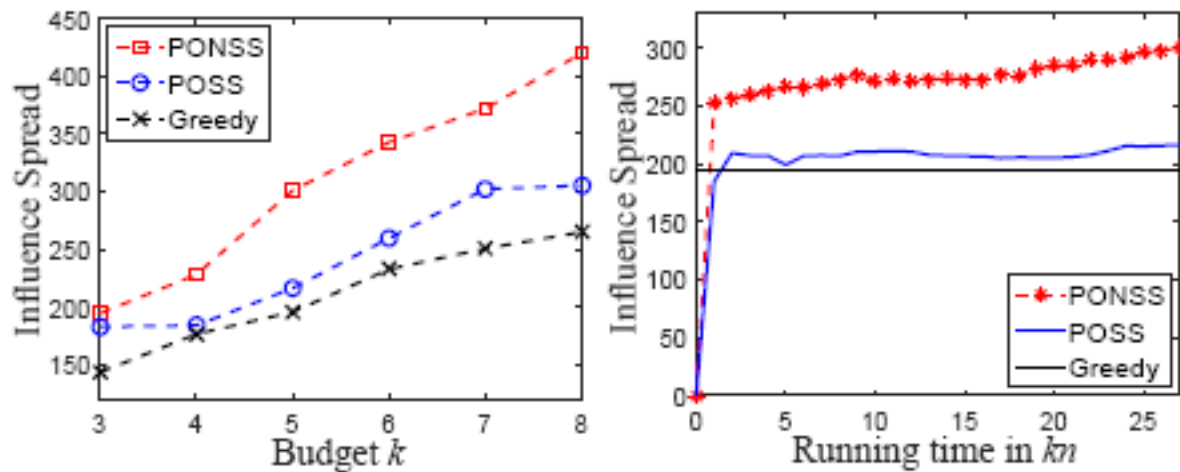
**PONSS**  $f(X) \geq \left(1 - \left(1 - \frac{\gamma}{k}\right)^k\right) \cdot OPT - 2\epsilon$  **better**

**POSS**  $f(X) \geq \left(1 - \left(1 - \frac{\gamma}{k}\right)^k\right) \cdot OPT - \left(\frac{2k}{\gamma} - \frac{2k}{\gamma} e^{-\gamma}\right) \epsilon$

# Experimental results - influence maximization

PONSS (red line) vs. POSS (blue line) vs. Greedy (black line):

- Noisy evaluation: the average of 10 independent Monte Carlo simulations
- The output solution: the average of 10,000 independent Monte Carlo simulations

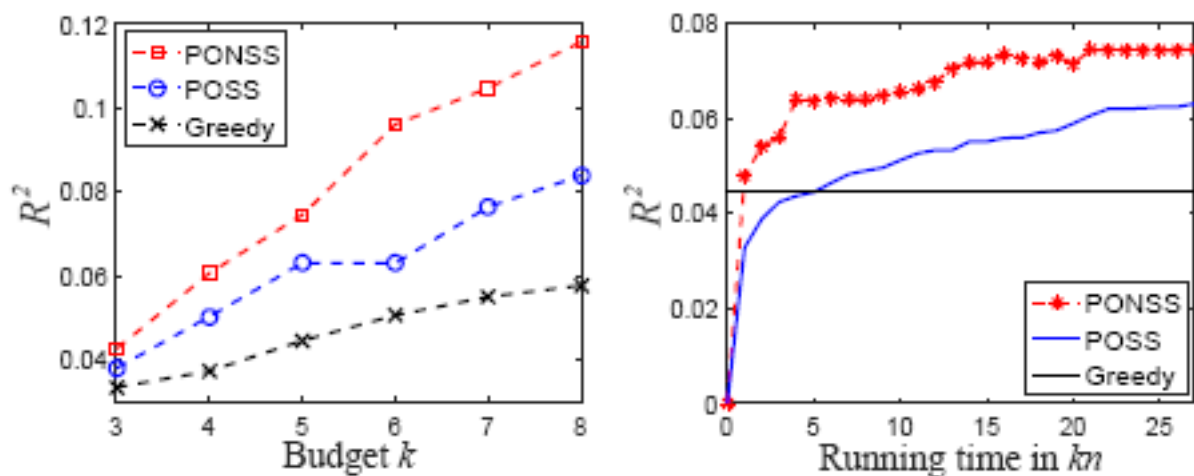


(b) Weibo (10,000 #nodes, 162,371 #edges)

# Experimental results - sparse regression

PONSS (red line) vs. POSS (blue line) vs. Greedy (black line):

- Noisy evaluation: a random sample of 1,000 instances
- The output solution: the whole data set



(a) *protein* (24,387 #inst, 357 #feat)

# Conclusion

---

Safe evolutionary optimization?



Develop running time analysis tools for EAs



Disclose theoretical properties of EAs  
for constrained and noisy optimization



Propose EAs with good polynomial-time  
approximation guarantees for (noisy) subset selection

Yes

A large, light gray curved arrow on the left side of the diagram, starting from the bottom step and pointing upwards towards the top step, indicating a feedback loop.

---

## Collaborators:

### Nanjing University:



Chao Bian



Yang Yu



Zhi-Hua Zhou

### SUSTech:



Ke Tang



Xin Yao

# For details

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- Y. Yu, C. Qian, Z.-H. Zhou. Switch analysis for running time analysis of evolutionary algorithms. *IEEE Transactions on Evolutionary Computation*, 2015, 19(6): 777-792.
- C. Qian, Y. Yu, Z.-H. Zhou. On constrained Boolean Pareto optimization. In: *Proceedings of the 24th International Joint Conference on Artificial Intelligence (IJCAI'15)*, Buenos Aires, Argentina, 2015.
- C. Qian, Y. Yu, Z.-H. Zhou. Subset selection by Pareto optimization. In: *Advances in Neural Information Processing Systems 28 (NIPS'15)*, Montreal, Canada, 2015.
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- C. Qian, Y. Yu, Z.-H. Zhou. Analyzing evolutionary optimization in noisy environments. *Evolutionary Computation*, 2018, 26(1): 1-41.
- C. Qian, Y. Yu, K. Tang, X. Yao, Z.-H. Zhou. Maximizing submodular or monotone approximately submodular functions by multi-objective evolutionary algorithms. *Artificial Intelligence*, 2019, 275: 279-294.

Codes available at <http://www.lamda.nju.edu.cn/qianc/>

# THANK YOU !