



Lecture 5. Online Convex Optimization

Advanced Optimization (Fall 2023)

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Outline

- Online Learning
- Online Convex Optimization
- Convex Functions
- Strongly Convex Functions
- Exp-concave Functions

A Brief Review of Statistical Learning

The fundamental goal of (supervised) learning: *Risk Minimization (RM)*,

$$\min_{h \in \mathcal{H}} \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}} [\ell(h(\mathbf{x}), y)],$$

where

- h denotes the hypothesis (model) from the hypothesis space \mathcal{H} .
- (\mathbf{x}, y) is an instance chosen from a unknown distribution \mathcal{D} .
- $\ell(h(\mathbf{x}), y)$ denotes the loss of using hypothesis h on the instance (\mathbf{x}, y) .

A Brief Review of Statistical Learning

Given a loss function f and distribution \mathcal{D} , the *expected risk* of predictor h is

$$R(h) = \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}}[\ell(h(\mathbf{x}), y)].$$

In practice, we can only access to a sample set $S = \{(\mathbf{x}_1, y_1), \dots, (\mathbf{x}_m, y_m)\}$.

Thus, the following *empirical risk* is naturally defined:

$$\hat{R}_S(h) = \frac{1}{m} \sum_{i=1}^m \ell(h(\mathbf{x}_i), y_i).$$

A Brief Review of Statistical Learning

- A successful story : characterization of sample complexity

- excess risk bound

$$R(\hat{h}) - \inf_{h \in \mathcal{H}} R(h) \leq \mathcal{O} \left(\frac{1}{\sqrt{m}} \right).$$

- generalization error bound

$$\hat{R}_S(\hat{h}) - R(\hat{h}) \leq \mathcal{O} \left(\frac{1}{\sqrt{m}} \right).$$

Offline Towards Online Learning

- Traditional statistical machine learning
 - The training data are available *offline*
 - Learning model is trained based on the offline data in a *batch* setting
- Online learning scenario
 - In real applications, data are in the form of *stream*
 - New data are being collected all the time: after observing a new data point, the model should be *online updated* at a constant cost



A Formulation of Online Learning

- We introduce a *game-theoretic view* to model online learning.
- Online learning is formulated as a repeated game between
 - *Player*: essentially the learner, or you can think as the “learning model”
 - *Environments*: an abstraction of all factors evaluating the model.

At each round $t = 1, 2, \dots$

- (1) the player first picks a model $\mathbf{w}_t \in \mathcal{W}$;
- (2) and simultaneously environments pick an online function $f_t : \mathcal{W} \rightarrow \mathbb{R}$;
- (3) the player suffers loss $f_t(\mathbf{w}_t)$, observes some information about f_t and updates the model.

Online Learning: Formulation

At each round $t = 1, 2, \dots$

- (1) the player first picks a model $\mathbf{w}_t \in \mathcal{W}$;
- (2) and simultaneously environments pick an online function $f_t : \mathcal{W} \rightarrow \mathbb{R}$;
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- An example of online function $f_t : \mathcal{W} \mapsto \mathbb{R}$.

Considering the task of *online classification*, we have

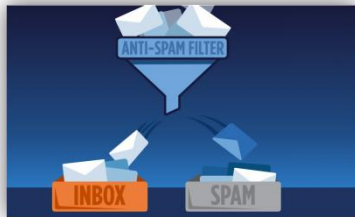
- (i) the loss $\ell : \hat{\mathcal{Y}} \times \mathcal{Y} \mapsto \mathbb{R}$, and
 - (ii) the hypothesis function $h : \mathcal{W} \times \mathcal{X} \mapsto \hat{\mathcal{Y}}$.
- $\implies f_t(\mathbf{w}) = \ell(h(\mathbf{w}; \mathbf{x}_t), y_t)$
 $= \ell(\mathbf{w}^\top \mathbf{x}_t, y_t)$ *for simplicity*

Online Learning: Formulation

At each round $t = 1, 2, \dots$

- (1) the player first picks a model $\mathbf{w}_t \in \mathcal{W}$;
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Spam filtering



(1) Player submits a spam classifier \mathbf{w}_t

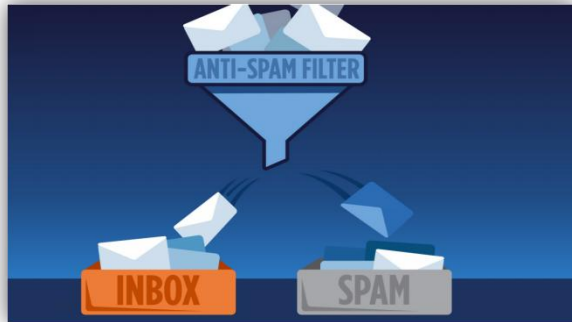


(2) A mail is revealed whether it is a spam



(3) Player suffers loss $f_t(\mathbf{w}_t)$ and updates model

Applications



spam detection (online classification/regression): At each time $t = 1, 2, \dots$

- receive an email $\mathbf{x}_t \in \mathbb{R}^d$;
- predict whether it is a spam $\hat{y}_t \in \{-1, +1\}$;
- see its true label $y_t \in \{-1, +1\}$.



aggregating weather prediction (the expert problem): At each day $t = 1, 2, \dots$

- obtain temperature predictions from N models;
- make the final prediction by randomly following a model according to the probability $p_t \in \Delta_N$;
- on the next day observe the loss of each model $f_t \in [0, 1]^N$.

Performance Measure

- Recall in the statistical learning: *risk*

$$R(h) \triangleq \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}} [\ell(h(\mathbf{x}), y)].$$

- In online learning: *sequential risk*

$$\hat{R}(\{\mathbf{w}_t\}_{t=1}^T) \triangleq \sum_{t=1}^T f_t(\mathbf{w}_t) = \sum_{t=1}^T \ell(\mathbf{w}_t^\top \mathbf{x}_t, y_t).$$

meaning: cumulative loss of models trained on growing data stream $S_t = \{(\mathbf{x}_1, y_1), \dots, (\mathbf{x}_t, y_t)\}$.

Performance Measure

- In offline learning, we use *excess risk* as measure for \hat{h} :

$$\begin{aligned} R(\hat{h}) - \min_{h \in \mathcal{H}} R(h) \\ = \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}}[\ell(\mathbf{w}^\top \mathbf{x}, y)] - \min_{\mathbf{w} \in \mathcal{W}} \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}}[\ell(\mathbf{w}^\top \mathbf{x}, y)] \end{aligned}$$

simply using a linear model \mathbf{w} to parametrize the hypothesis h

- In online learning, we define *regret* as measure for sequence $\{\mathbf{w}_t\}_{t=1}^T$:

$$\begin{aligned} R(\{\mathbf{w}_t\}_{t=1}^T) - \min_{\mathbf{w} \in \mathcal{W}} R(\{\mathbf{w}\}_{t=1}^T) \\ = \sum_{t=1}^T \ell(\mathbf{w}_t^\top \mathbf{x}_t, y_t) - \min_{\mathbf{w} \in \mathcal{W}} \sum_{t=1}^T \ell(\mathbf{w}^\top \mathbf{x}_t, y_t) \end{aligned}$$

Performance Measure

- In offline learning, we use *excess risk* as measure for \hat{h} :

$$R(\hat{h}) - \min_{h \in \mathcal{H}} R(h)$$

$$= \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}}[\ell(\mathbf{w}^\top \mathbf{x}, y)] - \min_{\mathbf{w} \in \mathcal{W}} \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}}[\ell(\mathbf{w}^\top \mathbf{x}, y)]$$

simply using a linear model \mathbf{w} to parametrize the hypothesis h

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$$= \sum_{t=1}^T \ell(\mathbf{w}_t^\top \mathbf{x}_t, y_t) - \min_{\mathbf{w} \in \mathcal{W}} \sum_{t=1}^T \ell(\mathbf{w}^\top \mathbf{x}_t, y_t)$$

$$\text{Regret}_T = \sum_{t=1}^T f_t(\mathbf{w}_t) - \min_{\mathbf{w} \in \mathcal{W}} \sum_{t=1}^T f_t(\mathbf{w})$$

benchmark performance with the offline model (optimal in hindsight)

Another View of Regret

- Ultimate goal: minimize the *cumulative loss* $\sum_{t=1}^T f_t(\mathbf{w}_t)$
- The cumulative loss highly depends on the loss function itself, so we need a benchmark:

$$\text{Regret}_T = \sum_{t=1}^T f_t(\mathbf{w}_t) - \min_{\mathbf{w} \in \mathcal{W}} \sum_{t=1}^T f_t(\mathbf{w})$$

- We hope the regret be sub-linear dependence with T .

$$\frac{\text{Regret}_T}{T} \rightarrow 0 \text{ as } T \rightarrow \infty$$

Hannan Consistency

ALT'16

Hannan Consistency in On-Line Learning
in Case of Unbounded Losses Under Partial
Monitoring^{*,**}

Chamy Allenberg¹, Peter Auer², László Györfi³, and György Ottucsák³

Compared with Statistical Offline Learning

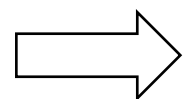
- Memory efficient
- Do not need i.i.d. assumption
 - the environment can be even adversarial
 - typically, the regret analysis does not need concentration
- Strictly harder than statistical learning
 - under non-i.i.d. assumption
 - online to batch conversion

Online Learning: Is it (provably) solvable?

- In general, the online learning formulation is *hard* to solve.

At each round $t = 1, 2, \dots$

- (1) the player first picks a model $\mathbf{w}_t \in \mathcal{W}$;
- (2) and simultaneously environments pick an online function $f_t : \mathcal{W} \rightarrow \mathbb{R}$;
- (3) the player suffers loss $f_t(\mathbf{w}_t)$, observes some information about f_t and updates the model.



A Trackable Case: *Online Convex Optimization*

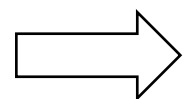
*requiring feasible domain and online functions to be **convex***

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A Trackable Case: *Online Convex Optimization*

*requiring feasible domain and online functions to be **convex***

Online Convex Optimization

- OCO framework
 - feasible domain is a convex set
 - online functions are convex

At each round $t = 1, 2, \dots$

- (1) the player first picks a model \mathbf{x}_t from a convex set $\mathcal{X} \subseteq \mathbb{R}^d$;
- (2) and environments pick an online convex function $f_t : \mathcal{X} \rightarrow \mathbb{R}$;
- (3) the player suffers loss $f_t(\mathbf{x}_t)$, observes some information about f_t and updates the model.

From this point forward, we use \mathbf{x} (and \mathcal{X}) instead of \mathbf{w} (and \mathcal{W}) for consistency with opt. language.

OCO: Different Setups

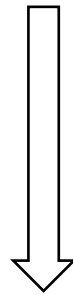
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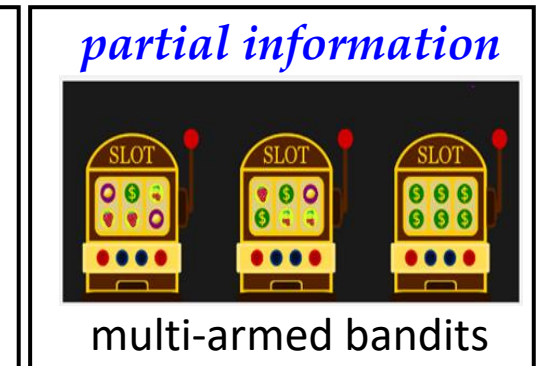
on the feedback information:

- *full information*: observe entire f_t (or at least gradient $\nabla f_t(\mathbf{x}_t)$)

- *partial information (bandits)*: observe function value $f_t(\mathbf{x}_t)$ only



less information



OCO: Different Setups

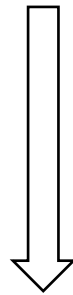
At each round $t = 1, 2, \dots$

- (1) the player first picks a model \mathbf{x}_t from a convex set $\mathcal{X} \subseteq \mathbb{R}^d$;
- (2) and **environments** pick an online convex function $f_t : \mathcal{X} \rightarrow \mathbb{R}$;
- (3) the player suffers loss $f_t(\mathbf{x}_t)$, observes some information about f_t and updates the model.

on the difficulty of environments:

- stochastic setting

- adversarial setting { oblivious
adaptive
(non-oblivious)



*less restricted
but harder*

oblivious adversary



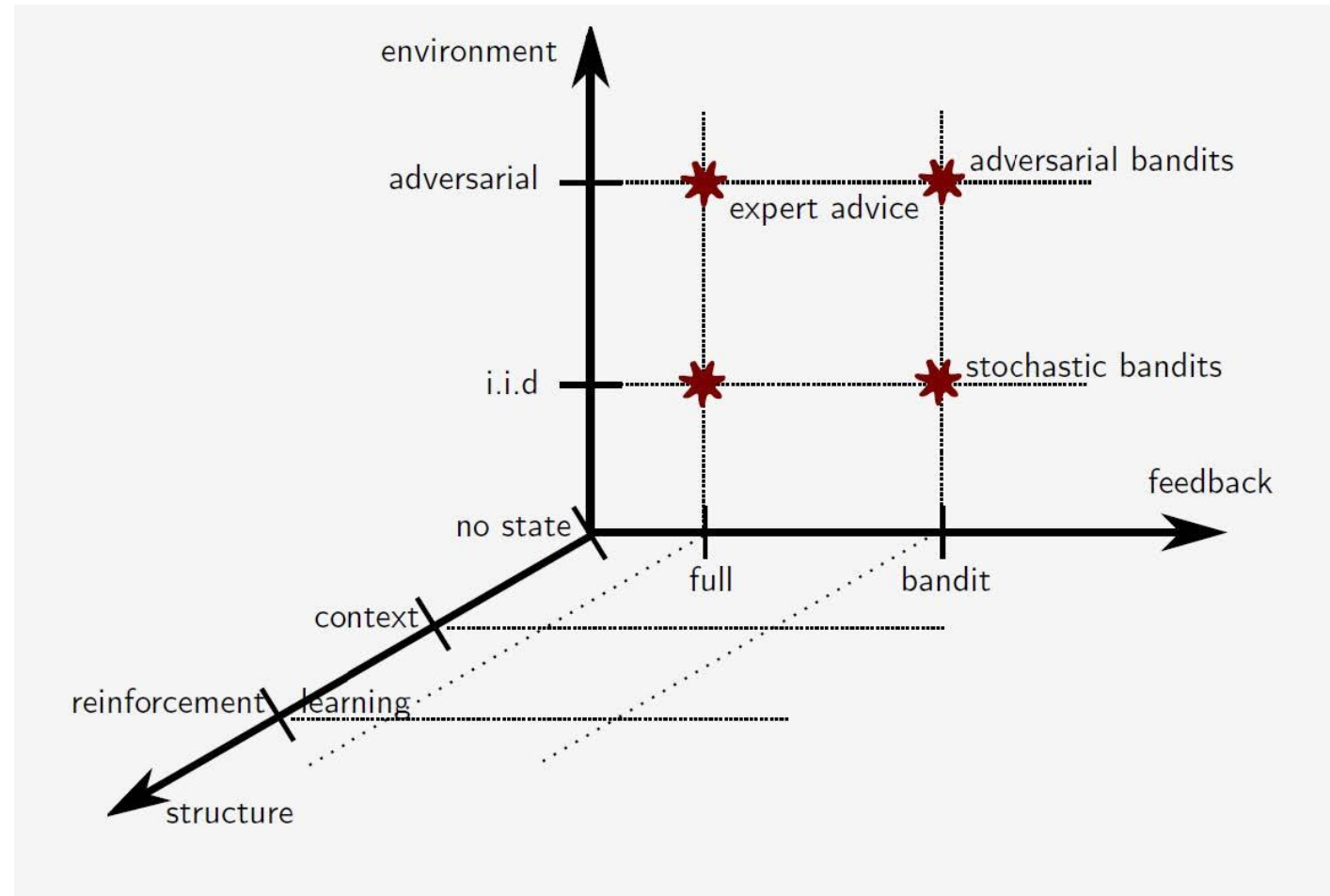
examination

adaptive adversary



interview

The Space of Online Learning Problems



Yevgeny Seldin. The Space of Online Learning Problems, ECML-PKDD, Porto, Portugal, 2015.

Online Learning

- Full-information setting:
 - Online Convex Optimization
 - Prediction with Expert Advice
 - ...

- Partial-information setting:
 - Multi-Armed Bandits
 - Linear Bandits/Parametric Bandits
 - Bandit Convex Optimization
 - ...

Online Learning

- **Full-information setting:**
 - Online Convex Optimization
 - Prediction with Expert Advice
 - ...

- Partial-information setting:
 - Multi-Armed Bandits
 - Linear Bandits/Parametric Bandits
 - Bandit Convex Optimization
 - ...

History: Two-Player Zero-Sum Games

Theory of repeated games



James Hannan
(1922–2010)



David Blackwell
(1919–2010)

Learning to play a game (1956)

Play a game repeatedly against a possibly suboptimal opponent

Zero-sum 2-person games played more than once

	1	2	...	M
1	$\ell(1,1)$	$\ell(1,2)$...	
2	$\ell(2,1)$	$\ell(2,2)$...	
\vdots	\vdots	\vdots	\ddots	
N				

$N \times M$ known loss matrix

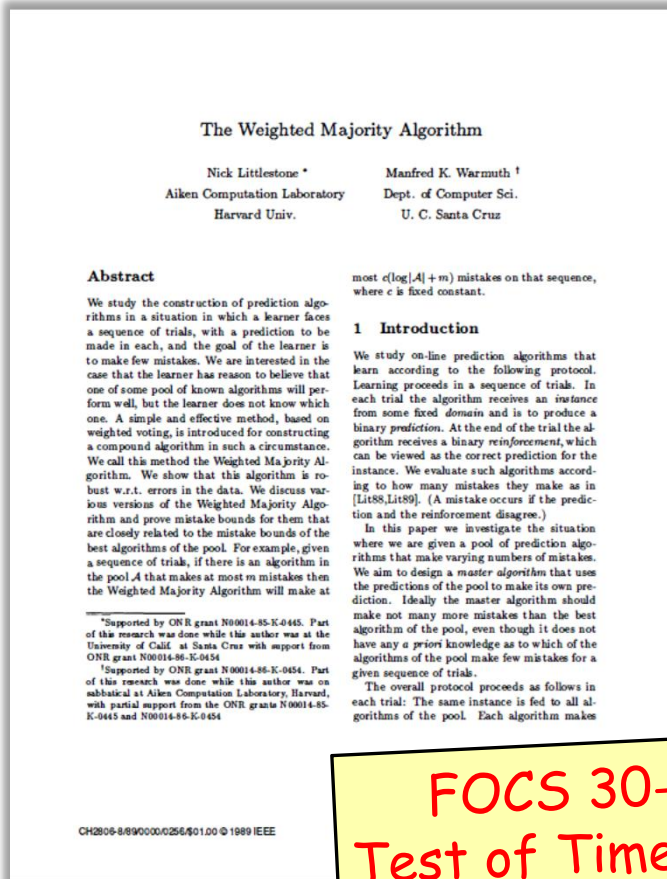
- Row player (**player**) has N actions
- Column player (**opponent**) has M actions

For each game round $t = 1, 2, \dots$

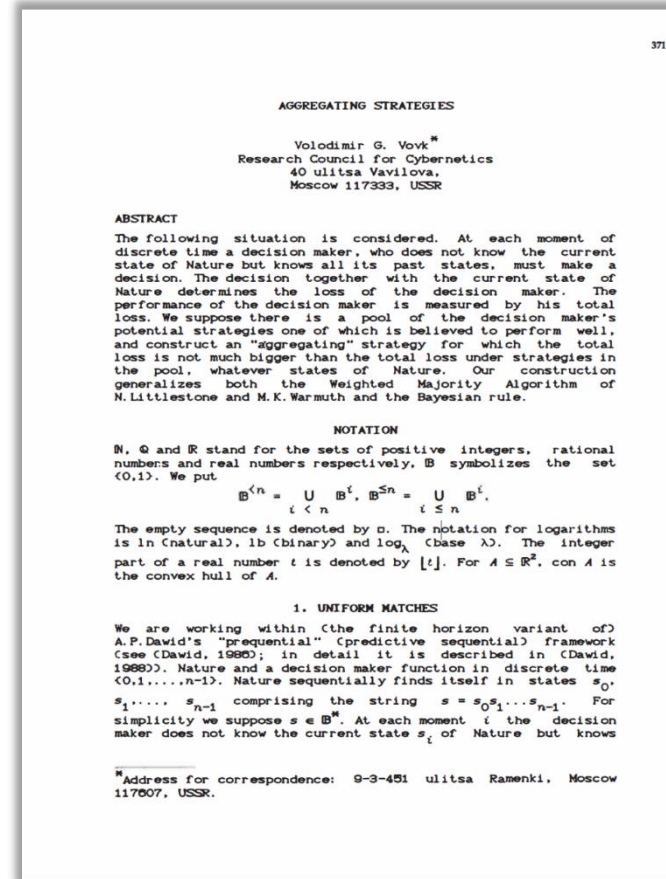
- Player chooses action i_t and opponent chooses action y_t
- The player suffers loss $\ell(i_t, y_t)$ (= gain of opponent)

Player can learn from opponent's history of past choices y_1, \dots, y_{t-1}

History: Prediction with Expert Advice



Nick Littlestone and Manfred K. Warmuth.
"The Weighted Majority Algorithm." FOCS 1989: 256-261.



Volodimir G. Vovk
Royal Holloway,
University of London

Volodimir G. Vovk. "Aggregating Strategies." COLT 1990: 371-383.

Online Convex Optimization

- Convex Functions
- Strongly Convex Functions
- Exponentially Concave Functions

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Online Optimization with Convex Functions

Definition 2 (Convex Function). A function $f : \mathcal{X} \mapsto \mathbb{R}$ is convex if for any $\mathbf{x}, \mathbf{y} \in \mathcal{X}$

$$\forall \alpha \in [0, 1], f((1 - \alpha)\mathbf{x} + \alpha\mathbf{y}) \leq (1 - \alpha)f(\mathbf{x}) + \alpha f(\mathbf{y}).$$

Equivalently, if f is differentiable, we have that $\forall \mathbf{x}, \mathbf{y} \in \mathcal{X}$,

$$f(\mathbf{y}) \geq f(\mathbf{x}) + \nabla f(\mathbf{x})^\top (\mathbf{y} - \mathbf{x}).$$

Online Gradient Descent

Online Gradient Descent (OGD)

At each round $t = 1, 2, \dots$

1. the player first picks a model $\mathbf{x}_t \in \mathcal{X}$;
2. and simultaneously environments pick a *convex loss function* $f_t : \mathcal{X} \rightarrow \mathbb{R}$;
3. the player suffers loss $f_t(\mathbf{x}_t)$, observes the information (loss) f_t and update the model according to $\mathbf{x}_{t+1} = \Pi_{\mathcal{X}}[\mathbf{x}_t - \eta_t \nabla f_t(\mathbf{x}_t)]$.

- $\Pi_{\mathcal{X}}[\mathbf{y}] = \arg \min_{\mathbf{x} \in \mathcal{X}} \|\mathbf{x} - \mathbf{y}\|_2$ denotes the Euclidean projection onto the feasible set \mathcal{X} .
- This belongs to the full-information setting, so player can access the gradient $\nabla f_t(\mathbf{x}_t)$.

But actually the gradient is the only required, so it's also called *gradient-feedback* OCO model.

OGD: Regret Analysis

- The following assumptions are required for standard analysis.

Assumption 1 (Convexity). The feasible set \mathcal{X} is closed and convex in Euclidean space, and f_1, \dots, f_T are convex functions.

Assumption 2 (Bounded Decision Set). The diameter of the set \mathcal{X} is upper bounded by D , i.e., $\forall \mathbf{x}, \mathbf{y} \in \mathcal{X}, \|\mathbf{x} - \mathbf{y}\| \leq D$.

Assumption 3 (Bounded Gradient). The norm of the subgradients of f is upper bounded by G , i.e., $\|\nabla f(\mathbf{x})\| \leq G$ for all $\mathbf{x} \in \mathcal{X}$.

OGD: Regret Analysis

Theorem 3 (Regret bound for OGD). *Under Assumptions 1, 2 and 3, online gradient descent (OGD) with step sizes $\eta_t = \frac{D}{G\sqrt{t}}$ for $t \in [T]$ guarantees:*

$$\text{Regret}_T = \sum_{t=1}^T f_t(\mathbf{x}_t) - \min_{\mathbf{x} \in \mathcal{X}} \sum_{t=1}^T f_t(\mathbf{x}) \leq \frac{3}{2}GD\sqrt{T} = \mathcal{O}(\sqrt{T}).$$

The First Gradient Descent Lemma

Lemma 1. *Suppose that f is proper, closed and convex; the feasible domain \mathcal{X} is nonempty, closed and convex. Let $\{\mathbf{x}_t\}_{t=1}^T$ be the sequence generated by the gradient descent method. Then for any $\mathbf{u} \in \mathcal{X}^*$ and $t \geq 0$,*

$$\|\mathbf{x}_{t+1} - \mathbf{u}\|^2 \leq \|\mathbf{x}_t - \mathbf{u}\|^2 - 2\eta_t(f_t(\mathbf{x}_t) - f_t(\mathbf{u})) + \eta_t^2 \|\nabla f_t(\mathbf{x}_t)\|^2.$$

Proof:

$$\begin{aligned} \|\mathbf{x}_{t+1} - \mathbf{u}\|^2 &= \|\Pi_{\mathcal{X}}[\mathbf{x}_t - \eta_t \nabla f_t(\mathbf{x}_t)] - \mathbf{u}\|^2 \quad (\text{GD}) \\ &\leq \|\mathbf{x}_t - \eta_t \nabla f_t(\mathbf{x}_t) - \mathbf{u}\|^2 \quad (\text{Pythagoras Theorem}) \\ &= \|\mathbf{x}_t - \mathbf{u}\|^2 - 2\eta_t \langle \nabla f_t(\mathbf{x}_t), \mathbf{x}_t - \mathbf{u} \rangle + \eta_t^2 \|\nabla f_t(\mathbf{x}_t)\|^2 \\ &\leq \|\mathbf{x}_t - \mathbf{u}\|^2 - 2\eta_t(f_t(\mathbf{x}_t) - f_t(\mathbf{u})) + \eta_t^2 \|\nabla f_t(\mathbf{x}_t)\|^2 \\ &\quad (\text{convexity: } f_t(\mathbf{x}_t) - f_t(\mathbf{u}) = f_t(\mathbf{x}_t) - f_t(\mathbf{u}) \leq \langle \nabla f_t(\mathbf{x}_t), \mathbf{x}_t - \mathbf{u} \rangle) \quad \square \end{aligned}$$

Proof for OGD Regret Bound

Proof: We use the first gradient descent lemma to analyze online gradient descent.

Lemma 1. *Suppose that f is proper, closed and convex; the feasible domain \mathcal{X} is nonempty, closed and convex. Let $\{\mathbf{x}_t\}_{t=1}^T$ be the sequence generated by the gradient descent method. Then for any $\mathbf{u} \in \mathcal{X}^*$ and $t \geq 0$,*

$$\|\mathbf{x}_{t+1} - \mathbf{u}\|^2 \leq \|\mathbf{x}_t - \mathbf{u}\|^2 - 2\eta_t(f_t(\mathbf{x}_t) - f_t(\mathbf{u})) + \eta_t^2 \|\nabla f_t(\mathbf{x}_t)\|^2.$$

By Lemma 1,

$$2(f_t(\mathbf{x}_t) - f_t(\mathbf{u})) \leq \frac{\|\mathbf{x}_t - \mathbf{u}\|^2 - \|\mathbf{x}_{t+1} - \mathbf{u}\|^2}{\eta_t} + \eta_t G^2$$

Proof for OGD Regret Bound

Proof: By setting $\eta_t = \frac{D}{G\sqrt{t}}$ (with $\frac{1}{\eta_0} := 0$), summing over T :

$$\begin{aligned} 2 \left(\sum_{t=1}^T f_t(\mathbf{x}_t) - f_t(\mathbf{u}) \right) &\leq \sum_{t=1}^T \frac{\|\mathbf{x}_t - \mathbf{u}\|^2 - \|\mathbf{x}_{t+1} - \mathbf{u}\|^2}{\eta_t} + G^2 \sum_{t=1}^T \eta_t && \text{(GD lemma)} \\ &\leq \sum_{t=1}^T \|\mathbf{x}_t - \mathbf{u}\|^2 \left(\frac{1}{\eta_t} - \frac{1}{\eta_{t-1}} \right) + G^2 \sum_{t=1}^T \eta_t && (\|\mathbf{x}_{T+1} - \mathbf{u}\|^2 \geq 0) \\ &\leq D^2 \sum_{t=1}^T \left(\frac{1}{\eta_t} - \frac{1}{\eta_{t-1}} \right) + G^2 \sum_{t=1}^T \eta_t \\ &\leq D^2 \frac{1}{\eta_T} + G^2 \sum_{t=1}^T \eta_t && (\eta_t = \frac{D}{G\sqrt{t}} \text{ and } \sum_{t=1}^T \frac{1}{\sqrt{t}} \leq 2\sqrt{T}) \\ &\leq 3DG\sqrt{T}. && \square \end{aligned}$$

Online Convex Optimization

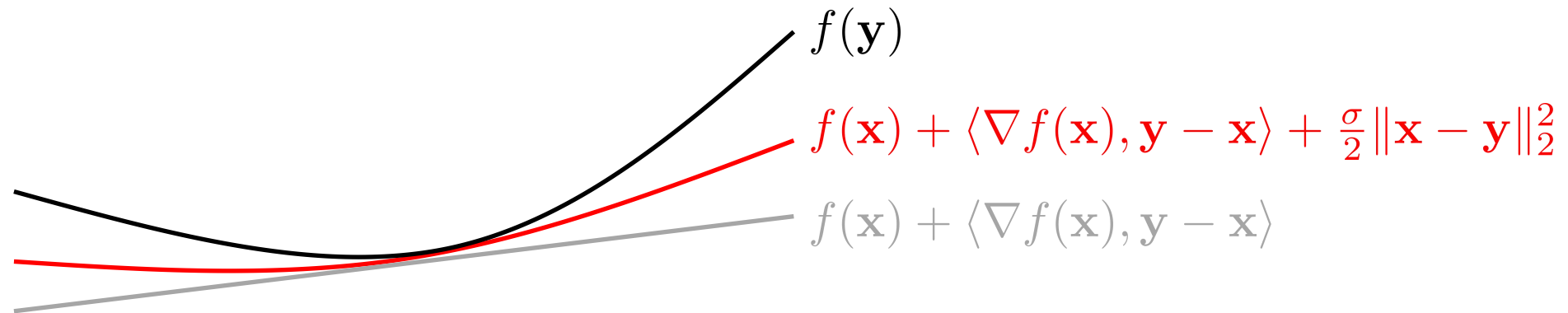
- Convex Functions
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OCO with Strongly Convex Functions

Definition 3 (Strong Convexity). A function f is σ -strongly convex if, for any $\mathbf{x}, \mathbf{y} \in \text{dom } f$,

$$f(\mathbf{y}) \geq f(\mathbf{x}) + \nabla f(\mathbf{x})^\top (\mathbf{y} - \mathbf{x}) + \frac{\sigma}{2} \|\mathbf{y} - \mathbf{x}\|^2,$$

or equivalently, $\nabla^2 f(\mathbf{x}) \succeq \alpha I$.



OGD for Strongly Convex Loss

Online Gradient Descent (OGD)

At each round $t = 1, 2, \dots$

- (1) the player first picks a model $\mathbf{x}_t \in \mathcal{X}$;
- (2) and simultaneously environments pick a **strongly convex loss** $f_t : \mathcal{X} \rightarrow \mathbb{R}$;
- (3) the player suffers loss $f_t(\mathbf{x}_t)$, observes the information (loss) f_t and update the model according to $\mathbf{x}_{t+1} = \Pi_{\mathcal{X}}[\mathbf{x}_t - \eta_t \nabla f_t(\mathbf{x}_t)]$.

- The learning rate for strongly convex OGD is set as $\eta_t = \frac{1}{\sigma t}$.

OGD for Strongly Convex Loss

Theorem 4 (Regret bound for strongly-convex functions). *Under Assumption 1 and Assumption 3, for σ -strongly convex loss functions, online gradient descent with step sizes $\eta_t = \frac{1}{\sigma t}$ achieves the following guarantee*

$$\text{Regret}_T \leq \frac{G^2}{2\sigma} (1 + \log T) = \mathcal{O}(\log T).$$

- Strongly convex case compared with convex case: $\mathcal{O}(\log T)$ vs. $\mathcal{O}(\sqrt{T})$
- A caveat is that we now don't need Assumption 2 (bounded domain).

OCO with Strongly Convex Functions

Proof: we start by extending *the first GD lemma* to strongly convex case.

Strongly convex case:

$$\begin{aligned}\|\mathbf{x}_{t+1} - \mathbf{u}\|^2 &\leq \|\mathbf{x}_t - \mathbf{u}\|^2 - 2\eta_t \langle \nabla f_t(\mathbf{x}_t), \mathbf{x}_t - \mathbf{u} \rangle + \eta_t^2 \|\nabla f_t(\mathbf{x}_t)\|^2 \\ &\leq \|\mathbf{x}_t - \mathbf{u}\|^2 - 2\eta_t \left(f_t(\mathbf{x}_t) - f_t(\mathbf{u}) + \frac{\sigma}{2} \|\mathbf{x}_t - \mathbf{u}\|^2 \right) + \eta_t^2 \|\nabla f_t(\mathbf{x}_t)\|^2 \\ &\quad \text{(strong convexity: } f_t(\mathbf{x}_t) - f_t(\mathbf{u}) + \frac{\sigma}{2} \|\mathbf{x}_t - \mathbf{u}\|^2 \leq \langle \nabla f_t(\mathbf{x}_t), \mathbf{x}_t - \mathbf{u} \rangle \text{)} \\ &\leq (1 - \sigma\eta_t) \|\mathbf{x}_t - \mathbf{u}\|^2 - 2\eta_t (f_t(\mathbf{x}_t) - f_t(\mathbf{u})) + \eta_t^2 \|\nabla f_t(\mathbf{x}_t)\|^2 \\ \implies f_t(\mathbf{x}_t) - f_t(\mathbf{u}) &\leq \frac{\eta_t^{-1} - \sigma}{2} \|\mathbf{x}_t - \mathbf{u}\|^2 - \frac{\eta_t^{-1}}{2} \|\mathbf{x}_{t+1} - \mathbf{u}\|^2 + \frac{\eta_t G^2}{2} \quad \text{(rearranging)}\end{aligned}$$

OCO with Strongly Convex Functions

Proof:
$$f_t(\mathbf{x}_t) - f_t(\mathbf{u}) \leq \frac{\eta_t^{-1} - \sigma}{2} \|\mathbf{x}_t - \mathbf{u}\|^2 - \frac{\eta_t^{-1}}{2} \|\mathbf{x}_{t+1} - \mathbf{u}\|^2 + \frac{\eta_t G^2}{2}$$

Summing from $t = 1$ to T , setting $\eta_t = \frac{1}{\sigma t}$ (define $\frac{1}{\eta_0} := 0$):

$$\begin{aligned} 2 \sum_{t=1}^T (f_t(\mathbf{x}_t) - f_t(\mathbf{u})) &\leq \sum_{t=1}^T \|\mathbf{x}_t - \mathbf{u}\|^2 \left(\frac{1}{\eta_t} - \frac{1}{\eta_{t-1}} - \sigma \right) + G^2 \sum_{t=1}^T \eta_t \quad \left(\frac{1}{\eta_0} := 0 \right) \\ &= 0 + G^2 \sum_{t=1}^T \frac{1}{\sigma t} \quad \left(\frac{1}{\eta_0} \triangleq 0, \|\mathbf{x}_{T+1} - \mathbf{u}\|^2 \geq 0, \frac{1}{\eta_t} - \frac{1}{\eta_{t-1}} - \sigma = 0 \right) \\ &\leq \frac{G^2}{\sigma} (1 + \log T). \quad \square \end{aligned}$$

Online Convex Optimization

- Convex Functions
- Strongly Convex Functions
- Exponentially Concave Functions

Comparison of (Strongly) Convex Problems

Convex Problem

Property: $f_t(\mathbf{x}) \geq f_t(\mathbf{y}) + \nabla f_t(\mathbf{y})^\top (\mathbf{x} - \mathbf{y})$

$$\text{OGD: } \mathbf{x}_{t+1} = \Pi_{\mathcal{X}} \left[\mathbf{x}_t - \frac{1}{\sqrt{t}} \nabla f_t(\mathbf{x}_t) \right]$$

$$\text{Regret}_T \leq \frac{3}{2} GD\sqrt{T}$$

Strongly Convex Problem

Property: $f_t(\mathbf{y}) \geq f_t(\mathbf{x}) + \nabla f_t(\mathbf{x})^\top (\mathbf{y} - \mathbf{x}) + \frac{\sigma}{2} \|\mathbf{y} - \mathbf{x}\|^2$

$$\text{OGD: } \mathbf{x}_{t+1} = \Pi_{\mathcal{X}} \left[\mathbf{x}_t - \frac{1}{\sigma t} \nabla f_t(\mathbf{x}_t) \right]$$

$$\text{Regret}_T \leq \frac{G^2}{2\sigma} (1 + \log T)$$

Can we explore more function class with a regret rate faster than \sqrt{T} ?

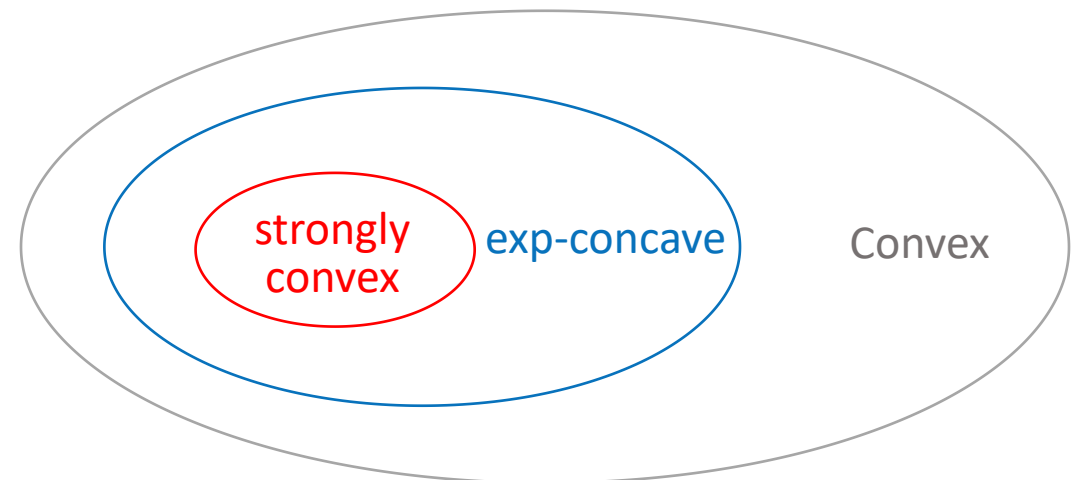
Exponentially-concave Function

Definition 2 (Exp-concavity). A convex function $f : \mathbb{R}^d \mapsto \mathbb{R}$ is defined to be α -exp-concave over $\mathcal{X} \subseteq \mathbb{R}^d$ if the function g is concave, where $g : \mathcal{X} \mapsto \mathbb{R}$ is defined as


$$g(\mathbf{x}) = e^{-\alpha f(\mathbf{x})}.$$

Directly employ OGD: Regret bound $\mathcal{O}(\sqrt{T})$.

But actually we can get a **tighter** bound!



An Example for Exp-concave Learning

- Universal Portfolio Selection 
 - a total of d stocks in the stock market.
 - each round, the player chooses stocks by a distribution $\mathbf{x}_t \in \Delta_d$.
 - the market returns the **price ratio** θ_t between iter t and $t + 1$,

$$\theta_t(i) = \frac{\text{price of stock}_i \text{ at time } t + 1}{\text{price of stock}_i \text{ at time } t}$$

which means that our final wealth W_T will be: $W_T = W_1 \cdot \prod_{t=1}^T \theta_t^\top \mathbf{x}_t$

\Rightarrow Our goal is to **maximize our wealth** at time T .

An Example for Exp-concave Learning

- Universal Portfolio Selection 

- we hope to maximize the logarithm of W_T
- using OCO framework,

$$\log \frac{W_T}{W_1} = \sum_{t=1}^T \log \boldsymbol{\theta}_t^\top \mathbf{x}_t$$

$$f_t(\mathbf{x}) = \log(\boldsymbol{\theta}_t^\top \mathbf{x})$$

At each round $t = 1, 2, \dots$

- (1) the player first picks a model $\mathbf{x}_t \in \Delta_d$;
- (2) and simultaneously environments pick an online function $f_t : \mathcal{X} \rightarrow \mathbb{R}$;
- (3) the player get a **gain** $f_t(\mathbf{x}_t) = \log(\boldsymbol{\theta}_t^\top \mathbf{x}_t)$, observes f_t and updates the model.

- Goal: $\text{Regret}_T = \max_{\mathbf{x}^* \in \Delta_d} \sum_{t=1}^T f_t(\mathbf{x}^*) - \sum_{t=1}^T f_t(\mathbf{x}_t)$

online function is exp-concave

Exponential-concave Function

Lemma 3 (Property of Exp-concavity). *Let $f : \mathcal{X} \rightarrow \mathbb{R}$ be an α -exp-concave function, and D, G denote the diameter of \mathcal{X} and a bound on the (sub)gradients of f respectively. The following holds for all $\gamma \leq \frac{1}{2} \min \left\{ \frac{1}{GD}, \alpha \right\}$ and all $\mathbf{x}, \mathbf{y} \in \mathcal{X}$:*

$$f(\mathbf{x}) \geq f(\mathbf{y}) + \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) + \frac{\gamma}{2} (\mathbf{x} - \mathbf{y})^\top \nabla f(\mathbf{y}) \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}).$$

Proof. Recall that f is α -exp-concave if and only if $e^{-\alpha f(\mathbf{x})}$ is concave.

As $2\gamma \leq \alpha$, $e^{-2\gamma f(\mathbf{x})} = (e^{-\alpha f(\mathbf{x})})^{2\gamma/\alpha}$ is also concave and thus is 2γ -exp-concave.

$$e^{-2\gamma f(\mathbf{x})} - e^{-2\gamma f(\mathbf{y})} \leq \left\langle \mathbf{x} - \mathbf{y}, -2\gamma e^{-2\gamma f(\mathbf{y})} \nabla f(\mathbf{y}) \right\rangle.$$

(concavity)

Exponential-concave Function

Lemma 3 (Property of Exp-concavity). *Let $f : \mathcal{X} \rightarrow \mathbb{R}$ be an α -exp-concave function, and D, G denote the diameter of \mathcal{X} and a bound on the (sub)gradients of f respectively. The following holds for all $\gamma \leq \frac{1}{2} \min \left\{ \frac{1}{GD}, \alpha \right\}$ and all $\mathbf{x}, \mathbf{y} \in \mathcal{X}$:*

$$f(\mathbf{x}) \geq f(\mathbf{y}) + \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) + \frac{\gamma}{2} (\mathbf{x} - \mathbf{y})^\top \nabla f(\mathbf{y}) \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}).$$

Proof. Dividing $e^{-2\gamma f(\mathbf{y})}$ at both sides achieves

$$f(\mathbf{y}) - f(\mathbf{x}) \leq \frac{1}{2\gamma} \log \left(1 + 2\gamma \langle \mathbf{y} - \mathbf{x}, \nabla f(\mathbf{y}) \rangle \right).$$

Our constructive condition $\gamma \leq \frac{1}{2} \min \left\{ \frac{1}{GD}, \alpha \right\}$ ensures $|2\gamma \langle \mathbf{y} - \mathbf{x}, \nabla f(\mathbf{y}) \rangle| \leq 1$,

$$f(\mathbf{y}) - f(\mathbf{x}) \leq \langle \mathbf{y} - \mathbf{x}, \nabla f(\mathbf{y}) \rangle - \frac{\gamma}{2} \langle \mathbf{y} - \mathbf{x}, \nabla f(\mathbf{y}) \rangle^2$$

($\log(1+x) \leq x - \frac{1}{4}x^2$) holds for ($|x| \leq 1$) \square

A Comparison of Different Curvatures

- Convex

$$f(\mathbf{x}) \geq f(\mathbf{y}) + \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y})$$

- Strongly Convex

$$f(\mathbf{x}) \geq f(\mathbf{y}) + \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) + \frac{\sigma}{2} \|\mathbf{x} - \mathbf{y}\|_2^2$$

- Exponentially Concave

$$\begin{aligned} f(\mathbf{x}) &\geq f(\mathbf{y}) + \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) + \frac{\gamma}{2} (\mathbf{x} - \mathbf{y})^\top \nabla f(\mathbf{y}) \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) \\ &= f(\mathbf{y}) + \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) + \frac{\gamma}{2} \|\mathbf{x} - \mathbf{y}\|_{\nabla f(\mathbf{y}) \nabla f(\mathbf{y})^\top}^2 \end{aligned}$$

Exponential-concave Function

Lemma 3 (Property of Exp-concavity). *Let $f : \mathcal{X} \rightarrow \mathbb{R}$ be an α -exp-concave function, and D, G denote the diameter of \mathcal{X} and a bound on the (sub)gradients of f respectively. The following holds for all $\gamma \leq \frac{1}{2} \min \left\{ \frac{1}{GD}, \alpha \right\}$ and all $\mathbf{x}, \mathbf{y} \in \mathcal{X}$:*

$$\begin{aligned} f(\mathbf{x}) &\geq f(\mathbf{y}) + \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) + \frac{\gamma}{2} (\mathbf{x} - \mathbf{y})^\top \nabla f(\mathbf{y}) \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) \\ &= f(\mathbf{y}) + \nabla f(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) + \frac{\gamma}{2} \|\mathbf{x} - \mathbf{y}\|_{\nabla f(\mathbf{y}) \nabla f(\mathbf{y})^\top}^2 \end{aligned}$$

Algorithmic intuition:

- For convex loss, we use 2-norm to encode the structure of the space.
- Can we exploit **local structures** of exp-concave loss to improve the regret?

Intuition

- Convex

$$f_t(\mathbf{x}) \geq f_t(\mathbf{x}_t) + \nabla f_t(\mathbf{x}_t)^\top (\mathbf{x} - \mathbf{x}_t)$$

$$\mathbf{x}_{t+1} = \Pi_{\mathcal{X}} [\mathbf{x}_t - \eta_t \nabla f_t(\mathbf{x}_t)] \quad \text{OGD with } \eta_t = \mathcal{O}(1/\sqrt{t})$$

- Strongly convex

$$f_t(\mathbf{x}) \geq f_t(\mathbf{x}_t) + \nabla f_t(\mathbf{x}_t)^\top (\mathbf{x} - \mathbf{x}_t) + \frac{\sigma}{2} \|\mathbf{x} - \mathbf{x}_t\|_2^2$$

$$\mathbf{x}_{t+1} = \Pi_{\mathcal{X}} [\mathbf{x}_t - \eta_t \nabla f_t(\mathbf{x}_t)] \quad \text{OGD with } \eta_t = \mathcal{O}(1/t)$$

- Exp-concave

$$f_t(\mathbf{x}) \geq f_t(\mathbf{x}_t) + \nabla f_t(\mathbf{x}_t)^\top (\mathbf{x} - \mathbf{x}_t) + \frac{\gamma}{2} \|\mathbf{x} - \mathbf{x}_t\|_{\nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top}^2$$

⇒ We may still GD update, but the step size should be “*data-dependent*”.

Intuitively, step size should be stretched heterogeneously in different directions, being smaller when $\nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top$ is “larger”.

ONS for Exp-concave Function

Online Newton Step

Input: parameters $\gamma, \varepsilon > 0$, matrix $A_0 = \varepsilon I_d$

At each round $t = 1, 2, \dots$

- (1) the player first picks a model $\mathbf{x}_t \in \mathcal{X} \subseteq \mathbb{R}^d$;
- (2) and simultaneously environments pick an *exp-concave loss function* $f_t : \mathcal{X} \rightarrow \mathbb{R}$;
- (3) the player suffers loss $f_t(\mathbf{x}_t)$, observes the information (loss) f_t and update:

$$\text{Update } A_t = A_{t-1} + \nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top$$

$$\text{Update } \mathbf{x}_{t+1} = \arg \min_{\mathbf{x} \in \mathcal{X}} \left\| \mathbf{x} - \left(\mathbf{x}_t - \frac{1}{\gamma} A_t^{-1} \nabla f_t(\mathbf{x}_t) \right) \right\|_{A_t}^2$$

ONS: In a View of Proximal Gradient

Convex Problem

Property: $f_t(\mathbf{x}) \geq f_t(\mathbf{y}) + \nabla f_t(\mathbf{y})^\top (\mathbf{x} - \mathbf{y})$

$$\text{OGD: } \mathbf{x}_{t+1} = \Pi_{\mathcal{X}} \left[\mathbf{x}_t - \frac{1}{\sqrt{t}} \nabla f_t(\mathbf{x}_t) \right]$$

Proximal type update:

$$\mathbf{x}_{t+1} = \arg \min_{\mathbf{x} \in \mathcal{X}} \langle \mathbf{x}, \nabla f_t(\mathbf{x}_t) \rangle + \frac{1}{2\eta_t} \|\mathbf{x} - \mathbf{x}_t\|_2^2$$

Exp-concave Problem

Property: $f_t(\mathbf{x}) \geq f_t(\mathbf{y}) + \nabla f_t(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) + \frac{\gamma}{2} \|\mathbf{x} - \mathbf{y}\|_{\nabla f_t(\mathbf{y}) \nabla f_t(\mathbf{y})^\top}^2$

$$\text{ONS: } A_t = A_{t-1} + \nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top$$

$$\mathbf{x}_{t+1} = \Pi_{\mathcal{X}}^{A_t} \left[\mathbf{x}_t - \frac{1}{\gamma} A_t^{-1} \nabla f_t(\mathbf{x}_t) \right]$$

Proximal type update:

$$\mathbf{x}_{t+1} = \arg \min_{\mathbf{x} \in \mathcal{X}} \langle \mathbf{x}, \nabla f_t(\mathbf{x}_t) \rangle + \frac{\gamma}{2} \|\mathbf{x} - \mathbf{x}_t\|_{A_t}^2$$

ONS: In a View of Proximal Gradient

Proof.

$$\begin{aligned}
 \mathbf{x}_{t+1} &= \Pi_{\mathcal{X}}^{A_t} \left[\mathbf{x}_t - \frac{A_t^{-1}}{\gamma} \mathbf{g}_t \right] \quad (\mathbf{g}_t \triangleq \nabla f_t(\mathbf{x}_t)) \\
 &= \arg \min_{\mathbf{x} \in \mathcal{X}} \left(\mathbf{x} - \mathbf{x}_t + \frac{A_t^{-1}}{\gamma} \mathbf{g}_t \right)^\top A_t \left(\mathbf{x} - \mathbf{x}_t + \frac{A_t^{-1}}{\gamma} \mathbf{g}_t \right) \\
 &= \arg \min_{\mathbf{x} \in \mathcal{X}} \left(\mathbf{x} - \mathbf{x}_t + \frac{A_t^{-1}}{\gamma} \mathbf{g}_t \right)^\top \left(A_t \mathbf{x} - A_t \mathbf{x}_t + \frac{\mathbf{g}_t}{\gamma} \right) \\
 &= \arg \min_{\mathbf{x} \in \mathcal{X}} (\mathbf{x} - \mathbf{x}_t)^\top A_t (\mathbf{x} - \mathbf{x}_t) + \cancel{(A_t^{-1})^\top \mathbf{g}_t^\top \mathbf{g}_t} \\
 &\quad + 2 \frac{\mathbf{g}_t^\top (\mathbf{x} - \mathbf{x}_t)}{\gamma} \quad (\text{constant}) \\
 &= \arg \min_{\mathbf{x} \in \mathcal{X}} \langle \mathbf{x}, \mathbf{g}_t \rangle + \frac{\gamma}{2} \|\mathbf{x} - \mathbf{x}_t\|_{A_t}^2
 \end{aligned}$$

Exp-concave Problem

Property: $f_t(\mathbf{x}) \geq f_t(\mathbf{y}) + \nabla f_t(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) + \frac{\gamma}{2} \|\mathbf{x} - \mathbf{y}\|_{\nabla f_t(\mathbf{y}) \nabla f_t(\mathbf{y})^\top}^2$

ONS: $A_t = A_{t-1} + \nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top$

$$\mathbf{x}_{t+1} = \Pi_{\mathcal{X}}^{A_t} \left[\mathbf{x}_t - \frac{1}{\gamma} A_t^{-1} \nabla f_t(\mathbf{x}_t) \right]$$

Proximal type update:

$$\mathbf{x}_{t+1} = \arg \min_{\mathbf{x} \in \mathcal{X}} \langle \mathbf{x}, \nabla f_t(\mathbf{x}_t) \rangle + \frac{\gamma}{2} \|\mathbf{x} - \mathbf{x}_t\|_{A_t}^2$$

ONS for Exp-concave Function

Theorem 5. *Under Assumptions 1, 2 and 3, for α -exp-concave online functions, the ONS algorithm with parameters $\gamma = \frac{1}{2} \min \left\{ \frac{1}{GD}, \alpha \right\}$ and $\varepsilon = \frac{1}{\gamma^2 D^2}$ (recall that the initial matrix is $A_0 = \varepsilon I_d$) guarantees*

$$\text{Regret}_T \leq \mathcal{O} \left(\left(\frac{1}{\alpha} + GD \right) d \log T \right) = \mathcal{O}(d \log T),$$

where d is the dimension of the feasible domain $\mathcal{X} \subseteq \mathbb{R}^d$.

Proof

$$\begin{aligned} \bullet A_t &= A_{t-1} + \nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top \\ \bullet \mathbf{x}_{t+1} &= \arg \min_{\mathbf{x} \in \mathcal{X}} \left\| \mathbf{x} - \left(\mathbf{x}_t - \frac{1}{\gamma} A_t^{-1} \mathbf{g}_t \right) \right\|_{A_t}^2 \end{aligned}$$

Extending *the first GD lemma* to *exp-concave case*:

Proof.

We use norm induced by A_t instead of 2-norm.

$$\begin{aligned} \|\mathbf{x}_{t+1} - \mathbf{u}\|_{A_t}^2 &= \left\| \Pi_{\mathcal{X}}^{A_t} \left[\mathbf{x}_t - \frac{1}{\gamma} A_t^{-1} \nabla f_t(\mathbf{x}_t) \right] - \mathbf{u} \right\|_{A_t}^2 && (\Pi_{\mathcal{X}}^A[\mathbf{y}] \triangleq \arg \min_{\mathbf{x} \in \mathcal{X}} \|\mathbf{x} - \mathbf{y}\|_A^2) \\ &\leq \left\| \mathbf{x}_t - \frac{1}{\gamma} A_t^{-1} \nabla f_t(\mathbf{x}_t) - \mathbf{u} \right\|_{A_t}^2 && (A_t \text{ is semidefinite matrix}) \\ &&& (\text{Pythagoras theorem}) \\ &= \left(\mathbf{x}_t - \frac{1}{\gamma} A_t^{-1} \nabla f_t(\mathbf{x}_t) - \mathbf{u} \right)^\top A_t \left(\mathbf{x}_t - \frac{1}{\gamma} A_t^{-1} \nabla f_t(\mathbf{x}_t) - \mathbf{u} \right) \\ &&& (\text{definition of } \|\cdot\|_{A_t}^2) \\ &= \left(\mathbf{x}_t - \mathbf{u} - \frac{1}{\gamma} A_t^{-1} \nabla f_t(\mathbf{x}_t) \right)^\top \left(A_t(\mathbf{x}_t - \mathbf{u}) - \frac{1}{\gamma} \nabla f_t(\mathbf{x}_t) \right) \end{aligned}$$

Proof

- $A_t = A_{t-1} + \nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top$
- $\mathbf{x}_{t+1} = \arg \min_{\mathbf{x} \in \mathcal{X}} \left\| \mathbf{x} - \left(\mathbf{x}_t - \frac{1}{\gamma} A_t^{-1} \mathbf{g}_t \right) \right\|_{A_t}^2$

Extending *the first GD lemma* to *exp-concave case*:

Proof.

$$\begin{aligned} \|\mathbf{x}_{t+1} - \mathbf{u}\|_{A_t}^2 &= \left(\mathbf{x}_t - \mathbf{u} - \frac{1}{\gamma} A_t^{-1} \nabla f_t(\mathbf{x}_t) \right)^\top \left(A_t (\mathbf{x}_t - \mathbf{u}) - \frac{1}{\gamma} \nabla f_t(\mathbf{x}_t) \right) \\ &= (\mathbf{x}_t - \mathbf{u})^\top A_t (\mathbf{x}_t - \mathbf{u}) - \frac{2}{\gamma} \nabla f_t(\mathbf{x}_t)^\top (\mathbf{x}_t - \mathbf{u}) + \frac{1}{\gamma^2} \nabla f_t(\mathbf{x}_t)^\top A_t^{-1} \nabla f_t(\mathbf{x}_t) \\ &\leq \|\mathbf{x}_t - \mathbf{u}\|_{A_t}^2 - \frac{2}{\gamma} (f_t(\mathbf{x}_t) - f_t(\mathbf{u})) + \frac{1}{\gamma^2} \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2 \\ &\quad - (\mathbf{x}_t - \mathbf{u})^\top \nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top (\mathbf{x}_t - \mathbf{u}) \end{aligned}$$

(Exp-concave: $f_t(\mathbf{x}) \geq f_t(\mathbf{y}) + \nabla f_t(\mathbf{y})^\top (\mathbf{x} - \mathbf{y}) + \frac{\gamma}{2} (\mathbf{x} - \mathbf{y})^\top \nabla f_t(\mathbf{y}) \nabla f_t(\mathbf{y})^\top (\mathbf{x} - \mathbf{y})$)

Proof

Proof. $\|\mathbf{x}_{t+1} - \mathbf{u}\|_{A_t}^2$

$$\leq \|\mathbf{x}_t - \mathbf{u}\|_{A_t}^2 - \frac{2}{\gamma} (f_t(\mathbf{x}_t) - f_t(\mathbf{u})) - \|\mathbf{x}_t - \mathbf{u}\|_{\nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top}^2 + \frac{1}{\gamma^2} \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2$$

$$\Leftrightarrow f_t(\mathbf{x}_t) - f_t(\mathbf{u}) \leq \frac{\gamma}{2} \|\mathbf{x}_t - \mathbf{u}\|_{A_t}^2 - \frac{\gamma}{2} \|\mathbf{x}_{t+1} - \mathbf{u}\|_{A_t}^2 - \frac{\gamma}{2} \|\mathbf{x}_t - \mathbf{u}\|_{\nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top}^2 + \frac{1}{2\gamma} \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2$$

(rearranging)

Summing from $t = 1$ to T , by telescoping:

$$\begin{aligned} \sum_{t=1}^T (f_t(\mathbf{x}_t) - f_t(\mathbf{u})) &\leq \frac{\gamma}{2} \sum_{t=1}^T \left(\|\mathbf{x}_t - \mathbf{u}\|_{A_t}^2 - \|\mathbf{x}_t - \mathbf{u}\|_{A_{t-1}}^2 \right) + \frac{1}{2\gamma} \sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2 \\ &\quad + \frac{\gamma}{2} \|\mathbf{x}_1 - \mathbf{u}\|_{A_0}^2 - \frac{\gamma}{2} \sum_{t=1}^T \|\mathbf{x}_t - \mathbf{u}\|_{\nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top}^2 \quad \text{cancellation} \\ &\leq \frac{\gamma}{2} \|\mathbf{x}_1 - \mathbf{u}\|_{A_0}^2 + \frac{1}{2\gamma} \sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2 \quad (A_t = A_{t-1} + \nabla f_t(\mathbf{x}_t) \nabla f_t(\mathbf{x}_t)^\top) \end{aligned}$$

Proof

Proof.
$$\sum_{t=1}^T (f_t(\mathbf{x}_t) - f_t(\mathbf{u})) \leq \frac{\gamma}{2} \|\mathbf{x}_1 - \mathbf{u}\|_{A_0}^2 + \frac{1}{2\gamma} \sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2$$

By the definition that $A_0 \triangleq \varepsilon I_d$, $\varepsilon = \frac{1}{\gamma^2 D^2}$ and the diameter $\|\mathbf{x}_1 - \mathbf{u}\|_2^2 \leq D^2$:

$$\begin{aligned} \sum_{t=1}^T (f_t(\mathbf{x}_t) - f_t(\mathbf{u})) &\leq \frac{\gamma}{2} (\mathbf{x}_1 - \mathbf{u})^\top A_0 (\mathbf{x}_1 - \mathbf{u}) + \frac{1}{2\gamma} \sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2 \\ &\leq \frac{1}{2\gamma} + \frac{1}{2\gamma} \sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2. \end{aligned}$$

Next, we bound the term $\sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2$.

Proof

Proof. Next, we bound the term $\sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2$.

Lemma 4 (Elliptical Potential Lemma). *For any sequence $\{X_1, \dots, X_T\} \in \mathbb{R}^{d \times T}$, suppose $U_0 = \lambda I$, $U_t = U_{t-1} + X_t X_t^\top$, and $\|X_t\|_2 \leq L$, then*

$$\sum_{t=1}^T \|X_t\|_{U_t^{-1}}^2 \leq d \log \left(1 + \frac{L^2 T}{\lambda d} \right)$$

Proof. $U_{t-1} = U_t - X_t X_t^\top = U_t^{\frac{1}{2}} \left(I - U_t^{-\frac{1}{2}} X_t X_t^\top U_t^{-\frac{1}{2}} \right) U_t^{\frac{1}{2}}$ (definition of U_t)

$$\det(U_{t-1}) = \det(U_t) \det \left(I - U_t^{-\frac{1}{2}} X_t X_t^\top U_t^{-\frac{1}{2}} \right) \quad (\text{determinant on both side})$$

Proof

Proof. Next, we bound the term $\sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2$.

Lemma 5. For any $\mathbf{v} \in \mathbb{R}^d$, we have

$$\det(I - \mathbf{v}\mathbf{v}^\top) = 1 - \|\mathbf{v}\|_2^2$$

Proof.

- (i) $(I - \mathbf{v}\mathbf{v}^\top) \mathbf{v} = (1 - \|\mathbf{v}\|_2^2) \mathbf{v}$, therefore, \mathbf{v} is its eigenvector with $(1 - \|\mathbf{v}\|_2^2)$ as eigenvalue;
- (ii) $(I - \mathbf{v}\mathbf{v}^\top) \mathbf{v}^\perp = \mathbf{v}^\perp$, therefore, $\mathbf{v}^\perp \perp \mathbf{v}$ is its eigenvector with 1 as the eigenvalue. □

Proof

Proof. Next, we bound the term $\sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2$.

Lemma 4 (Elliptical Potential Lemma). *For any sequence $\{X_1, \dots, X_T\} \in \mathbb{R}^{d \times T}$, suppose $U_0 = \lambda I$, $U_t = U_{t-1} + X_t X_t^\top$, and $\|X_t\|_2 \leq L$, then*

$$\sum_{t=1}^T \|X_t\|_{U_t^{-1}}^2 \leq d \log \left(1 + \frac{L^2 T}{\lambda d} \right)$$

Proof. $\det(U_{t-1}) = \det(U_t) \det \left(I - U_t^{-\frac{1}{2}} X_t X_t^\top U_t^{-\frac{1}{2}} \right) = \det(U_t) \left(1 - \left\| U_t^{-\frac{1}{2}} X_t \right\|_2^2 \right)$
(by Lemma 5)

$$\implies \|X_t\|_{U_t^{-1}}^2 = \left\| U_t^{-\frac{1}{2}} X_t \right\|_2^2 = 1 - \frac{\det(U_{t-1})}{\det(U_t)} \quad (\text{rearranging, } U \text{ is a symmetric matrix})$$

Proof

Proof. Next, we bound the term $\sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2$.

Lemma 4 (Elliptical Potential Lemma). *For any sequence $\{X_1, \dots, X_T\} \in \mathbb{R}^{d \times T}$, suppose $U_0 = \lambda I$, $U_t = U_{t-1} + X_t X_t^\top$, and $\|X_t\|_2 \leq L$, then*

$$\sum_{t=1}^T \|X_t\|_{U_t^{-1}}^2 \leq d \log \left(1 + \frac{L^2 T}{\lambda d} \right)$$

Proof.

$$\Rightarrow \sum_{t=1}^T X_t^\top U_t^{-1} X_t = \sum_{t=1}^T \left(1 - \frac{\det(U_{t-1})}{\det(U_t)} \right) \leq \sum_{t=1}^T \log \frac{\det(U_t)}{\det(U_{t-1})} \quad (\forall x > 0, 1 - x \leq -\log x)$$

$$= \log \frac{\det(U_T)}{\det(U_0)} = d \log \left(1 + \frac{L^2 T}{\lambda d} \right) \quad \square \quad \begin{array}{l} \text{Tr}(U_T) \leq \text{Tr}(U_0) + L^2 T = \lambda d + L^2 T \\ \Rightarrow \det(U_T) \leq (\lambda + L^2 T/d)^d \end{array}$$

Proof

Proof. Next, we bound the term $\sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2$.

Lemma 4 (Elliptical Potential Lemma). *For any sequence $\{X_1, \dots, X_T\} \in \mathbb{R}^{d \times T}$, suppose $U_0 = \lambda I$, $U_t = U_{t-1} + X_t X_t^\top$, and $\|X_t\|_2 \leq L$, then*

$$\sum_{t=1}^T \|X_t\|_{U_t^{-1}}^2 \leq d \log \left(1 + \frac{L^2 T}{\lambda d} \right)$$

Therefore, by Lemma 4, we have

$$\sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2 \leq d \log \left(1 + \frac{D^2 T}{\varepsilon d} \right).$$

Proof

Proof. To conclude,

$$\sum_{t=1}^T (f_t(\mathbf{x}_t) - f_t(\mathbf{u})) \leq \underbrace{\frac{\gamma}{2} \|\mathbf{x}_1 - \mathbf{u}\|_{A_0}^2}_{\leq \frac{1}{2\gamma} \text{ (bounded domain)}} + \underbrace{\frac{1}{2\gamma} \sum_{t=1}^T \|\nabla f_t(\mathbf{x}_t)\|_{A_t^{-1}}^2}_{\leq \frac{d}{2\gamma} \log \left(1 + \frac{D^2 T}{\varepsilon d}\right) \text{ (elliptical potential lemma)}}.$$

Recall that $\gamma = \frac{1}{2} \min \left\{ \frac{1}{GD}, \alpha \right\}$ and $\varepsilon = \frac{1}{\gamma^2 D^2}$,

$$\text{Regret}_T \leq \mathcal{O} \left(\left(\frac{1}{\alpha} + GD \right) d \log T \right). \quad \square$$

Lower Bounds

- A natural question: whether previous regret can be improved?
- Lower bound argument:

minimax bound: smallest possible worst-case regret of any algorithm:

$$\min_A \max_{\ell_1, \dots, \ell_T} \text{Regret}_T$$

Theorem 7 (Lower Bound for OCO). *Any algorithm for online convex optimization incurs $\Omega(DG\sqrt{T})$ regret in the worst case. This is true even if the cost functions are generated from a fixed stationary distribution.*

Lower Bounds

Theorem 7 (Lower Bound for OCO). *Any algorithm for online convex optimization incurs $\Omega(DG\sqrt{T})$ regret in the worst case. This is true even if the cost functions are generated from a fixed stationary distribution.*

Proof Sketch.

Construct a **'hard'** environment:

- Binary classification, loss functions in each iteration are chosen at random
- Similar results can be obtained for strongly convex and exp-concave cases

Comparison

	Algorithm	Upper Bound	Lower Bound
Convex	OGD	$\mathcal{O}(\sqrt{T})$	$\Omega(\sqrt{T})$
σ -Strongly Convex	OGD	$\mathcal{O}\left(\frac{\log T}{\sigma}\right)$	$\Omega\left(\frac{\log T}{\sigma}\right)$
α -Exp-concave	ONS	$\mathcal{O}\left(\frac{d \log T}{\alpha}\right)$	$\Omega\left(\frac{d \log T}{\alpha}\right)$

Back to Exp-concave Learning

- Universal Portfolio Selection



Algorithm	Regret	Runtime (per round)
Universal Portfolios	$d \log(T)$	$d^4 T^{14}$
Online Gradient Descent	$G_2 \sqrt{T}$	d
Exponentiated Gradient	$G_\infty \sqrt{T \log(d)}$	d
Online Newton Step (ONS)	$G_\infty d \log(T)$	$d^2 + \text{generalized projection on } \Delta_d$
Soft-Bayes	$\sqrt{dT \log(d)}$	d
Ada-BARRONS	$d^2 \log^4(T)$	$d^{2.5} T$
BISONS	$d^2 \log^2(T)$	$\text{poly}(d)$
AdaMix+DONS	$d^2 \log^5(T)$	d^3
VB-FTRL	$d \log(T)$	$d^2 T$

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Open Problem: Fast and Optimal Online Portfolio Selection

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Abstract

Online portfolio selection has received much attention in the COLT community since its introduction by Cover, but all state-of-the-art methods fall short in at least one of the following ways: they are either i) computationally infeasible; or ii) they do not guarantee optimal regret; or iii) they assume the gradients are bounded, which is unnecessary and cannot be guaranteed. We are interested in a natural follow-the-regularized-leader (FTRL) approach based on the log barrier regularizer, which is computationally feasible. The open problem we put before the community is to formally prove whether this approach achieves the optimal regret. Resolving this question will likely lead to new techniques to analyse FTRL algorithms. There are also interesting technical connections to self-concordance, which has previously been used in the context of bandit convex optimization.

1. Introduction

Online portfolio selection (Cover, 1991) may be viewed as an instance of online convex optimization (OCO) (Hazan et al., 2016): in each of $t = 1, \dots, T$ rounds, a learner has to make a prediction w_t in a convex domain \mathcal{W} before observing a convex loss function $f_t: \mathcal{W} \rightarrow \mathbb{R}$. The goal is to obtain a guaranteed bound on the regret $\text{Regret}_T = \sum_{t=1}^T f_t(w_t) - \min_{w \in \mathcal{W}} \sum_{t=1}^T f_t(w)$ that holds for any possible sequence of loss functions f_t . Online portfolio selection corresponds to the special case that the domain $\mathcal{W} = \{w \in \mathbb{R}_+^d \mid \sum_{i=1}^d w_i = 1\}$ is the probability simplex and the loss functions are restricted to be of the form $f_t(w) = -\ln(w^T x_t)$ for vectors $x_t \in \mathbb{R}_+^d$. It was introduced by Cover (1991) with the interpretation that $x_{t,i}$ represents the factor by which the value of an asset $i \in \{1, \dots, d\}$ grows in round t and $w_{t,i}$ represents the fraction of our capital we re-invest in asset i in round t . The factor by which our initial capital grows over T rounds then becomes $\prod_{t=1}^T w_t^T x_t = e^{-\sum_{t=1}^T f_t(w_t)}$. An alternative interpretation in terms of mixture learning is given by Oseanu et al. (2017).

For an extensive survey of online portfolio selection we refer to Li and Hoi (2014). Here we review only the results that are most relevant to our open problem. Cover (1991); Cover and Ordentlich (1996) show that the best possible guarantee on the regret is of order $\text{Regret}_T = O(d \ln T)$ and that this is achieved by choosing w_{t+1} as the mean of a continuous exponential weights distribution $dP_{t+1}(w) \propto e^{-\sum_{s=1}^t f_s(w)} d\pi(w)$ with Dirichlet-prior π (and learning rate $\eta = 1$). Unfortunately, this approach has a run-time of order $O(T^d)$, which scales exponentially in the number

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[COLT 2020 Open Problem]

⇒ still an important open problem: **efficiency and optimality**

Connection with Offline Learning

- Application to Stochastic Optimization
- Online-to-Batch Conversion

Application to Stochastic Optimization

- Consider the following *convex optimization* problem:

$$\min_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x})$$

- Stochastic optimization method

Computational oracle: only access *noisy* gradient oracle, namely, $\mathbf{g}(\mathbf{x})$, such that

$$\mathbb{E}[\mathbf{g}(\mathbf{x})] = \nabla f(\mathbf{x}), \text{ and } \mathbb{E}[\|\mathbf{g}(\mathbf{x})\|^2] \leq G^2$$

for some $G > 0$.

Example (large-scale opt.). Given dataset $S = \{(\mathbf{x}_1, y_1), \dots, (\mathbf{x}_m, y_m)\}$, ERM optimizes

$$\min_{h \in \mathcal{H}} \sum_{i=1}^m \ell(h(\mathbf{x}_i), y_i) \quad \Rightarrow$$

full gradient computation requires a pass of *all data*

stochastic method only uses a *mini batch* at each round

Stochastic Gradient Descent

- Consider the following *convex optimization* problem:

$$\min_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x})$$

Algorithm 2 Stochastic Gradient Descent

Input: noisy gradient oracle $g(\cdot)$, step sizes $\{\eta_t\}$

1: **for** $t = 1, \dots, T$ **do**

2: Obtain noisy gradient $g(\mathbf{x}_t)$

3: Update the model $\mathbf{x}_{t+1} = \Pi_{\mathcal{X}} [\mathbf{x}_t - \eta_t g(\mathbf{x}_t)]$

4: **end for**

5: **return** $\bar{\mathbf{x}}_T = \frac{1}{T} \sum_{t=1}^T \mathbf{x}_t$

$$\mathbb{E}[g(\mathbf{x})] = \nabla f(\mathbf{x})$$

$$\mathbb{E}[\|g(\mathbf{x})\|^2] \leq G^2$$

Stochastic Gradient Descent

Theorem 7 (Convergence of SGD). *Suppose the domain $\mathcal{X} \subseteq \mathbb{R}^d$ has a diameter $D > 0$, and the noisy gradient oracle is unbiased and variance bounded by G^2 . SGD with step size $\eta_t = \frac{D}{G\sqrt{t}}$ guarantees*

$$\mathbb{E}[f(\bar{\mathbf{x}}_T)] - \min_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x}) \leq \frac{3GD}{2\sqrt{T}} = \mathcal{O}\left(\frac{1}{\sqrt{T}}\right),$$

where $\bar{\mathbf{x}}_T \triangleq \frac{1}{T} \sum_{t=1}^T \mathbf{x}_t$ is the output of the SGD algorithm.

Proof of SGD Convergence

Proof. First, we rephrase SGD from lens of *online convex optimization*.

To see this, we define *linear function* $h_t(\mathbf{x}) \triangleq \mathbf{g}_t^\top \mathbf{x}$, where $\mathbf{g}_t = \mathbf{g}(\mathbf{x}_t)$.

Claim: deploying OGD over the online functions $\{h_t(\mathbf{x})\}$ is equivalent to SGD proposed in the earlier page.

$$\begin{aligned} \text{OGD: } \mathbf{x}_{t+1} &= \Pi_{\mathcal{X}}[\mathbf{x}_t - \eta_t \nabla h_t(\mathbf{x}_t)] \\ &= \Pi_{\mathcal{X}}[\mathbf{x}_t - \eta_t \mathbf{g}(\mathbf{x}_t)] \end{aligned}$$

Algorithm 2 Stochastic Gradient Descent

Input: noisy gradient oracle $\mathbf{g}(\cdot)$, step sizes $\{\eta_t\}$

- 1: **for** $t = 1, \dots, T$ **do**
 - 2: Obtain noisy gradient $\mathbf{g}(\mathbf{x}_t)$
 - 3: Update the model $\mathbf{x}_{t+1} = \Pi_{\mathcal{X}}[\mathbf{x}_t - \eta_t \mathbf{g}(\mathbf{x}_t)]$
 - 4: **end for**
 - 5: **return** $\bar{\mathbf{x}}_T = \frac{1}{T} \sum_{t=1}^T \mathbf{x}_t$
-

Proof of SGD Convergence

Proof.

$$\begin{aligned} \mathbb{E}[f(\bar{\mathbf{x}}_T)] - f(\mathbf{x}^*) &\leq \mathbb{E}\left[\frac{1}{T} \sum_{t=1}^T f(\mathbf{x}_t)\right] - f(\mathbf{x}^*) && (\mathbf{x}^* = \arg \min_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x})) \\ &&& \text{(Jensen's inequality)} \\ &\leq \frac{1}{T} \mathbb{E}\left[\sum_{t=1}^T \nabla f(\mathbf{x}_t)^\top (\mathbf{x}_t - \mathbf{x}^*)\right] && \text{(convexity)} \\ &= \frac{1}{T} \mathbb{E}\left[\sum_{t=1}^T \mathbf{g}_t^\top (\mathbf{x}_t - \mathbf{x}^*)\right] \end{aligned}$$

Theorem 3 (Regret bound for OGD). Under Assumption 1, 2 and 3, online gradient descent (OGD) with step sizes $\eta_t = \frac{D}{G\sqrt{t}}$ for $t \in [T]$ guarantees:

$$\text{Regret}_T = \sum_{t=1}^T f_t(\mathbf{x}_t) - \min_{\mathbf{x} \in \mathcal{X}} \sum_{t=1}^T f_t(\mathbf{x}) \leq \frac{3}{2}GD\sqrt{T}.$$

Proof:

Define $\boldsymbol{\xi}_t \triangleq \nabla f(\mathbf{x}_t) - \mathbf{g}_t$. We know $\mathbb{E}[\boldsymbol{\xi}_t \mid \mathbf{x}_t] = 0$.

We have $\mathbb{E}[\nabla f(\mathbf{x}_t)^\top \mathbf{x}_t] = \mathbb{E}[\boldsymbol{\xi}_t^\top \mathbf{x}_t] + \mathbb{E}[\mathbf{g}_t^\top \mathbf{x}_t]$

$$\mathbb{E}[\boldsymbol{\xi}_t^\top \mathbf{x}_t] = \mathbb{E}[\mathbb{E}[\boldsymbol{\xi}_t^\top \mathbf{x}_t \mid \mathbf{x}_t]] = \mathbb{E}[\mathbb{E}[\mathbb{E}[\boldsymbol{\xi}_t \mid \mathbf{x}_t]^\top \mathbf{x}_t \mid \mathbf{x}_t]] = 0.$$

Therefore, we have proved that $\mathbb{E}[\nabla f(\mathbf{x}_t)^\top \mathbf{x}_t] = \mathbb{E}[\mathbf{g}_t^\top \mathbf{x}_t]$.

Similar argument shows $\mathbb{E}[\nabla f(\mathbf{x}_t)^\top \mathbf{x}] = \mathbb{E}[\mathbf{g}_t^\top \mathbf{x}]$ for any fixed \mathbf{x} . \square

Proof of SGD Convergence

Proof.

$$\mathbb{E} [f(\bar{\mathbf{x}}_T)] - f(\mathbf{x}^*) \leq \mathbb{E} \left[\frac{1}{T} \sum_{t=1}^T f(\mathbf{x}_t) \right] - f(\mathbf{x}^*)$$

($\mathbf{x}^* = \arg \min_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x})$)

(Jensen's inequality)

$$\leq \frac{1}{T} \mathbb{E} \left[\sum_{t=1}^T \nabla f(\mathbf{x}_t)^\top (\mathbf{x}_t - \mathbf{x}^*) \right]$$

(convexity)

$$= \frac{1}{T} \mathbb{E} \left[\sum_{t=1}^T \mathbf{g}_t^\top (\mathbf{x}_t - \mathbf{x}^*) \right]$$

$$= \frac{1}{T} \mathbb{E} \left[\sum_{t=1}^T h_t(\mathbf{x}_t) - h_t(\mathbf{x}^*) \right]$$

(definition of $f_t(\cdot)$)

$$\leq \frac{\text{Regret}_T}{T}$$

(SGD = OGD over $\{f_t(\cdot)\}$)

(regret bound of OGD)

$$\leq \frac{3GD}{2\sqrt{T}}$$

□

(regret of OGD algorithm)

Theorem 3 (Regret bound for OGD). Under Assumption 1, 2 and 3, online gradient descent (OGD) with step sizes $\eta_t = \frac{D}{G\sqrt{t}}$ for $t \in [T]$ guarantees:

$$\text{Regret}_T = \sum_{t=1}^T f_t(\mathbf{x}_t) - \min_{\mathbf{x} \in \mathcal{X}} \sum_{t=1}^T f_t(\mathbf{x}) \leq \frac{3}{2}GD\sqrt{T}.$$

Stochastic Gradient Descent

Theorem 7 (Convergence of SGD). *Suppose the domain $\mathcal{X} \subseteq \mathbb{R}^d$ has a diameter $D > 0$, and the noisy gradient oracle is unbiased and variance bounded by G^2 . SGD with step size $\eta_t = \frac{D}{G\sqrt{t}}$ guarantees*

$$\mathbb{E}[f(\bar{\mathbf{x}}_T)] \leq \min_{\mathbf{x} \in \mathcal{X}} f(\mathbf{x}) + \frac{3GD}{2\sqrt{T}},$$

where $\bar{\mathbf{x}}_T \triangleq \frac{1}{T} \sum_{t=1}^T \mathbf{x}_t$ is the output of the SGD algorithm.

- We define the linear function $h_t(\mathbf{x}) \triangleq \mathbf{g}_t^\top \mathbf{x} = \mathbf{g}(\mathbf{x}_t)^\top \mathbf{x}$ and run OGD on $\{h_t\}_{t=1}^T$.
- Note that function h_t *depends* on the decision \mathbf{x}_t , which actually reveals that OGD regret can hold even against *adaptive adversary*.

More bits of OGD

- We define the linear function $h_t(\mathbf{x}) \triangleq \mathbf{g}_t^\top \mathbf{x} = \mathbf{g}(\mathbf{x}_t)^\top \mathbf{x}$ and run OGD on $\{h_t\}_{t=1}^T$.
- Note that function h_t *depends* on the decision \mathbf{x}_t , which actually reveals that OGD regret can hold even against *adaptive adversary*.

At each round $t = 1, 2, \dots$

- (1) the player first picks a model $\mathbf{x}_t \in \mathcal{X}$;
- (2) and *simultaneously* environments pick an online function $f_t : \mathcal{X} \rightarrow \mathbb{R}$;
- (3) the player suffers loss $f_t(\mathbf{x}_t)$, observes some information about f_t and updates the model.

oblivious adversary



examination

adaptive adversary



interview

The “simultaneously” requirement can be sometimes not necessary!

OGD for full-info OCO can handle the case when online functions *depend on \mathbf{x}_t* !

Online-to-Batch Conversion

- An alternative way to solve statistic learning:
 - use the data in a sequential way
 - run any online algorithm minimizing the regret
 - return the final model as the average

Algorithm 1 Online-to-Batch Conversion

Input: Data $\{(\mathbf{x}_1, y_1), \dots, (\mathbf{x}_T, y_T)\}$ **i.i.d.** sampled from the distribution \mathcal{D} , a **bounded** loss function $\ell : \mathcal{Y} \times \mathcal{Y} \rightarrow [0, 1]$, an online learning algorithm \mathcal{A}

- 1: **for** $t = 1, \dots, T$ **do**
 - 2: let \mathbf{w}_t be the output of algorithm \mathcal{A} for this round
 - 3: Feed algorithm \mathcal{A} with loss function $f_t(\mathbf{w}) = \ell(\mathbf{w}^\top \mathbf{x}_t), y_t$
 - 4: **end for**
 - 5: **return** $\bar{\mathbf{w}} = \frac{1}{T} \sum_{t=1}^T \mathbf{w}_t$
-

Online-to-Batch Conversion

Theorem 2 (Online-to-Batch Conversion). *If the risk $R(\mathbf{w})$ is convex w.r.t. \mathbf{w} with a **bounded** loss function $\ell : \mathcal{Y} \times \mathcal{Y} \rightarrow [0, 1]$, and the data $\{(\mathbf{x}_1, y_1), \dots, (\mathbf{x}_T, y_T)\}$ are **i.i.d.** sampled from the distribution \mathcal{D} , then with probability at least $1 - \delta$, the excess risk of the output of Algorithm 1 satisfies*

$$R(\bar{\mathbf{w}}) - \min_{\mathbf{w} \in \mathcal{W}} R(\mathbf{w}) \leq \frac{\text{Regret}_T}{T} + 2\sqrt{\frac{2 \ln(2/\delta)}{T}}$$

where $R(\mathbf{w}) \triangleq \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}}[\ell(h(\mathbf{w}; \mathbf{x}), y)]$ is the expected risk, and $\text{Regret}_T \triangleq \sum_{t=1}^T f_t(\mathbf{w}_t) - \min_{\mathbf{w} \in \mathcal{W}} \sum_{t=1}^T f_t(\mathbf{w})$ is the regret of the online algorithm \mathcal{A} after T rounds.

Concentration Inequalities

Lemma 1 (Hoeffding's inequality). Let $X_1, \dots, X_T \in [-B, B]$ for some $B > 0$ be independent random variables such that $\mathbb{E}[X_t] = 0$ for all $t \in [T]$, then for all $\delta \in (0, 1)$,

$$\Pr \left[\sum_{t=1}^T X_t \geq B \sqrt{2T \ln \frac{1}{\delta}} \right] \leq \delta$$

Lemma 2 (Azuma's inequality). Let $X_1, \dots, X_T \in [-B, B]$ for some $B > 0$ be a martingale difference sequence (i.e., $\forall t \in [T], \mathbb{E}[X_t | X_{t-1}, \dots, X_1] = 0$), then $\forall \delta > 0$,

$$\Pr \left[\sum_{t=1}^T X_t \geq B \sqrt{2T \ln \frac{1}{\delta}} \right] \leq \delta$$

Online-to-Batch Conversion

Theorem 2 (Online-to-Batch Conversion). *If the risk $R(\mathbf{w})$ is convex w.r.t. \mathbf{w} with a **bounded** loss function $\ell : \mathcal{Y} \times \mathcal{Y} \rightarrow [0, 1]$, and the data $\{(\mathbf{x}_1, y_1), \dots, (\mathbf{x}_T, y_T)\}$ are **i.i.d.** sampled from the distribution \mathcal{D} , then with probability at least $1 - \delta$, the excess risk of the output of Algorithm 1 satisfies*

$$R(\bar{\mathbf{w}}) - \min_{\mathbf{w} \in \mathcal{W}} R(\mathbf{w}) \leq \frac{\text{Regret}_T}{T} + 2\sqrt{\frac{2 \ln(2/\delta)}{T}}$$

where $R(\mathbf{w}) \triangleq \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}}[\ell(h(\mathbf{w}; \mathbf{x}), y)]$ is the expected risk, and $\text{Regret}_T \triangleq \sum_{t=1}^T f_t(\mathbf{w}_t) - \min_{\mathbf{w} \in \mathcal{W}} \sum_{t=1}^T f_t(\mathbf{w})$ is the regret of the online algorithm \mathcal{A} after T rounds.

Proof Sketch.

$$R(\hat{\mathbf{w}}) \stackrel{\text{Jensen's inequality}}{\leq} \frac{1}{T} \sum_{t=1}^T R(\mathbf{w}_t) \stackrel{\text{Azuma's inequality}}{\leq} \frac{1}{T} \sum_{t=1}^T f_t(\mathbf{w}_t) + \sqrt{\frac{2 \ln(2/\delta)}{T}}$$

$$R(\mathbf{w}^*) + \sqrt{\frac{2 \ln(2/\delta)}{T}} \stackrel{\text{Hoeffding's inequality}}{\geq} \frac{1}{T} \sum_{t=1}^T f_t(\mathbf{w}^*) \geq \frac{1}{T} \min_{\mathbf{w} \in \mathcal{W}} \sum_{t=1}^T f_t(\mathbf{w})$$

Online-to-Batch Conversion

Proof.

$$\begin{aligned} R(\hat{\mathbf{w}}) &= \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}} [\ell(h(\hat{\mathbf{w}}; \mathbf{x}), y)] \\ &= \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}} \left[\ell\left(h\left(\frac{1}{T} \sum_{t=1}^T \mathbf{w}_t; \mathbf{x}\right), y\right) \right] \\ &\leq \frac{1}{T} \sum_{t=1}^T \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}} [\ell(h(\mathbf{w}_t; \mathbf{x}), y)] && \text{(Jensen's inequality)} \\ &= \frac{1}{T} \sum_{t=1}^T f_t(\mathbf{w}_t) + \sqrt{\frac{2 \ln(2/\delta)}{T}} && \begin{array}{l} \text{(Azuma's inequality} \\ \text{with } X_t = \mathbb{E}_{(\mathbf{x}, y) \sim \mathcal{D}} [\ell(h(\mathbf{w}_t; \mathbf{x}), y)] - f_t(\mathbf{w}_t)) \end{array} \end{aligned}$$

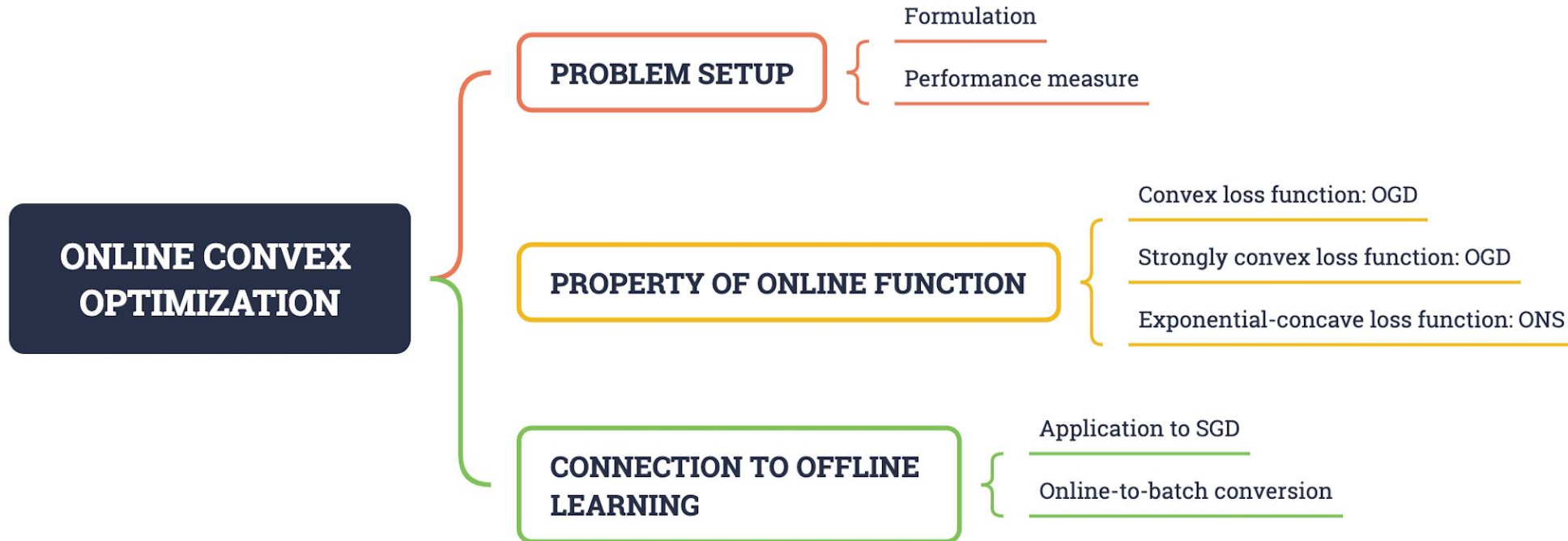
Online-to-Batch Conversion

Proof.
$$R(\hat{\mathbf{w}}) \leq \frac{1}{T} \sum_{t=1}^T f_t(\mathbf{w}_t) + \sqrt{\frac{2 \ln(2/\delta)}{T}}$$

$$= \min_{\mathbf{w} \in \mathcal{W}} \frac{1}{T} \sum_{t=1}^T f_t(\mathbf{w}) + \frac{\text{Regret}_T}{T} + \sqrt{\frac{2 \ln(2/\delta)}{T}} \quad (\text{definition of regret})$$
$$\leq \frac{1}{T} \sum_{t=1}^T f_t(\mathbf{w}^*) + \frac{\text{Regret}_T}{T} + \sqrt{\frac{2 \ln(2/\delta)}{T}}$$
$$\leq R(\mathbf{w}^*) + \frac{\text{Regret}_T}{T} + 2\sqrt{\frac{2 \ln(2/\delta)}{T}} \quad (\text{Hoeffding's inequality with } X_t = f_t(\mathbf{w}^*) - R(\mathbf{w}^*))$$

□

Summary



Q & A

Thanks!